

A Near-Optimal Source Location Privacy Scheme for Wireless Sensor Networks



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Outline

- 1. Introduction
- 2. Privacy and Attacker Models
- 3. Integer Linear Programming Model and Results
- 4. Near-optimal Heuristic
- 5. Conclusions

What is a Wireless Sensor Network?

A wireless sensor network (WSN) is a collection of computing devices called nodes, they have:

- ▶ a short range wireless radio
- ▶ an array of sensors such as light, heat and humidity
- ► a simple low powered CPU
- ► a battery with **limited** power supply

Applications include:

- Tracking
- Monitoring (Environment, Assets, ...)



What is Context Privacy?

- Privacy threats can be classified as either content-based or context-based
- Content-based threats have been widely addressed (using cryptography)
- Context-based threats are varied
 - Location of event source
 - Location of base station
 - Time at which the event occurred
- ▶ We focus on protecting the **location** context of the **event source**



The Problem of Source Location Privacy (SLP)

Given:

- A WSN that detects valuable assets
- A node broadcasting information about an asset

Found:

- An attacker can find the source node by backtracking the messages sent through the network.
- So by deploying a network to monitor a valuable asset, a way has been provided for it to be captured.

The Problem:

- Panda-Hunter Game
- Difficult



Privacy Model

Aim of an SLP protocol: Prevent the attacker from capturing an asset through information the WSN leaks.

- ► A stationary asset cannot be protected as an attacker can perform an exhaustive search.
- A mobile asset will only stay in detection range of a WSN node for a certain amount of time.
- ▶ The SLP problem can only be considered when it is time-bounded.
- ▶ The *safety period* is how long the asset will be protected for.

Aim of an Attacker: to reach the source within the safety period.

The attacker:

- ▶ is present in the network
- ► is mobile
- ► has a limited range
- starts at the sink
- follows the first new packet it receives

How to Develop a Technique?

Previous approaches: Have idea \rightarrow implement it \rightarrow test performance

- Developed solutions have had good properties proven about them
- Optimal solutions have been found for global attackers

Better approach:

Find optimal solution \rightarrow create heuristic with similar output \rightarrow test performance

- ▶ Performance needs to be tested as the optimal solution may not be implementable
- Optimal solution doesn't give optimal algorithm, so some creativity is still needed to develop a heuristic
- ► Can use different optimality measures to look for different solutions

Integer Linear Programming Model I

- Integer Linear Programming (ILP) allows an optimal solution to be found to certain problems
- ▶ Express SLP-aware routing as an ILP optimisation problem
- ► A solution is obtained using IBM's ILOG CPLEX

Aim to obtain the best utility for this objective function:

maximise	The distance from the attacker's final position to the source	
subject to	Routing Constraints ctR1 to ctR6,	(1)
	Attacker Constraints ctA1 to ctA7.	

Integer Linear Programming Model II

- ctR1 At t = 0, no messages are broadcasted.
- ctR2 From t > 0, each source node generates a message every P_{src} until the safety period is reached.
- ctR3 A node sends one message or no messages in a time slot.
- ctR4 Once a message is broadcasted by a node it is not broadcasted by that node again.
- ctR5 A node can broadcast a message only after a neighbour broadcasted that message in a previous time slot.
- ctR6 All messages sent by the source(s) must reach the sink.

Integer Linear Programming Model III

ctA1 At t = 0 the attacker moves from the attacker's starting position to that same position.

- ctA2 The attacker makes exactly one move each time slot.
- ctA3 A move must be from the attacker's current location.
- ctA4 If the attacker moves to *n* from *m* at time τ , then it must be because at time τ the node *n* broadcasted a message.
- ctA5 If the attacker receives a message that it has not previously moved in response to, then the attacker moves in response to that message.
- ctA6 If the attacker moved in response to a message at time τ , then at no time $\tau' > \tau$ will the attacker move in response to that message again.
- ctA7 If no neighbour broadcasts a message the attacker stays where it is.

Model Result I

- ► Source at node 1 (top left)
- Sink at node 13 (centre)
- Attacker starts at the sink (centre)
- ▶ 9 slots per second
- ► 7 Messages sent
- ► Source Period 1 second/msg
- Safety Period of 7 seconds

Model Result II

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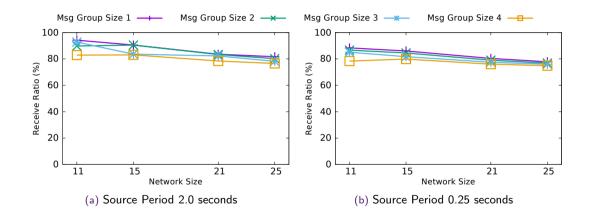
Inspired Heuristic

- ▶ Implementing the optimal solution requires global knowledge, so is unsuitable for a WSN.
- ► A heuristic was implemented instead based on these observations:
- 1. The routing path should go around the sink and approach from the opposite direction to the source.
- 2. Some routes should take the shortest path from the source to the sink.
- 3. Messages should be delayed so multiple messages are grouped together.
- 4. Messages should be delayed as late as possible with respect to the safety period.

Routing Algorithm

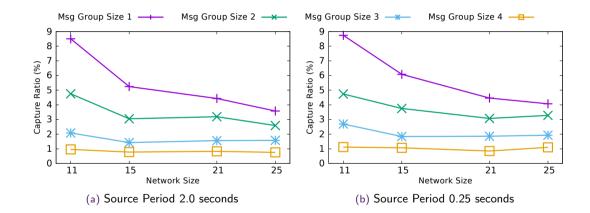
- 1. AvoidSink: In this first stage messages are routed around the sink.
- 2. **Backtrack**: Messages may end up attempting to go towards the sink and not having any valid routes, so they need to backtrack.
- 3. **ToSink**: Once a message has finished routing to avoid the sink it needs to be delivered to the sink.
- 4. **FromSink**: Finally, the message is sent in a starburst from the sink.

Results – Receive Ratio



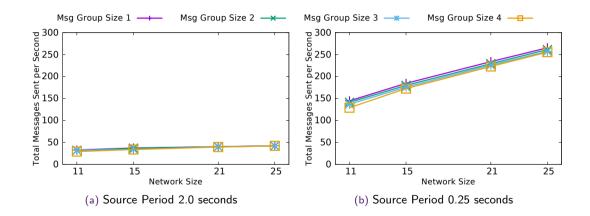
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Results – Capture Ratio



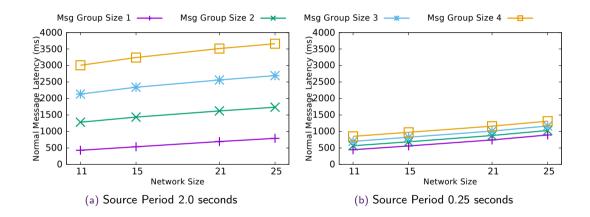
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Results – Messages Sent per Second



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Results – Message Latency



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Discussion

► A different objective function could be optimised for, such as:

- Latency
- Messages Sent
- ▶ High latency may make this technique unsuitable for some applications
 - for military it may be too high
 - for animal monitoring it should be acceptable
- Different contraints could be used
 - We require delivery of all messages, some applications may not require this
- ▶ The ILP model could be extended to support other SLP techniques (e.g., fake sources)

Summary

- Presented a way to model SLP-aware routing as an ILP optimisation problem
- Obtained an optimal solution from that model
- Implemented a near-optimal heuristic based on the optimal model result
- \blacktriangleright 1% capture ratio can be achieved by trading off for a higher latency

Thank You for Listening

Any Questions?