

Cryptography and Its Applications

Content

- Classical Cryptography
- Symmetric key encryption
 - Stream Cipher
 - Block Cipher
- Message Authentication Code
- JWT
- Public Key Cryptography
- Public Key Infrastructure (PKI)

Goals of Cryptography

- The most fundamental problem cryptography addresses: **ensure security of communication over insecure medium**
- What does secure communication mean?
 - confidentiality (privacy, secrecy)
 - only the intended recipient can see the communication
 - integrity (authenticity)
 - the communication is generated by the alleged sender
- What does insecure medium mean?
 - Two possibilities:
 - Passive attacker: the adversary can eavesdrop
 - Active attacker: the adversary has full control over the communication channel

Approaches to Secure Communication

- Steganography
 - “covered writing”
 - hides the existence of a message
 - depends on secrecy of method
- Cryptography
 - “hidden writing”
 - hide the meaning of a message
 - depends on secrecy of a short key, not method

Basic Terminology

- **Plaintext**: origin message
- **Ciphertext** : transformed message
- **Key**: secret used in transformation
- **Encryption/Decryption**: forward/reverse transformation
- **Cipher**: Algorithm of transformation
- **Cryptanalysis**: science and practice of breaking cryptographic systems/schemes

Shift Cipher

- The Key Space:
 - $[0 \dots 25]$
- Encryption given a key K :
 - each letter in the plaintext P is replaced with the K 'th letter following corresponding number (shift right)
- Decryption given K :
 - shift left

History: $K = 3$, Caesar's cipher



Shift Cipher: Cryptanalysis

- Can an attacker find K?
 - YES: by a brute force attack through exhaustive key search,
 - key space is small (≤ 26 possible keys).
- Lessons:
 - Cipher key space needs to be large enough.
 - Exhaustive key search can be effective.

Mono-alphabetic Substitution Cipher

- The key space: all permutations of $\Sigma = \{A, B, C, \dots, Z\}$
- Encryption given a key π :
 - each letter X in the plaintext P is replaced with $\pi(X)$
- Decryption given a key π :
 - each letter Y in the ciphertext P is replaced with $\pi^{-1}(Y)$

Example:

$\Sigma =$ A B C D E F G H I J K L M N O P Q R S T U V W X Y Z

$\pi =$ B A D C Z H W Y G O Q X S V T R N M L K J I P F E U

BECAUSE $\Rightarrow$ AZDBJSZ

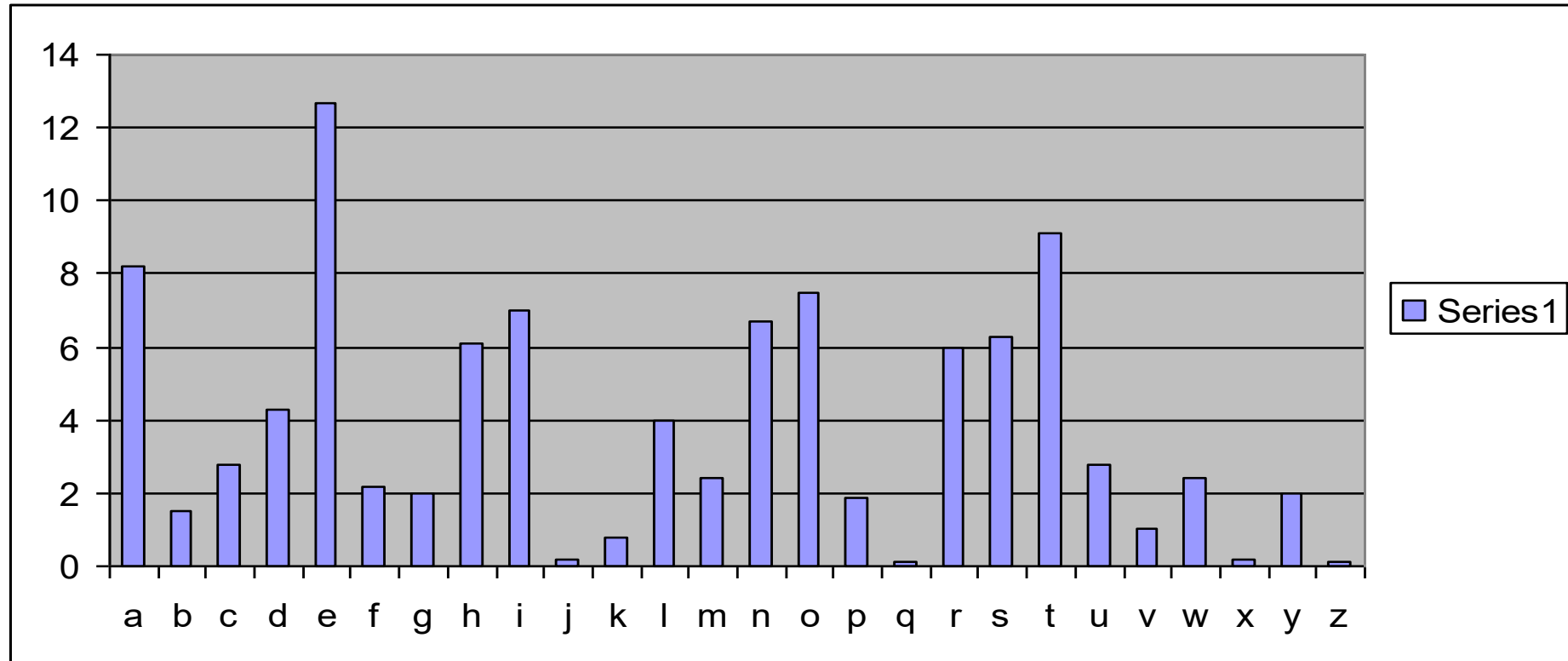
Strength of the Mono-alphabetic Substitution Cipher

- Exhaustive search is difficult
 - key space size is $26! \sim 4 \times 10^{26} \sim 2^{88}$
- Dominates the art of secret writing throughout the first millennium A.D.
- Thought to be unbreakable by many back then
- How to break it?

Cryptanalysis of Substitution Ciphers: Frequency Analysis

- Basic ideas:
 - Each language has certain features: frequency of letters, or of groups of two or more letters.
 - Substitution ciphers preserve the language features.
 - Substitution ciphers are vulnerable to frequency analysis attacks.

Frequency of Letters in English



Adversarial Models for Ciphers

- The language of the plaintext and the nature of the cipher are assumed to be known to the adversary.
- **Ciphertext-only attack:** The adversary knows only a number of ciphertexts.
- **Known-plaintext attack:** The adversary knows some pairs of ciphertext and corresponding plaintext.
- **Chosen-plaintext attack:** The adversary can choose a number of messages and obtain the ciphertexts
- **Chosen-ciphertext attack:** The adversary can choose a number of ciphertexts and obtain the plaintexts.

Security Principles

- **Kerckhoffs's Principle:**
 - A cryptosystem should be secure even if everything about the system, except the key, is public knowledge.
- **Shannon's maxim:**
 - "The enemy knows the system."
- Security by obscurity doesn't work
- Should assume that the adversary knows the algorithm; the only secret the adversary is assumed to not know is the key
- What is the difference between the algorithm and the key?

Notation for Symmetric-key Encryption

- A symmetric-key encryption scheme is comprised of three algorithms
 - **Gen** the key generation algorithm
 - The algorithm must be probabilistic/randomized
 - Output: a key k
 - **Enc** the encryption algorithm
 - Input: key k , plaintext m
 - Output: ciphertext $c := \mathbf{Enc}_k(m)$
 - **Dec** the decryption algorithm
 - Input: key k , ciphertext c
 - Output: plaintext $m := \mathbf{Dec}_k(c)$

Requirement: $\forall k \forall m [\mathbf{Dec}_k(\mathbf{Enc}_k(m)) = m]$

Stream Ciphers

- Plaintext P
- Ciphertext C
- Keystream Ks:
 - $\text{Len}(Ks) = \text{Len}(P)$
 - Each P has its own Ks
- Encryption & Decryption:
 - $C = P \text{ XOR } Ks$ (XOR bit by bit)
 - $P = C \text{ XOR } Ks$ (XOR bit by bit)
- Why does it work?:
 - Sender and receiver use the same way to generate K for each message.

