



# SMART CONTRACT AUDIT REPORT

for

## RockX ETH Staking



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June 12, 2022

## Document Properties

Client	RockX
Title	Smart Contract Audit Report
Target	RockX ETH Staking
Version	1.0
Author	Xuxian Jiang
Auditors	Xiaotao Wu, Xuxian Jiang
Reviewed by	Xiaomi Huang
Approved by	Xuxian Jiang
Classification	Public

## Version Info

Version	Date	Author(s)	Description
1.0	June 12, 2022	Xuxian Jiang	Final Release
1.0-rc1	June 10, 2022	Xuxian Jiang	Release Candidate #1

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# 1 | Introduction

Given the opportunity to review the design document and related smart contract source code of the staking support of RockX, we outline in the report our systematic approach to evaluate potential security issues in the smart contract implementation, expose possible semantic inconsistencies between smart contract code and design document, and provide additional suggestions or recommendations for improvement. Our results show that the given version of smart contracts can be further improved due to the presence of several issues related to either security or performance. This document outlines our audit results.

## 1.1 About RockX

RockX is a blockchain fintech company that helps our customers embrace Web 3.0 effortlessly through the development of innovative products and infrastructure. It also strives to enable institutions and disruptors in the financial and internet sectors to gain seamless access to blockchain data, crypto yield products and best-in-class key management solutions in a sustainable way. This audit covers the staking support for ETH 2.0 in allowing users to deposit any number of ethers to the staking contract, and get back equivalent value of uniETH token (decided by real-time exchange ratio). The basic information of the audited protocol is as follows:

Table 1.1: Basic Information of The RockX ETH Staking

Item	Description
Name	RockX
Website	<a href="https://www.rockx.com/">https://www.rockx.com/</a>
Type	EVM Smart Contract
Platform	Solidity
Audit Method	Whitebox
Latest Audit Report	June 12, 2022

In the following, we show the Git repository of reviewed files and the commit hash values used in this audit.

- <https://github.com/RockX-SG/stake.git> (735f12c)

And here is the commit ID after fixes for the issues found in the audit have been checked in:

- <https://github.com/RockX-SG/stake.git> (f1d0bef)

## 1.2 About PeckShield

PeckShield Inc. [11] is a leading blockchain security company with the goal of elevating the security, privacy, and usability of current blockchain ecosystems by offering top-notch, industry-leading services and products (including the service of smart contract auditing). We are reachable at Telegram (<https://t.me/peckshield>), Twitter (<http://twitter.com/peckshield>), or Email ([contact@peckshield.com](mailto:contact@peckshield.com)).

Table 1.2: Vulnerability Severity Classification

Impact	High	Critical	High	Medium
	Medium	High	Medium	Low
	Low	Medium	Low	Low
		High	Medium	Low
		Likelihood		

## 1.3 Methodology

To standardize the evaluation, we define the following terminology based on OWASP Risk Rating Methodology [10]:

- Likelihood represents how likely a particular vulnerability is to be uncovered and exploited in the wild;
- Impact measures the technical loss and business damage of a successful attack;
- Severity demonstrates the overall criticality of the risk.

Likelihood and impact are categorized into three ratings: *H*, *M* and *L*, i.e., *high*, *medium* and *low* respectively. Severity is determined by likelihood and impact and can be classified into four categories accordingly, i.e., *Critical*, *High*, *Medium*, *Low* shown in Table 1.2.

Table 1.3: The Full List of Check Items

Category	Check Item
<b>Basic Coding Bugs</b>	Constructor Mismatch
	Ownership Takeover
	Redundant Fallback Function
	Overflows & Underflows
	Reentrancy
	Money-Giving Bug
	Blackhole
	Unauthorized Self-Destruct
	Revert DoS
	Unchecked External Call
	Gasless Send
	Send Instead Of Transfer
	Costly Loop
	(Unsafe) Use Of Untrusted Libraries
	(Unsafe) Use Of Predictable Variables
	Transaction Ordering Dependence
Deprecated Uses	
<b>Semantic Consistency Checks</b>	Semantic Consistency Checks
<b>Advanced DeFi Scrutiny</b>	Business Logics Review
	Functionality Checks
	Authentication Management
	Access Control & Authorization
	Oracle Security
	Digital Asset Escrow
	Kill-Switch Mechanism
	Operation Trails & Event Generation
	ERC20 Idiosyncrasies Handling
	Frontend-Contract Integration
	Deployment Consistency
	Holistic Risk Management
<b>Additional Recommendations</b>	Avoiding Use of Variadic Byte Array
	Using Fixed Compiler Version
	Making Visibility Level Explicit
	Making Type Inference Explicit
	Adhering To Function Declaration Strictly
Following Other Best Practices	

