# GeoServ: A Distributed Urban Sensing Platform

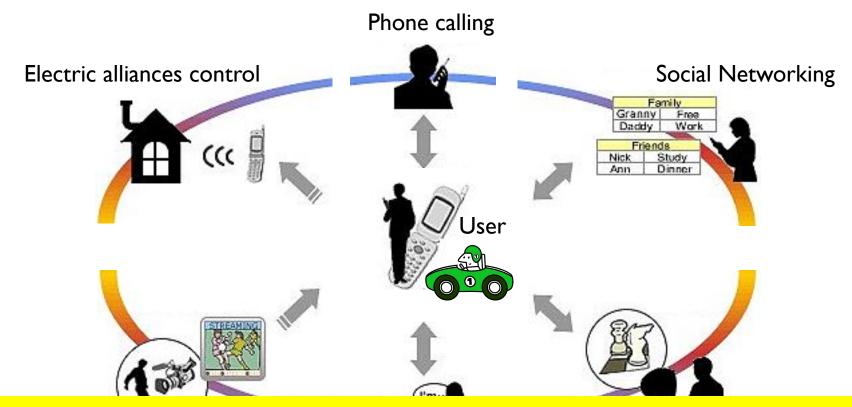
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# Mobile Computing Trends in US

- Desktop, laptop, netbook:
  - Now, I billion users
  - In 2015, 2 billion users
- Mobile phones
  - Now, 4.3 billion users
  - In 2013, 5.8 billion users
- Smartphones
  - Now, 150 million users
  - In 2020, billions of users???
  - Urban Sensing Enablers



# Urban Sensing with Smartphones

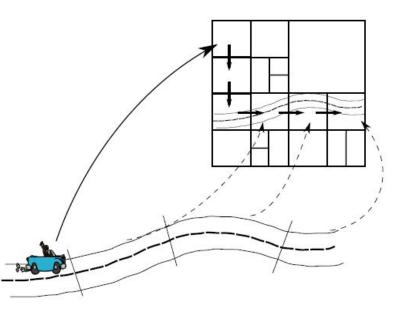


- I0 million mobile users generate sensor data updates and queries with the rate of IKB/sec per user
  - Bandwidth required : >80Gbps, Storage space required: > 36TB/hr

## Design Issues

## A mobile-to-mobile overlay network

- PeerTIS: message routing via DHT [1]
- Issues:
  - Network address translation (NAT)
  - Additional services such as session initiation protocol (SIP) or P2P proxy servers.
  - Intolerable delays
  - Significant resource consumption
    - e.g., battery, processing power, and bandwidth



 Therefore, consider using the Internet servers for large-scale urban sensing!

Design Issues (cont.)

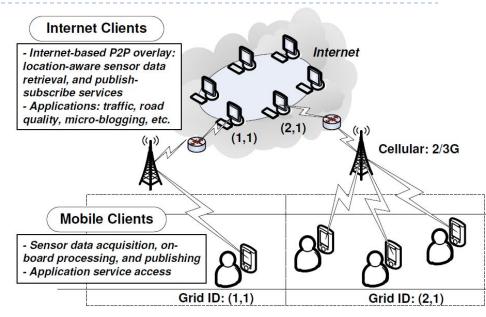
### Existing Approach (in using Internet servers)

- Centralized multi-tier architecture
  - Provide scalable services by provisioning more servers and bandwidth at the datacenters.
- Semi-hierarchical architecture
  - Each organization maintains database servers for its own "stationary sensors."
  - Appropriate for participatory sensing by sharing users' computing resources.
- Still, need a scalable architecture to support billions of mobile users!
- Consider a two-tier sensor networking architecture for large-scale participatory sensing.

# GeoServ: Two-Tier Architecture

### Two-tier Architecture

- I<sup>st</sup> tier: Internet-based fixed servers : form a distributed P2P sensor networking overlay
- 2<sup>nd</sup> tier: mobile users can publish/access sensor data through Internet servers.



The design emphasis for urban sensing applications is scalability and location-preservation for sensor data.

# Two Services using GeoServ

- 1. Location-aware sensor data retrieval service
- 2. Location-aware publish-subscribe service





#### Applications

- Street-level traffic flow information service
- Vehicular safety warning service
- Ride quality monitoring service
- GeoServ@CCGRID 2011

System Design Requirements

- 1. Seamlessly access to the sensor storage based on geographic locations.
- 2. Load balancing due to skewed distribution of vehicles on the road.
- 3. Publish-subscribe services
- Dynamic mobile user association to the Internet server based on one's current location.

