Convex Partitioning, Jaehne, Kahl, Kerner, Murrey Institut für Informatik



Approximation of Minimum Convex Partitioning

software project and competition 2019/20



Agenda

- 1. Introduction and Overview
- 2. DCEL
- 3. Nested Hulls
- 4. Single Convex Waves
- 5. Merged Convex Waves
- 6. Pass based
- 7. Start points
- 8. Solutions



1. Introduction and Overview

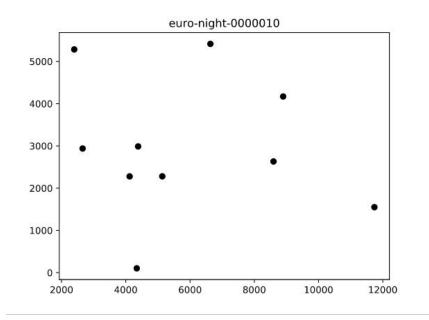


1. The CG Challenge 2020

- CG:SHOP = Computational Geometry Solving Hard Optimization Problems
- Part of the CG Week in Zurich (June 22-26)
- Open Class contest
- Opened: September 30
- Closes: February 14

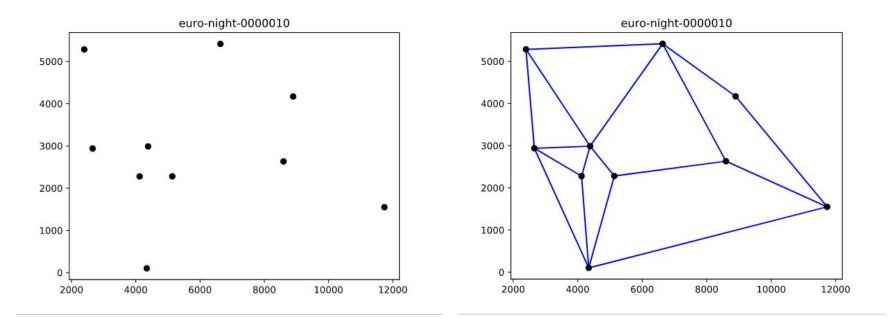


1. The Minimum Convex Partition Problem





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1. The Minimum Convex Partition Problem

- Complexity unknown
- At start: 247 instances
- Jan 21: 99 additional instances
- 4 types:
 - \circ uniform
 - edge
 - illumination
 - orthogonally collinear points
- Tiebreaker: Time



1. Workflow

- Language: Python
- Communication: Slack
- Repository: GitHub
- Team meetings every Wednesday



23.10.19 Algorithm conception and proof of concept



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6.11.19 Initial prototype



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Common interface specification



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- 6.11.19 Initial prototype Common interface specification
- 13.11.19 Multiple program runs



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- **13.11.19 Multiple program runs**
- 20.11.19 Baseline results



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 - **Common interface specification**
- **13.11.19 Multiple program runs**
- **20.11.19 Baseline results**





- 8.12.19 Alternative algorithms
 - Nested convex hulls



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 - Removing edges from triangulation



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- 25.12.19 Result comparison



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- 8.1.20 User Interface



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- 31.1.20 Miscellaneous improvements / alternatives



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- **31.1.20** Miscellaneous improvements / alternatives

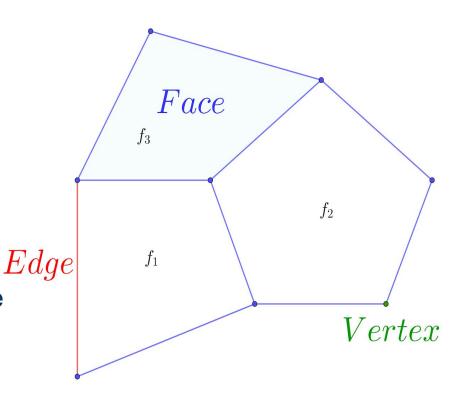


- Nested convex hulls
- Removing edges from triangulation
- Linear integer programming
- 25.12.19 Result comparison
- 8.1.20 User Interface
- **31.1.20** Miscellaneous improvements / alternatives
- 15.2.20 Contingency buffer