To evaluate the risk, we go through a list of check items and each would be labeled with a severity category. For one check item, if our tool or analysis does not identify any issue, the contract is considered safe regarding the check item. For any discovered issue, we might further deploy contracts on our private testnet and run tests to confirm the findings. If necessary, we would additionally build a PoC to demonstrate the possibility of exploitation. The concrete list of check items is shown in Table 1.3.

In particular, we perform the audit according to the following procedure:

- Basic Coding Bugs: We first statically analyze given smart contracts with our proprietary static code analyzer for known coding bugs, and then manually verify (reject or confirm) all the issues found by our tool.
- Semantic Consistency Checks: We then manually check the logic of implemented smart contracts and compare with the description in the white paper.
- Advanced DeFi Scrutiny: We further review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.
- Additional Recommendations: We also provide additional suggestions regarding the coding and development of smart contracts from the perspective of proven programming practices.

To better describe each issue we identified, we categorize the findings with Common Weakness Enumeration (CWE-699) [9], which is a community-developed list of software weakness types to better delineate and organize weaknesses around concepts frequently encountered in software development. Though some categories used in CWE-699 may not be relevant in smart contracts, we use the CWE categories in Table 1.4 to classify our findings.

## 1.4 Disclaimer

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Note that this security audit is not designed to replace functional tests required before any software release, and does not give any warranties on finding all possible security issues of the given smart contract(s) or blockchain software, i.e., the evaluation result does not guarantee the nonexistence of any further findings of security issues. As one audit-based assessment cannot be considered comprehensive, we always recommend proceeding with several independent audits and a public bug bounty program to ensure the security of smart contract(s). Last but not least, this security audit should not be used as investment advice.

Table 1.4: Common Weakness Enumeration (CWE) Classifications Used in This Audit

Category	Summary
<b>Configuration</b>	Weaknesses in this category are typically introduced during the configuration of the software.
<b>Data Processing Issues</b>	Weaknesses in this category are typically found in functionality that processes data.
<b>Numeric Errors</b>	Weaknesses in this category are related to improper calculation or conversion of numbers.
<b>Security Features</b>	Weaknesses in this category are concerned with topics like authentication, access control, confidentiality, cryptography, and privilege management. (Software security is not security software.)
<b>Time and State</b>	Weaknesses in this category are related to the improper management of time and state in an environment that supports simultaneous or near-simultaneous computation by multiple systems, processes, or threads.
<b>Error Conditions, Return Values, Status Codes</b>	Weaknesses in this category include weaknesses that occur if a function does not generate the correct return/status code, or if the application does not handle all possible return/status codes that could be generated by a function.
<b>Resource Management</b>	Weaknesses in this category are related to improper management of system resources.
<b>Behavioral Issues</b>	Weaknesses in this category are related to unexpected behaviors from code that an application uses.
<b>Business Logics</b>	Weaknesses in this category identify some of the underlying problems that commonly allow attackers to manipulate the business logic of an application. Errors in business logic can be devastating to an entire application.
<b>Initialization and Cleanup</b>	Weaknesses in this category occur in behaviors that are used for initialization and breakdown.
<b>Arguments and Parameters</b>	Weaknesses in this category are related to improper use of arguments or parameters within function calls.
<b>Expression Issues</b>	Weaknesses in this category are related to incorrectly written expressions within code.
<b>Coding Practices</b>	Weaknesses in this category are related to coding practices that are deemed unsafe and increase the chances that an exploitable vulnerability will be present in the application. They may not directly introduce a vulnerability, but indicate the product has not been carefully developed or maintained.