# GeoServ Approach

## GeoTable

- □ Key Based Routing (KBR): key space is geographic location.
- Apps running on Internet servers and mobile clients communicate one another via GeoTable.

## GeoPS

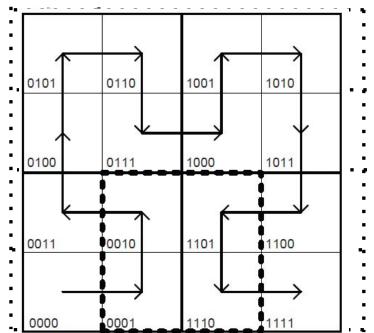
An efficient publish-subscribe routing method based on the geographic routing via GeoTable.

## Other Components

- GeoServDB
  - Manages sensor data of the grid space.
- GeoServMobile
  - Provides control and transparent access of sensing resources.

## LOCATION-AWARE SENSOR DATA RETRIEVAL SERVICE

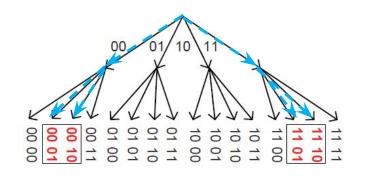
- In GeoServ, we divide the geographic area of interest into fixed size grids.
  2<sup>M</sup>x2<sup>M</sup>grids
  - *M* is the smallest exponent that covers the entire area.
  - E.g., size of the contiguous U.S. is approximated as 3000km X 3000km
    - 2<sup>13</sup> X 2<sup>13</sup> fixed grids where R is given as I km.

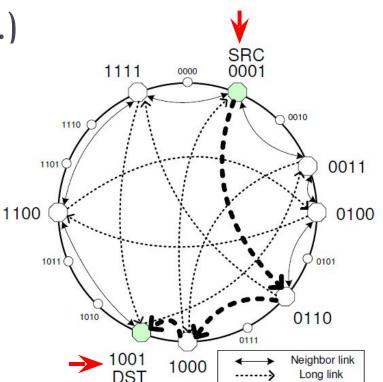


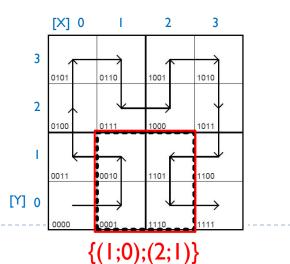
- How to construct overlay? Hilbert space filling curve!
  - A linear mapping function where successive points are nearest neighbors in the 2D grid.
  - The construction of the Hilbert curve is recursively defined.

## LOCATION-AWARE SENSOR DATA RETRIEVAL SERVICE (cont.)

- Routing Semantics
  - Geocasting to a single grid point
  - Geocasting to multiple grid points
    - Simple rectangular area based addressing as {(x1;y1); (x2;y2)}.
  - Concurrent Geocasting (for GeoPS)
    - Due to overhead of sequential search, we utilize how the Hilbert curve is constructed.







## LOCATION-AWARE PUBLISH-SUBSCRIBE SERVICE

### GeoPS

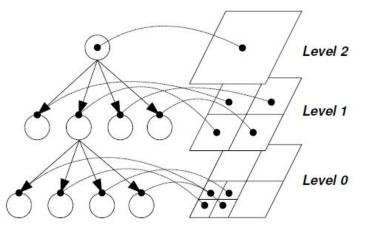
A publish-subscribe service where the data updates on a region are published subscribers.

### A key design issue

To build a multicast tree a multicast tree that exploits the geographic locality of the group members.

## The idea of hierarchical geographic location

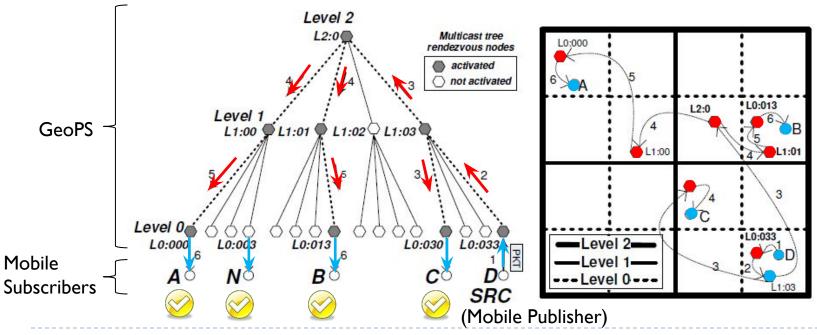
Inspired by hierarchical geographic location services (HGLS) in mobile ad hoc networks such as GLS [2] and HIGH-GRADE [3].