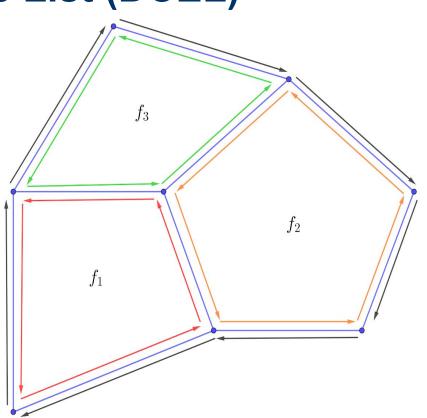


- Most commonly used representations for planar subdivisions
- It links together three sets of records:
 - > Vertex
 - ≻ Edge
 - ➤ Face
- It provides the ability of traversing the faces of planar subdivision, visiting all the edges around a given vertex

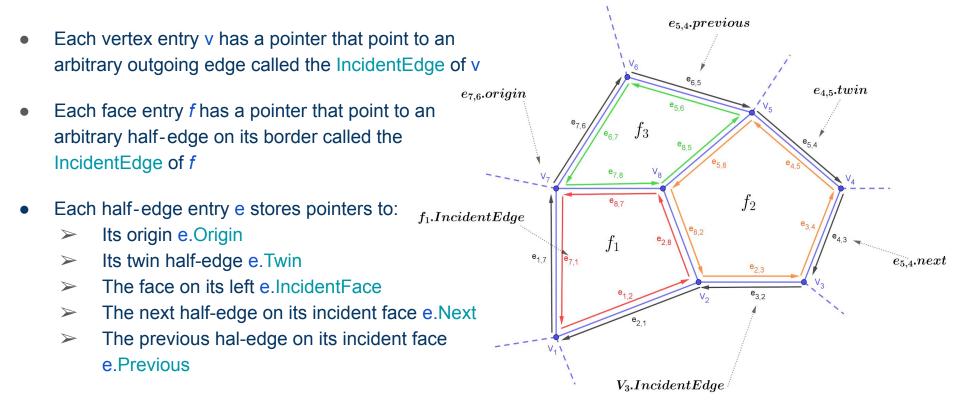




- Edges are oriented counterclockwise inside each face
- Each edge is a border between two faces, and is therefore represented by two half-edges, one for each face









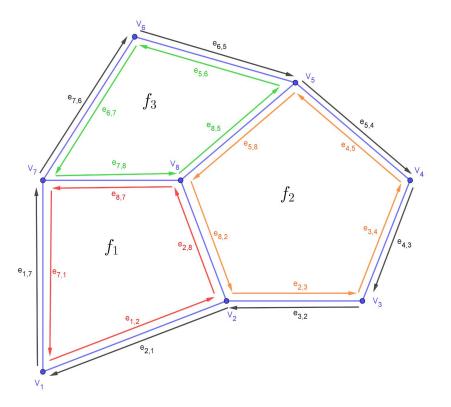
2. Doubly Connected Edge List (DCEL)

Vertex	Coordinates	IncidentEdge	
v ₁	(x ₁ ,y ₁)	e _{1,2}	
v ₂	(x ₂ ,y ₂)	e _{2,8}	

Face	Edge	
f ₁	e _{8,7}	
f ₂	e _{4,5}	

Half-edge	Origin	Twin	IncidentFace	Next	Previous
е _{6,7}	V ₆	е _{7,6}	f ₃	е _{7,8}	e _{5,6}
e _{5,8}	v ₅	e _{8,5}	f ₂	e _{8,2}	e _{4,5}

In our implementation of DCEL we excluded the faces table, as we did not need it.