Stream Ciphers characteristics

- Encrypt/decrypt bit by bit
 - Prevent frequency analysis
- Very fast
- Keystream generation algorithm based on Pseudo Random Number Generation
- Become vulnerable if a keystream is used twice.

Stream Cipher applications and standards

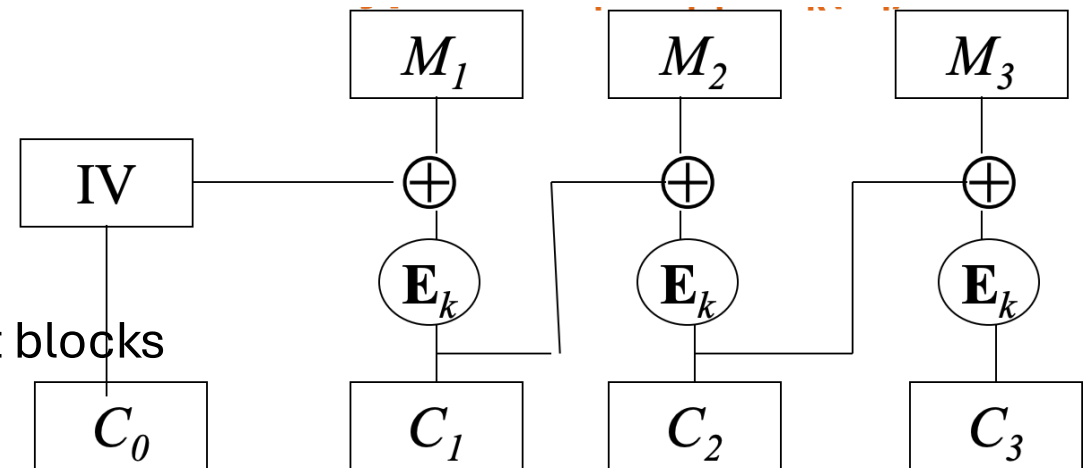
- Online content broadcasting
 - Content scramble
 - DRM (Digital Right Management)
- Typical standards
 - RC4
 - DVD – CCA (DVD Copy Control Association)

Block Ciphers

- Encrypt block-by-block
 - n-bit Plaintext --> n-bit ciphertext
- Common notations:
 - $P : \{0,1\}^n$
 - $C : \{0,1\}^n$
 - $K : \{0,1\}^s$
 - n: block size
 - s: key size
- Prevent frequency analysis

Block Cipher Encryption Modes:

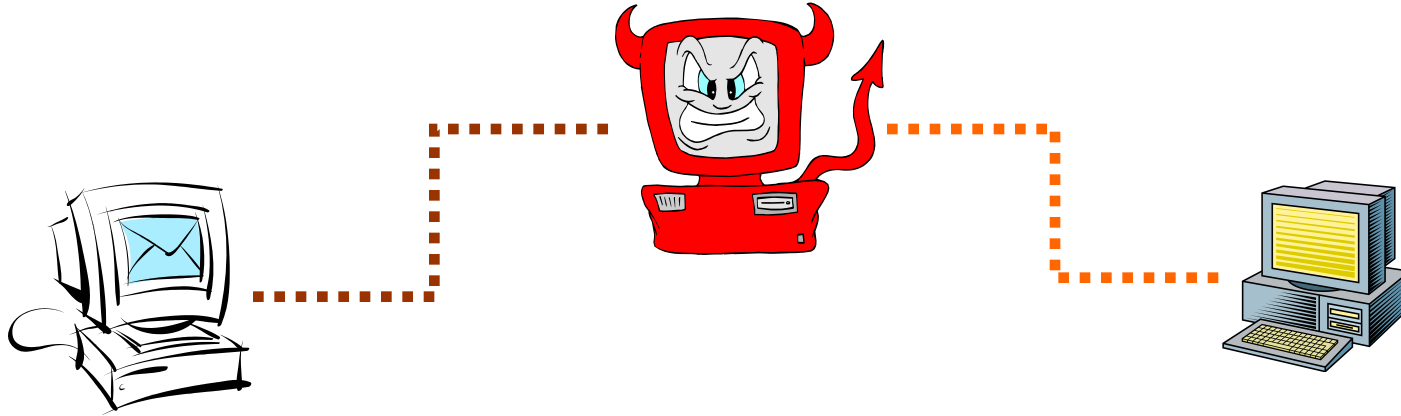
- ECB : Electronic Code Book
 - Each block is encrypted/decrypted separately
 - $C_i = E(k, X_i)$
 - $X_i = D(k, C_i)$
 - Not recommended for encrypting more than one block
- CBC: Cipher Block Chaining
 - Random Initial Vector (IV)
 - Next input depends on previous output
 - Randomized:
 - Same input block can have different output blocks