## LOCATION-AWARE PUBLISH-SUBSCRIBE SERVICE (cont.)

## A multicast tree construction in GeoPS

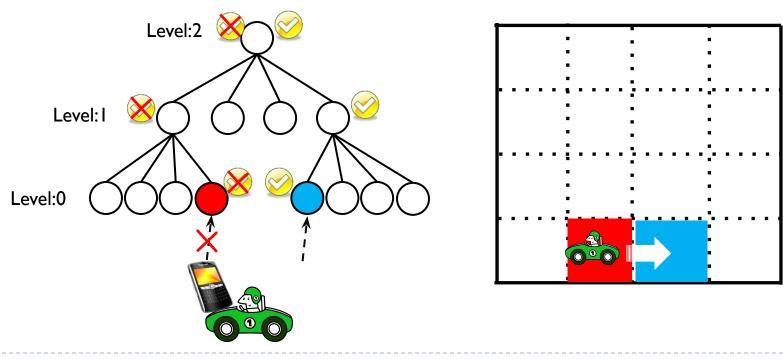
- Build a multicast tree over this geographic hierarchy
- Use our geocasting algorithm over the tree to preserve geographic locality.



GeoServ@CCGRID 2011

## LOCATION-AWARE PUBLISH-SUBSCRIBE SERVICE (cont.)

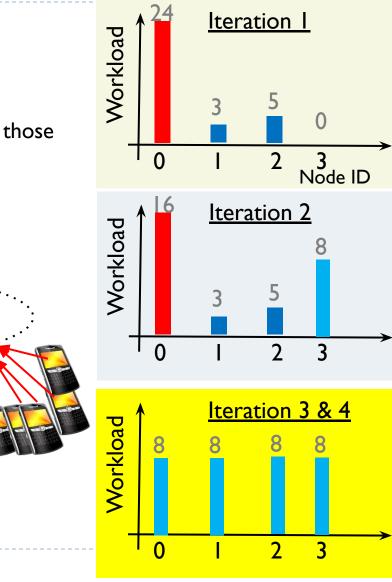
- Mobility handling in GeoPS
  - A mobile client's subscription needs to be updated (to upper layers) whenever the client crosses the level boundary (via explicit leave and join).
  - When mobile client C moves to the adjacent grid on the left (crossing level I boundary), rendezvous points at level 0, 1, 2 are updated.

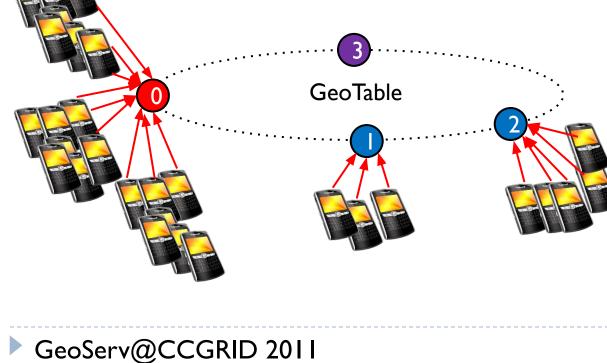


# LOCATION-AWARE PUBLISH-SUBSCRIBE SERVICE (cont.)

#### Load balancing in GeoPS

- Dynamic load balancing
  - In highly populated regions with mobile clients,
  - GeoTable re-organizes overlay nodes such that those regions are served by more number of nodes.





# **Evaluation List**

- 1. Location-aware data retrieval
- 2. Impact of query region sizes and concurrent geocasting
- 3. Load balancing
- 4. Location-aware publish-subscribe service

## **Evaluation Setup**

### 1. An event-driven discrete-time simulator in C#

- Supports dynamic node generation/join/leave, load balancing, and publish/subscribe features.
- Long links are set to 5, as recommended in Symphony DHT [4]
- 2. A large-scale participatory vehicular sensing scenario
  - I00 different simulations
- 3. Realistic mobility generation:VanetMobiSim