3. Nested Convex-Hulls Approach



3. Nested Convex-Hulls Approach

- 1. Iteratively keep computing c-hulls:
 - 1. Compute the c-hull of all points in the data set
 - 2. Subtract data points of the computed c-hull from the data set
 - 3. Repeat 1.1 & 1.2 until we get an empty data set
- Connect each two sequential c-hulls in such a way that none of the added edges can be removed, unless we violate the convexity conditions
- 3. Except for the most outer c-hull, for each c-hull check for each edge if it can be removed



3. Nested Convex-Hulls, An Example "stars-000020"

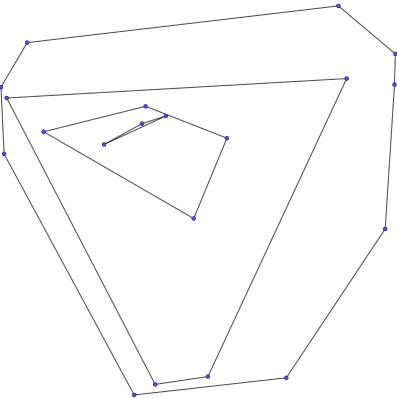


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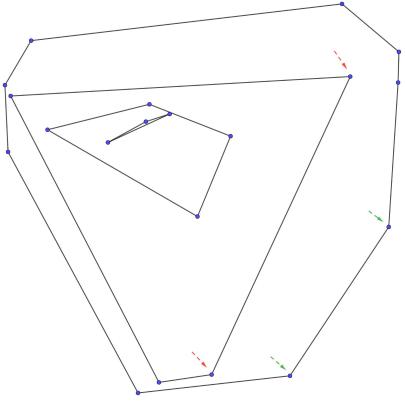
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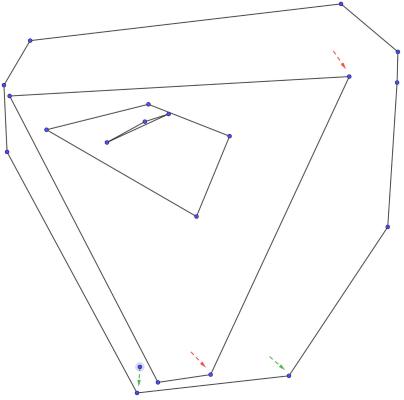
3. Example: Constructing C-Hulls