Block Cipher Standards and applications

- DES: Data Encryption Standard
 - Obsolete: 1977 – 2001
 - Block size: 64 bits
 - Key size: 56 bits
 - Vulnerable for brute-force attack (short key size)
 - Designed for hardware
 - Obsolete
- AES: Advanced Encryption Standard
 - Replace DES
 - Block size: 128 bits
 - Key size: 128, 192, 256 bits
 - Applications: Data/Disk encryption, WiFi, VPN, HTTPs

Data Integrity and Source Authentication



- Encryption does not protect data from modification by another party.
 - **Why?**
- Need a way to ensure that data arrives at destination in its original form as sent by the sender and it is coming from an authenticated source.

Hash Functions

- A hash function maps a message of an arbitrary length to a m-bit output
 - output known as the **fingerprint** or the **message digest**
- What is an example of hash functions?
 - Give a hash function that maps Strings to integers in $[0, 2^{\{32\}}-1]$
- Cryptographic hash functions are hash functions with additional security requirements

Common Hash Function

- MD5
 - 128-bit output
- SHA1
 - 160-bit output
- SHA2 (SHA-224, SHA-256, SHA-384, SHA-512)
 - 224-,256-,384-, 512-bit outputs

Usages of Cryptographic Hash Functions

- Software integrity
- File integrity
- Message integrity:
 - MAC (HMAC): Message Authentication Code

Message Authentication Code

- Original Message M
- MAC is a code calculated from M (using a secret K)
 - $\text{MAC}(K, M) = H_K(M)$
- Data to send : $D = M \parallel \text{MAC}(K, M)$
- Upon received:
 - Verify received $\text{MAC}(K, M)$ against M and K
- $\Rightarrow$ Message integrity verification
- MAC can be calculated using hash function
 - $\text{MAC}(K, M) = H(K \parallel M)$

HMAC

- $\text{HMAC}(K, M) = \text{H}(K \parallel \text{H}(K \parallel M))$
 - Two hashes: Inner hash and outer hash
- Why single hash is not enough?
 - With simple MAC, the original message can be added and a new digest can be recalculated without knowing the secret key.
- HMAC Application :
 - Message Authentication
 - JWT

JSON Web Token

- JWT = secure, compact, open standard for information exchange between client and server
- Applications:
 - Authentication: verify "who you are"
 - and Authorization: verify "what you are allowed to do"
- Structure:
 - HEADER . PAYLOAD . SIGNATURE

JSON Web Token (cont.)

- **HEADER:**

- Base64 encoded JSON object
- Contain metadata including "algorithm" and "type"
- Example:
 - `{"alg":"HS256","typ":"JWT"}` <-- base64 --> `eyJhbGciOiJIUzI1NiIsInR5cCI6IkpXVCJ9`

- **PAYLOAD:**

- Base64 encoded JSON object
- Contain "claims" (application data)
- Registered claims (optional): iat, sub, iss, aud, exp
- User claims: what ever you want

JSON Web Token (cont.)