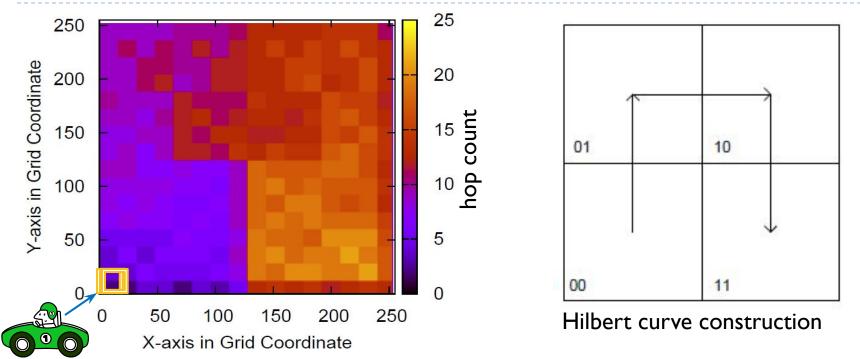


Westwood (Tigermap)

# Simulation Setup: (1) Location-aware data retrieval

- Goal: To show that GeoTable preserves geographic locality
- □ Set-up:
  - Grid size: 256x256
  - Place a querying node at a grid point (0; 0)
  - Measure the hop count of a remote query with the square area of size 4x4,
  - □ Vary the location of the query from (0,0) to (252,252).
  - The number of overlay nodes is set to 1000.

# Simulation Result: (1) Location-aware data retrieval



As distance from (0; 0) increases, the hop counts increases (getting brighter).

- Non-uniform colors is that some degree of locality is lost after linearization, and long links are randomly assigned.
- Due to the recursive construction property of the Hilbert curve, the average hop count increases as we move clockwise.
- Locality is preserved at the higher level.

Simulation Setup: (2) Impact of query region sizes and concurrent geocasting

- Goal: to show the sensitivity of routing cost with different query region sizes: 2x2 and 6x6.
- □ Set-up:
  - $\Box$  The querying node is located at grid (0, 0).
  - And it sends queries with different area sizes.
  - Vary the distance between the querying node and the query region from 0 to 255.
  - Compare the performance of various schemes:.
    - Conventional DHT with consistent hashing (CH)
      - Each grid is randomly mapped into the key space;
      - 2x2 and 6x6 queries require 4 and 36 unicasts, respectively.
    - Sequential geocasting (SG)
    - Concurrent geocasting (CG)

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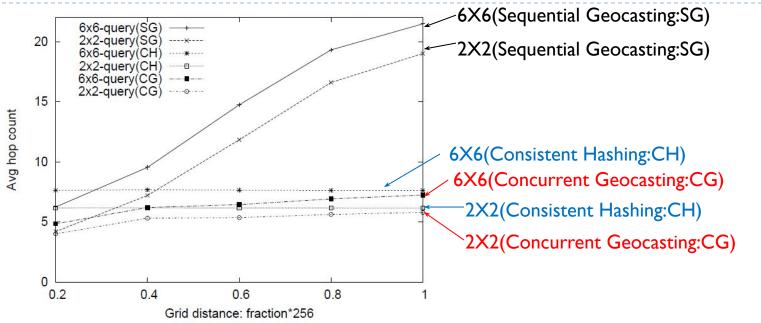
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distance

# Simulation Result: (2) Impact of query region sizes and concurrent geocasting



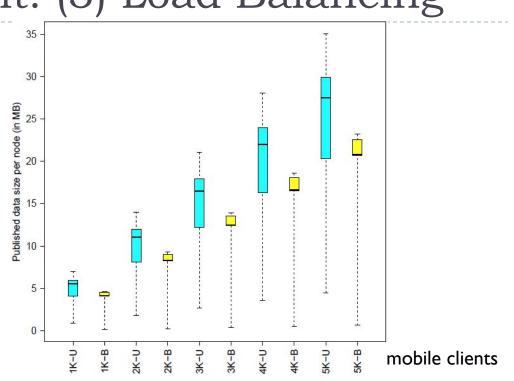
- As the grid distance increases, the average hop count of SG and CG increases (CG is better than SG).
  - But, no change on CH with the distance due to lack of locality.
- Note that delay improvement of CG against SG comes at the cost of more packet forwarding.
  - Still, far more efficient than issuing individual unicast to each segment.