## 2 | Findings

### 2.1 Summary

Here is a summary of our findings after analyzing the implementation of the `staking` support in `RockX`. During the first phase of our audit, we study the smart contract source code and run our in-house static code analyzer through the codebase. The purpose here is to statically identify known coding bugs, and then manually verify (reject or confirm) issues reported by our tool. We further manually review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.

Severity	# of Findings	
Critical	0	
High	1	■
Medium	1	■
Low	2	■ ■
Informational	1	■
Total	5	

We have so far identified a list of potential issues: some of them involve subtle corner cases that might not be previously thought of, while others refer to unusual interactions among multiple contracts. For each uncovered issue, we have therefore developed test cases for reasoning, reproduction, and/or verification. After further analysis and internal discussion, we determined a few issues of varying severities that need to be brought up and paid more attention to, which are categorized in the above table. More information can be found in the next subsection, and the detailed discussions of each of them are in [Section 3](#).

## 2.2 Key Findings

Overall, these smart contracts are well-designed and engineered, though the implementation can be improved by resolving the identified issues (shown in Table 2.1), including 1 high-severity vulnerability, 1 medium-severity vulnerability, 2 low-severity vulnerabilities, and 1 informational suggestion.

Table 2.1: Key RockX ETH Staking Audit Findings

ID	Severity	Title	Category	Status
PVE-001	High	<a href="#">Incorrect Validator Replacement Logic in replaceValidator()</a>	Business Logic	Resolved
PVE-002	Low	<a href="#">Suggested Adherence Of The Checks-Effects-Interactions Pattern</a>	Time and State	Resolved
PVE-003	Low	<a href="#">Proper minToMint Enforcement in mint()</a>	Business Logic	Resolved
PVE-004	Informational	<a href="#">Proper Event Emission in mint()</a>	Status Codes	Resolved
PVE-005	Medium	<a href="#">Trust Issue of Admin Keys</a>	Security Features	Mitigated

Besides recommending specific countermeasures to mitigate these issues, we also emphasize that it is always important to develop necessary risk-control mechanisms and make contingency plans, which may need to be exercised before the mainnet deployment. The risk-control mechanisms need to kick in at the very moment when the contracts are being deployed in mainnet. Please refer to Section 3 for details.

## 3 | Detailed Results

### 3.1 Incorrect Validator Replacement Logic in replaceValidator()

- ID: PVE-001
- Severity: High
- Likelihood: Medium
- Impact: High
- Target: RockXStaking
- Category: Business Logic [6]
- CWE subcategory: CWE-841 [4]

#### Description

The RockXStaking contract of RockX provides an external `replaceValidator()` function that allows for replacing a validator in case of misconfiguration. Our analysis with this function shows the current logic has a flawed implementation that needs to be corrected.

To elaborate, we show below its current implementation. By design, the current logic updates the internal accounting to remove the `oldpubkey`-mapped `pubkeyIndices` and add the new `pubkey`-mapped `pubkeyIndices`. However, there is a need to validate the `pubkeyIndices` of the `oldpubkey` is indeed present. The current implementation accidentally ensures the non-presence of `oldpubkey` (line 271)! Moreover, the current implementation can also be improved by validating the new `pubkey` does not exist, which is currently missing.

```
265     function replaceValidator(bytes calldata oldpubkey, bytes calldata pubkey, bytes
266         calldata signature) external onlyRole(REGISTRY_ROLE) {
267         require(pubkey.length == PUBKEY_LENGTH, "INCONSISTENT_PUBKEY_LEN");
268         require(signature.length == SIGNATURE_LENGTH, "INCONSISTENT_SIG_LEN");
269
270         // mark old pub key to false
271         bytes32 oldPubKeyHash = keccak256(oldpubkey);
272         require(pubkeyIndices[oldPubKeyHash] == 0, "PUBKEY_NOT_EXSITS");
273         uint256 index = pubkeyIndices[oldPubKeyHash] - 1;
274         delete pubkeyIndices[index];
275
276         // set new pubkey
277         bytes32 pubkeyHash = keccak256(pubkey);
```

```
277     validatorRegistry[index] = ValidatorCredential({pubkey:pubkey, signature:
           signature, stopped:false});
278     pubkeyIndices[pubkeyHash] = index+1;
279 }
```

Listing 3.1: `RockXStaking::replaceValidator()`

**Recommendation** Revise the validator replacement logic in the above routine to ensure the old one is correctly replaced by the new one.