- SIGNATURE:
 - HMACSHA256(HEADER . PAYLOAD, SECRET)
 - Can use other signing algorithm (RSA256)
 - It is used for verification
- DEMO:
 - base64 command
 - xxd command
 - openssl command to generate HMAC

Symmetric key cryptography review

- Sender and receiver must share a single secret key
- The key use for:
 - Encryption/Decryption : AES, Stream Cipher
 - Message Authentication: HMAC
- Because of the share keys:
 - Secure channels for key distribution
 - Many keys required for multiple parties ($O(n^2)$ keys for n parties)
 - Two parties should have prior relationship (key exchange) --> not convenient

Solution: Public Key Cryptography

- Assume multiple party communications:
- One party (party A) has:
 - Public key (Pub_A): known by others
 - Private key (Priv_A): kept in secret (only A knows the key)
- Encryption application: B --- send M ---> A
 - At B: $C = \text{encrypt}(\text{Pub}_A, M)$
 - At A: $M = \text{decrypt}(\text{Priv}_A, C)$
- Digital Signature application: A --- send (M || signature) ---> B
 - At A: $\text{signature} = \text{sign}(\text{Priv}_A, M)$
 - At B: $\text{verify}(\text{signature}, \text{Pub}_A)$

Real world application of Public Key Cryptography

- Public key cryptographic algorithms are slow --> Do not apply it on the data
- Combine Public Key encryption and symmetric keys
 - Before communication session: senders send symmetric keys (encryption keys) to receivers using Public Key encryption
 - Use the encryption keys for secured data exchange
- Practical digital signature:
 - Messages are hashed --> hash value
 - Hash value is signed with signing key (public key)

HMAC vs Digital Signature

- HMAC guarantees integrity and authenticity
- Digital Signature guarantees integrity, authenticity and non-repudiation
 - The signer cannot deny that he/she didn't do this.
- Digital Signature is slower than HMAC

Public Key Application : SSH Authentication

- At SSH server :
 - `~/.ssh/authorized_keys`: contains public keys of authorized clients/hosts
- SSH client :
 - Use its private for data encryption to communicate with the server

Public Key Infrastructure

- A framework for applying Public Key Cryptography
- It helps exchanging public keys in a trusted way.
 - Publishing "certificates" to a centralized repository
- Digital Certificate: A document containing:
 - Subject identifier
 - Public Key
 - Digital Signature to protect the document.

Certificate related operations

- Issuing a certificate:
 - Done by a CA
 - To bind a public key with an identity
- Storing a certificate:
 - At your own devices (just like persons keeping their ID cards)
- Sharing and verifying certificate:
 - Others can download certificates and verify it with VA to make sure its validity (Trust Model)

OpenSSL

- Implement core cryptographic components of PKI
 - Keypair generation
 - Certificate Signing Requests (CSR)
 - Certificate issuance
 - ...

Certificate Issuance procedure

- [Client] Client generate key pair (private-public keys)
 - OpenSSL , Certbot
- [Client] Client generate CSR (Code Signing Request), a text file
 - CSR contains: client information and client's public key
 - OpenSSL , Certbot
- [Client] client uploads CSR to [CA]
 - Manual via a web form / Automatic via ACME client
- [CA] CA validates CSR and issues an SSL Certificate
- [CA] CA sends SSL certificate to client

ACME - SSL certificate issuance protocol

- Certificates issuance procedure can be automatic via ACME protocol
 - Automated Certificate Management Environment
- Most CAs provide ACME interface
 - Let's Encrypt
 - Google Trust Services
 - ZeroSSL
 - GlobalSign
 - Sectigo

Self-signed SSL certificate

- Generate keypair
 - Expected input: private key file location. Sent via command line switch
 - Output:
 - private key file created
 - Private key file contains : both privatekey and publickey
 - E.g: openssl genrsa -out private.key 2048
- Generate CSR:
 - Expected Inputs: client information (CN: your domain name or ip address)
 - E.g: openssl req -new -key private.key -out certificate.csr
- Upload certificate.csr to CA : skipped
- Generate certificate:
 - E.g.: openssl x509 -req -days 365 -in certificate.csr -signkey private.key -out certificate.crt -notes

Using the self-signed certificate

- [Server] Using the certificate to config nginx
 - you need: crt file and private.key file
- [Client] Import the certificate to client's certificate store
 - Make the client trust the certificate without CA assistance.
- Import certificates at client