# Simulation Setup: (3) Load balancing

- Goal: to show how heterogeneous distribution of mobile clients influences the overall load imbalance
- □ Set-up
  - Mobile clients publish sensor data to the overlay nodes.
  - Road topology information from Tiger Los Angeles map.
    - □ The area size is 12,800m X 12,800m, centered at the UCLA campus.
    - The grid size of 50m X 50m, composed of 256x256 grids.
  - Mobile Clients
    - The different numbers of mobile clients from 1000 to 5000 with a gap of 1000 nodes.
    - Data generation with the rate of 128 Bytes/sec.
  - Overlay nodes
    - The number of overlay nodes is fixed to 1000 nodes.
  - Measurement: the total published data size per node to see how per-node workload is unbalanced.

# Simulation Result: (3) Load Balancing

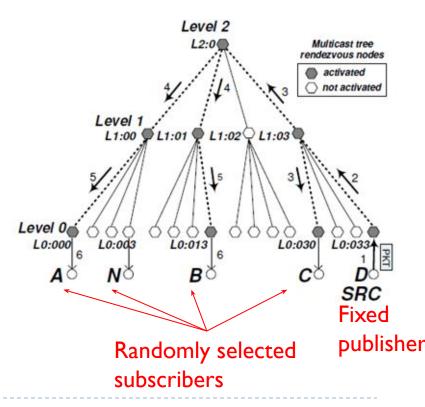
- boxplot shows min, 25%, median, 75%, and max;
- U: unbalanced, B: balanced
- □ IK: 1000



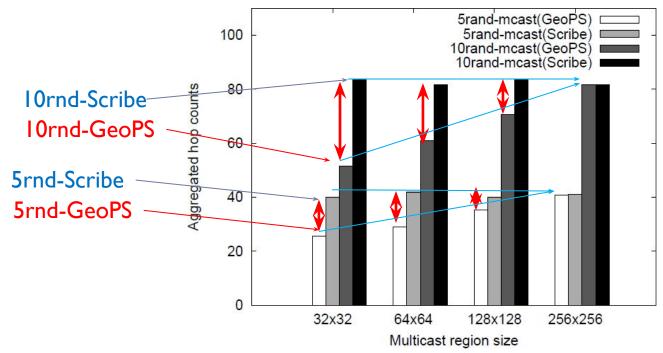
- As expected, the total data size increases linearly, as the number of mobile clients increases.
- The case without load balancing shows much higher variation as opposed to the case with load balancing.
- Still, minor variations in the balanced.
  - Because load balancing requires several iterations of leave/join-based load balancing operations.

# Simulation Setup: (4) Location-aware publish-subscribe service

- Goal: to show the geographic locality of our subscription-based multicast routing.
- □ Set-up:
  - The region size ranges from 32x32 to 256x256 grids.
  - Vary the origin of the region from (0, 0) to the maximum allowable.
    - E.g., for 32x32, it is (224, 224), and report the average hop count.
  - Randomly choose 5 or 10 random grids
  - Measure the aggregated number of hop counts.
  - Compare the performance of GeoPS with Scribe.
    - Note that Scribe destroys the locality by using consistent hashing.



# Simulation Result: (4) Location-aware publish-subscribe service



- GeoPS exploits the locality of receivers.
- As the area size (where the subscribers lie) increases, geographic locality among subscribers disappears,
- Accordingly, the cost of GeoPS increases, converging to that of Scribe in the case of 256x256.

# **Concluding Remarks**

## A scalable sensor networking system design

To enable location-relevant sensor data sharing among mobile users with smartphones.

## A two-tier sensor networking system: GeoServ

- Exploits the Internet infrastructure.
- Mobile users publish/access sensor information through the Internet-based distributed sensor storage.

## Mathematical Proofs in the paper

We proved that GeoServ protocols preserve geographic locality and validated their performance via extensive simulations.

## References

[1] J. Rybicki, B. Scheuermann, M. Koegel, and M. Mauve. PeerTIS: a Peer-to-Peer Traffic Information System. In VANET, 2009.

[2] J. Li, J. Jannotti, D. S. J. D. Couto, D. R. Karger, and R. Morris. A Scalable Location Service for Geographic Ad Hoc Routing. In MobiCom, 2000.

[3] Y. Yu, G.-H. Lu, and Z.-L. Zhang. Enhancing Location Service Scalability with HIGH-GRADE. In MASS, 2004.

[4] G. S. Manku, M. Bawa, and P. Raghavan. Symphony: Distributed Hashing in a Small World. In *USITS*, 2003.

[5] USA to Add 80 Million New Smartphone Users by 2011. http://twittown.com/mobile/mobile-blog/ usa-add-80-million-new-smartphone-users-2011.