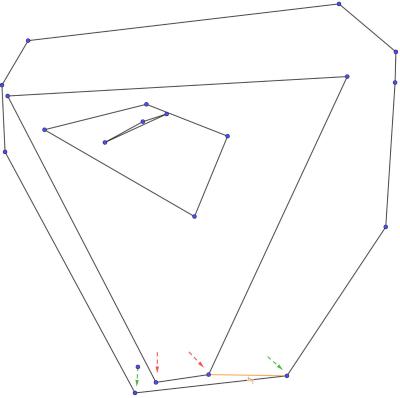




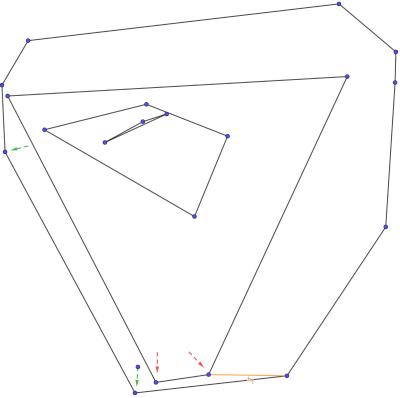




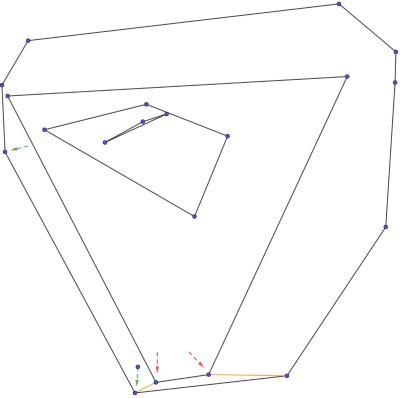




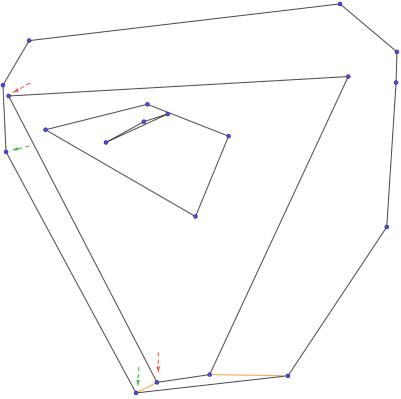




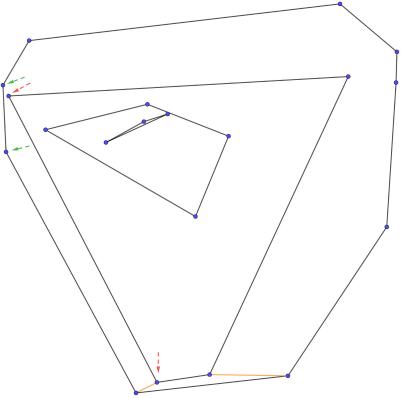




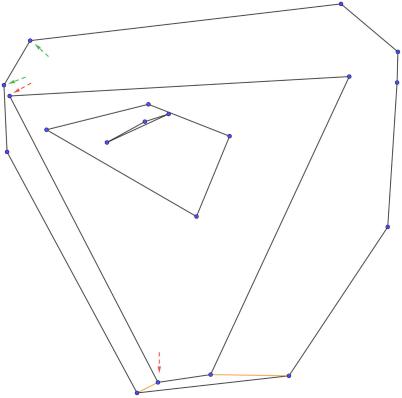




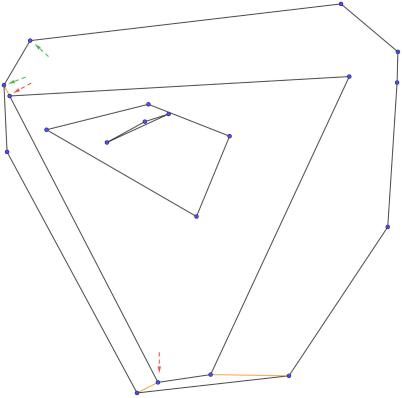




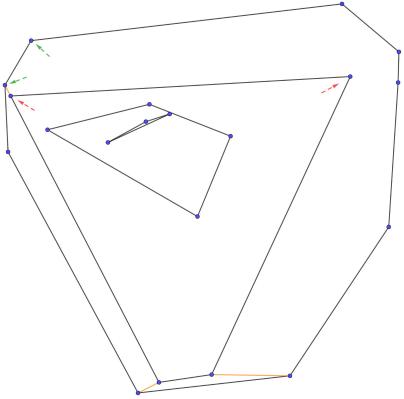




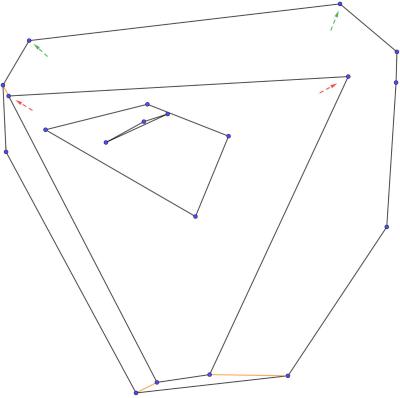




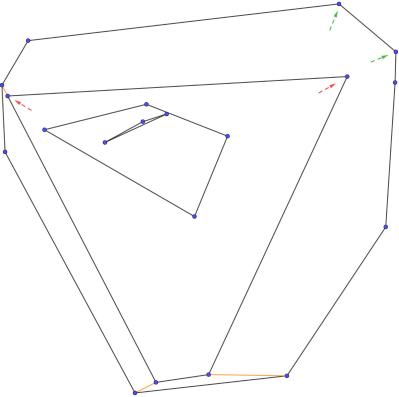




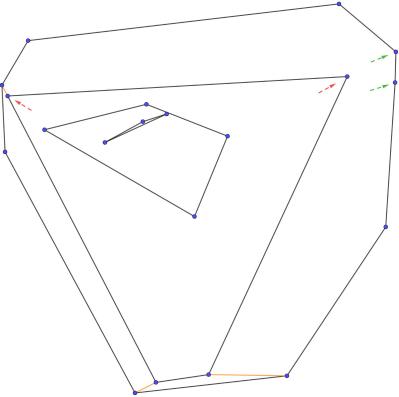




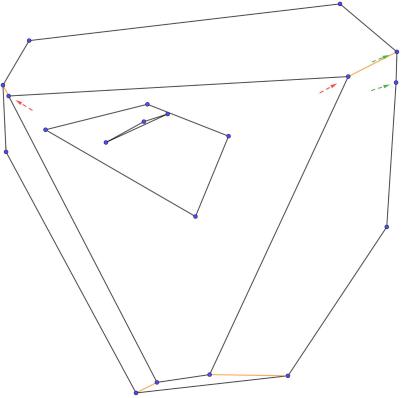




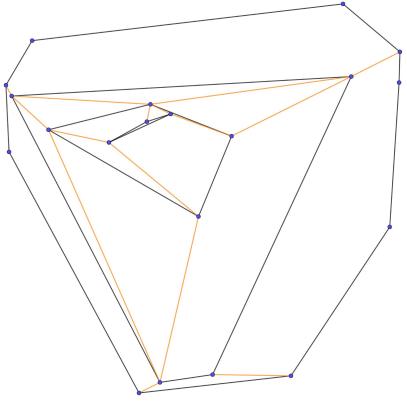




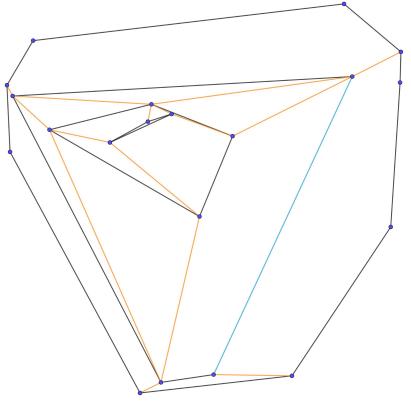




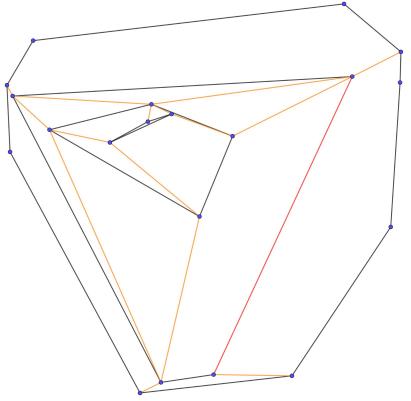




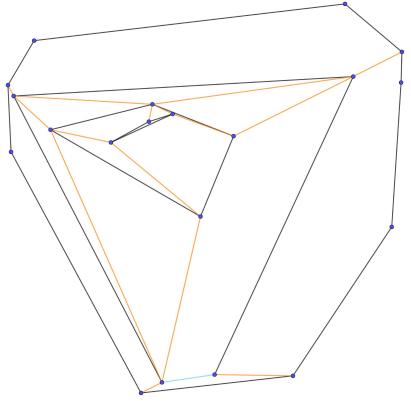