**Status** This issue has been fixed in the following commit: `e0bd595`.

## 3.2 Suggested Adherence of Checks-Effects-Interactions

- ID: PVE-002
- Severity: Low
- Likelihood: Low
- Impact: Low
- Target: `RockXStaking`
- Category: Time and State [7]
- CWE subcategory: CWE-663 [3]

### Description

A common coding best practice in Solidity is the adherence of `checks-effects-interactions` principle. This principle is effective in mitigating a serious attack vector known as `re-entrancy`. Via this particular attack vector, a malicious contract can be reentering a vulnerable contract in a nested manner. Specifically, it first calls a function in the vulnerable contract, but before the first instance of the function call is finished, second call can be arranged to re-enter the vulnerable contract by invoking functions that should only be executed once. This attack was part of several most prominent hacks in Ethereum history, including the `DAO` [13] exploit, and the `Uniswap/Lendf.Me` hack [12].

We notice an occasion where the `checks-effects-interactions` principle is violated. Using the `RockXStaking` as an example, the `withdrawManagerFee()` function (see the code snippet below) is provided to externally interact with a given `to` address to transfer assets. However, the invocation of an external address requires extra care in avoiding the above `re-entrancy`.

Apparently, the interaction with the external address (line 357) starts before effecting the update on internal state (line 358 – 360), hence violating the principle. In this particular case, if the external address is an contract with certain hidden logic that may be capable of launching `re-entrancy` via the very same `withdrawManagerFee()` function. Note that there is no harm that may be caused to current protocol. However, it is still suggested to follow the known `checks-effects-interactions` best practice.

```
354     function withdrawManagerFee(uint256 amount, address to) external nonReentrant
        onlyRole(MANAGER_ROLE) {
355         require(amount <= accountedManagerRevenue, "WITHDRAW_EXCEEDED_MANAGER_REVENUE");
356         require(amount <= _currentEthersReceived(), "INSUFFICIENT_ETHERS");
357         payable(to).sendValue(amount);
358         accountedBalance -= int256(amount);
359         // track manager's revenue
360         accountedManagerRevenue -= amount;
361         emit ManagerFeeWithdrawn(amount, to);
362     }
```

Listing 3.2: RockXStaking::withdrawManagerFee()

In the meantime, we should mention that the supported tokens in the protocol do implement rather standard ERC20 interfaces and their related token contracts are not vulnerable or exploitable for re-entrancy.

**Recommendation** Apply necessary reentrancy prevention by following the checks-effects-interactions best practice.

**Status** This issue has been fixed in the following commit: e0bd595.

### 3.3 Proper minToMint Enforcement in mint()

- ID: PVE-003
- Severity: Low
- Likelihood: Low
- Impact: Low
- Target: RockXStaking
- Category: Business Logic [6]
- CWE subcategory: CWE-841 [4]

#### Description

The `RockXStaking` contract allows to deposit any number of ethers to the staking contract, and get back equivalent value of `uniETH` token (decided by real-time exchange ratio). In addition, the contract allows the user to specify the minimum `uniETH` amount (`minToMint`) that will be returned. While examining the related exchange logic, we observe the `minToMint` enforcement needs to be improved.

To elaborate, we show below the related `mint()` function. It comes to our attention that the `minToMint` parameter is enforced to ensure `require(toMint >= minToMint)` (line 669), which may not represent the actual amount of the returned `uniETH`. In other words, the current enforcement only ensures that the exchange ratio is at least 1. However, the proper enforcement requires the use of the final `toMint` of `totalXETH * msg.value / totalEthers` (line 672)!

```

658     function mint(uint256 minToMint, uint256 deadline) external payable nonReentrant
        whenNotPaused {
659         require(block.timestamp < deadline, "TRANSACTION_EXPIRED");
660         require(msg.value > 0, "MINT_ZERO");
661
662         // track balance
663         _balanceIncrease(msg.value);
664
665         // mint xETH while keeping the exchange ratio invariant
666         uint256 totalXETH = IERC20(xETHAddress).totalSupply();
667         uint256 totalEthers = currentReserve();
668         uint256 toMint = 1 * msg.value; // default exchange ratio 1:1
669         require(toMint >= minToMint, "EXCHANGE_RATIO_MISMATCH");
670
671         if (totalEthers > 0) { // avert division overflow
672             toMint = totalXETH * msg.value / totalEthers;
673         }
674         // mint xETH
675         IMintableContract(xETHAddress).mint(msg.sender, toMint);
676         totalPending += msg.value;
677
678         // spin up n nodes
679         uint256 numValidators = totalPending / DEPOSIT_SIZE;
680         for (uint256 i = 0; i < numValidators; i++) {
681             if (nextValidatorId < validatorRegistry.length) {
682                 _spinup();
683             }
684         }
685     }