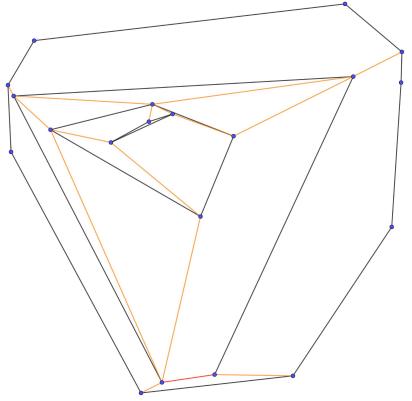




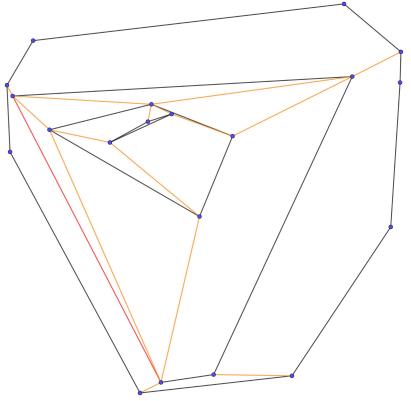




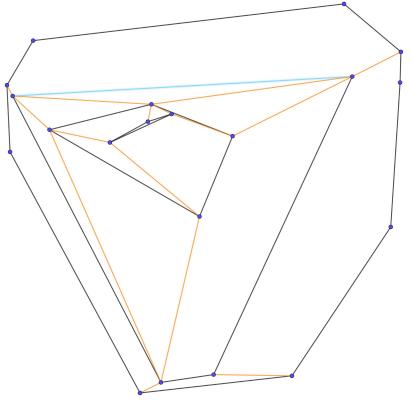




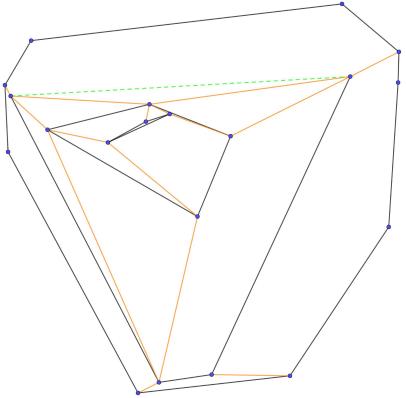




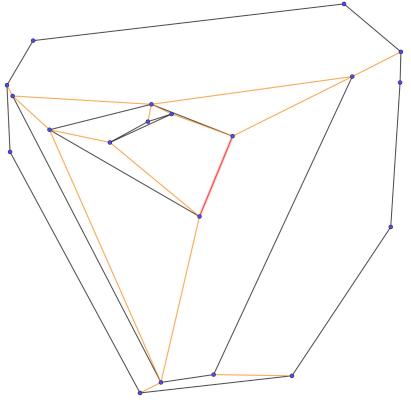




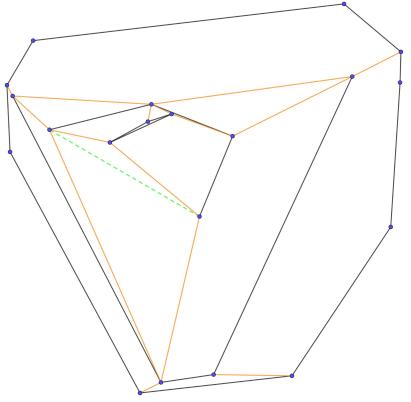




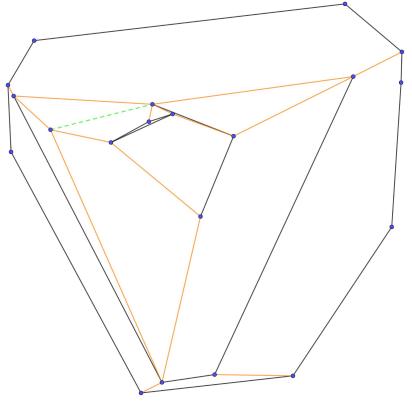




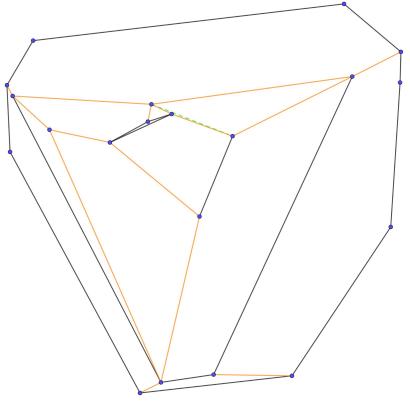




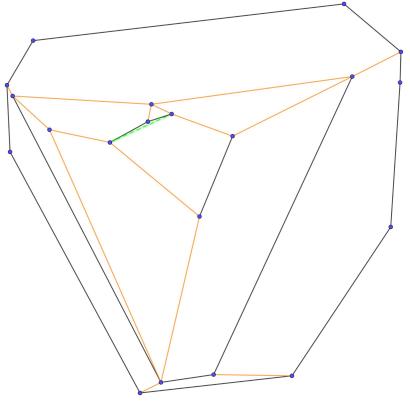




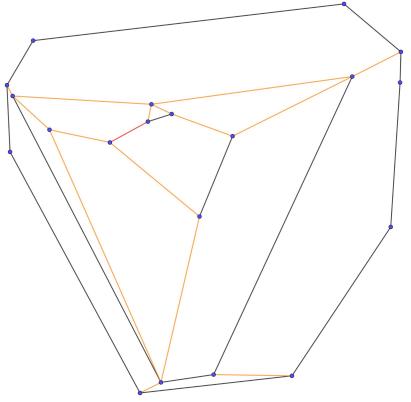




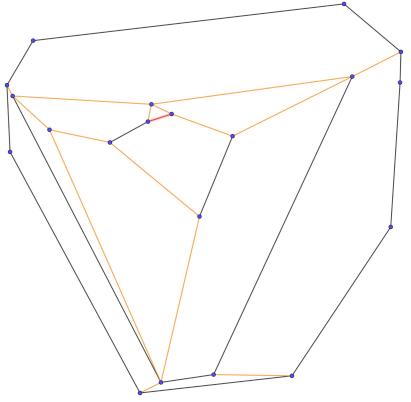






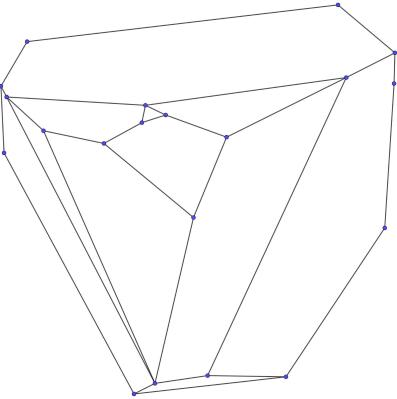








3. Example: Final Result



20 Vertices, 30 Edges & 11 Faces



3. Nested Convex-Hulls, An Upper Bound

For V to be the number of vertices, we have:

- When constructing nested c-hulls, we add at most V edges
- When connecting two c-hulls, we connect each vertex of the inner hull to at most 2 vertices of the outer hull, except when only one vertex left as last c-hull, which need to be connected to at most 3 vertices of the outer hull, so the worst case would be:
 - If the most outer hull is of size 3 and the most inner hull is of size 1, then for connecting hulls we add at most 2*(V-3) +1 edges
- Suppose in the deletion step no edge was eligible to be deleted

Then the max number of edges that can be added is $V + 2^{(V-3)+1} = 3V-5$ edges

*In practice: 2V - ~ 20% " 20% of 2V"



3. Nested Convex-Hulls, Run-Time

Sorting:O(nlog(n))Constructing Convex-Hulls : $O(n^2)$ Connecting Convex-Hulls :O(n)Deleting Edges:O(n)

 $O(n^2)$





Sort by distance to *startpoint* $\rightarrow Q$

First three points in $Q \rightarrow H$

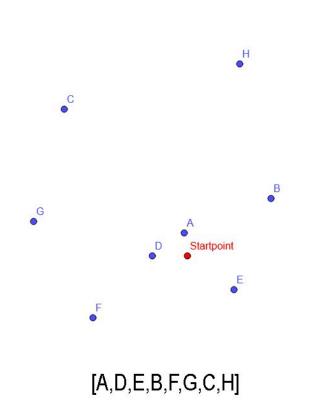
For each point p in Q:

Get visible bounds of p to H

Connect p to all points in-between

Remove redundant edges

Update H to the convex hull





Sort by distance to *startpoint* $\rightarrow Q$

First three points in $Q \rightarrow H$

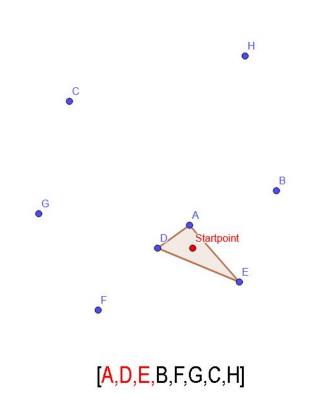
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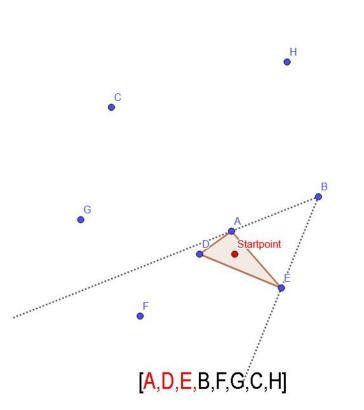
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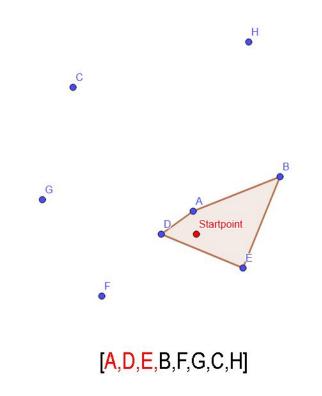