```

Listing 3.3: RockXStaking::mint()

**Recommendation** Revisit the above logic to properly enforce the `minToMint`. An example revision is shown in the following:

```

658     function mint(uint256 minToMint, uint256 deadline) external payable nonReentrant
        whenNotPaused {
659         require(block.timestamp < deadline, "TRANSACTION_EXPIRED");
660         require(msg.value > 0, "MINT_ZERO");
661
662         // track balance
663         _balanceIncrease(msg.value);
664
665         // mint xETH while keeping the exchange ratio invariant
666         uint256 totalXETH = IERC20(xETHAddress).totalSupply();
667         uint256 totalEthers = currentReserve();
668         uint256 toMint = 1 * msg.value; // default exchange ratio 1:1
669
670         if (totalEthers > 0) { // avert division overflow
671             toMint = totalXETH * msg.value / totalEthers;
672         }
673     }

```

```

674     require(toMint >= minToMint, "EXCHANGE_RATIO_MISMATCH");
675
676     // mint xETH
677     IMintableContract(xETHAddress).mint(msg.sender, toMint);
678     totalPending += msg.value;
679
680     // spin up n nodes
681     uint256 numValidators = totalPending / DEPOSIT_SIZE;
682     for (uint256 i = 0; i < numValidators; i++) {
683         if (nextValidatorId < validatorRegistry.length) {
684             _spinup();
685         }
686     }
687 }

```

Listing 3.4: Revised `RockXStaking::mint()`

**Status** This issue has been fixed in the following commit: [e0bd595](#).

### 3.4 Proper Event Emission in `mint()`

- ID: PVE-004
- Severity: Informational
- Likelihood: N/A
- Impact: N/A
- Target: `RockXStaking`
- Category: Status Codes [8]
- CWE subcategory: CWE-391 [2]

#### Description

In Ethereum, the `event` is an indispensable part of a contract and is mainly used to record a variety of runtime dynamics. In particular, when an `event` is emitted, it stores the arguments passed in transaction logs and these logs are made accessible to external analytics and reporting tools. Events can be emitted in a number of scenarios. One particular case is when system-wide parameters or settings are being changed. Another case is when tokens are being minted, transferred, or burned.

In the following, we show the set of events defined in `RockXStaking`. While examining the list of events, we notice the following three of them are defined, but not used: `RevenueWithdrawedFromValidator` (line 876), `Redeemed` (line 881), and `RedeemFromValidator` (line 882). Therefore, there is a need to properly emit them when respective operations are performed. For example, we may emit `RedeemFromValidator` within `redeemFromValidators()` and emit `Redeemed` within `_payDebts()`.

```

873     event ValidatorActivated(uint256 node_id);
874     event ValidatorStopped(uint256 stoppedCount, uint256 stoppedBalance);
875     event RevenueAccounted(uint256 amount);
876     event RevenueWithdrawedFromValidator(uint256 amount);
877     event ValidatorSlashedStopped(uint256 stoppedCount, uint256 slashed);

```

```

878     event ManagerAccountSet(address account);
879     event ManagerFeeSet(uint256 milli);
880     event ManagerFeeWithdrawed(uint256 amount, address);
881     event Redeemed(uint256 amountXETH, uint256 amountETH);
882     event RedeemFromValidator(uint256 amountXETH, uint256 amountETH);
883     event WithdrawCredentialSet(bytes32 withdrawCredential);
884     event DebtQueued(address creditor, uint256 amountEther);
885     event XETHContractSet(address addr);
886     event DepositContractSet(address addr);
887     event RedeemContractSet(address addr);
888     event BalanceSynced(uint256 diff);

```

Listing 3.5: The Events Defined in `RockXStaking`

**Recommendation** Properly emit the above-mentioned events within respective functions to timely reflect state changes. This is very helpful for external analytics and reporting tools.

**Status** This issue has been fixed in the following commit: `e0bd595`.

### 3.5 Trust Issue of Admin Keys

- ID: PVE-005
- Severity: Medium
- Likelihood: Medium
- Impact: Medium
- Target: Multiple contracts
- Category: Security Features [5]
- CWE subcategory: CWE-287 [1]

#### Description

In `RockXStaking`, there is a privileged administrative account, i.e., the account with the `DEFAULT_ADMIN_ROLE` role. The administrative account plays a critical role in governing and regulating the staking-wide operations. It also has the privilege to control or govern the flow of assets within the protocol contracts. Our analysis shows that this privileged account needs to be scrutinized. In the following, we use the `RockXStaking` contract as an example and show the representative functions potentially affected by the privileges of the administrative account.

```

301     function setManagerFeeShare(uint256 milli) external onlyRole(DEFAULT_ADMIN_ROLE) {
302         require(milli >=0 && milli <=1000, "SHARE_OUT_OF_RANGE");
303         managerFeeShare = milli;
304
305         emit ManagerFeeSet(milli);
306     }
307
308     /**
309     * @dev set xETH token contract address
310     */

```

```
311     function setXETHContractAddress(address _xETHContract) external onlyRole(
312         DEFAULT_ADMIN_ROLE) {
313         xETHAddress = _xETHContract;
314         emit XETHContractSet(_xETHContract);
315     }
316
317     /**
318     * @dev set eth deposit contract address
319     */
320     function setETHDepositContract(address _ethDepositContract) external onlyRole(
321         DEFAULT_ADMIN_ROLE) {
322         ethDepositContract = _ethDepositContract;
323         emit DepositContractSet(_ethDepositContract);
324     }
325
326     /**
327     * @dev set redeem contract
328     */
329     function setRedeemContract(address _redeemContract) external onlyRole(
330         DEFAULT_ADMIN_ROLE) {
331         redeemContract = _redeemContract;
332         emit RedeemContractSet(_redeemContract);
333     }
334
335     /**
336     @dev set withdraw credential to receive revenue, usually this should be the
337         contract itself.
338     */
339     function setWithdrawCredential(bytes32 withdrawalCredentials_) external onlyRole(
340         DEFAULT_ADMIN_ROLE) {
341         withdrawalCredentials = withdrawalCredentials_;
342         emit WithdrawCredentialSet(withdrawalCredentials);
343     }
```

Listing 3.6: Example Privileged Operations in `RockXStaking`

We understand the need of the privileged functions for contract maintenance, but at the same time the extra power to the owner may also be a counter-party risk to the protocol users. It would be worrisome if the privileged administrative account is a plain EOA account. Note that a multi-sig account could greatly alleviate this concern, though it is still far from perfect. Specifically, a better approach is to eliminate the administration key concern by transferring the role to a community-governed DAO.

**Recommendation** Promptly transfer the privileged account to the intended DAO-like governance contract. All changes to privileged operations may need to be mediated with necessary timelocks. Eventually, activate the normal on-chain community-based governance life-cycle and ensure the in-

tended trustless nature and high-quality distributed governance.

**Status** This issue has been mitigated as the team confirms the use of Aragon DAO to use these administrative functions.



## 4 | Conclusion

In this audit, we have analyzed the design and implementation of the `staking` support in `RockX`, which makes it possible for anyone to access efficient and reliable mining and staking services. The staking contract allows users to deposit any number of ethers to the staking contract of `ETH 2.0`, and get back equivalent value of `uniETH` token (decided by real-time exchange ratio). The current code base is well structured and neatly organized. Those identified issues are promptly confirmed and addressed.

Meanwhile, we need to emphasize that `Solidity`-based smart contracts as a whole are still in an early, but exciting stage of development. To improve this report, we greatly appreciate any constructive feedbacks or suggestions, on our methodology, audit findings, or potential gaps in scope/coverage.



## References

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