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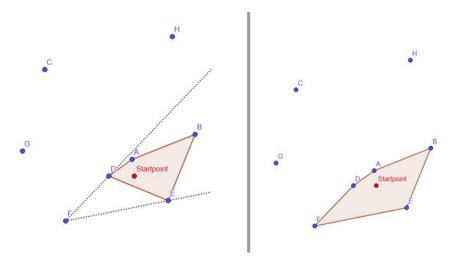
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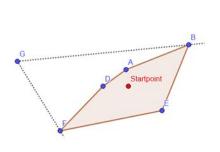
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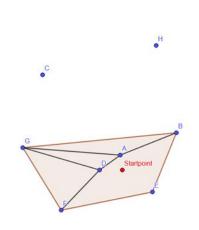
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H





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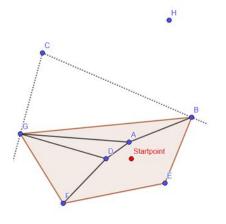
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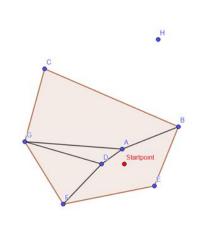
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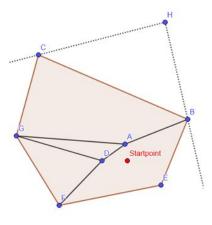
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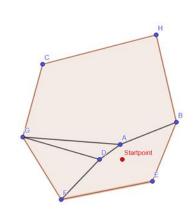
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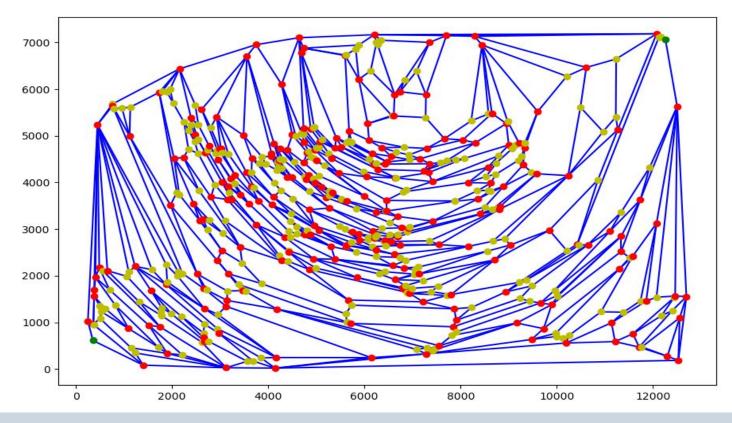
Remove redundant edges

Update *H* to the convex hull





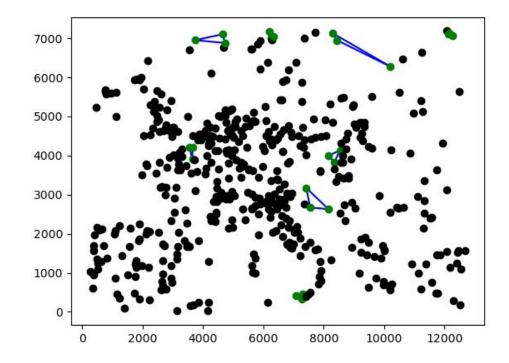




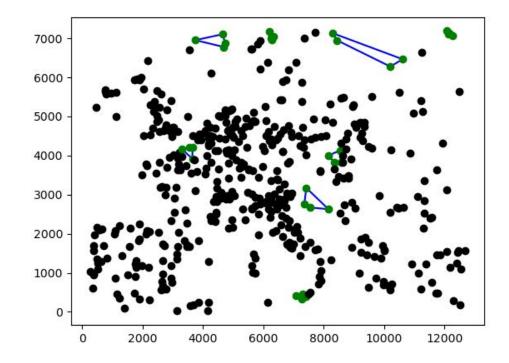


- Perform a convex wave for each startpoint
- Merge two waves on collision

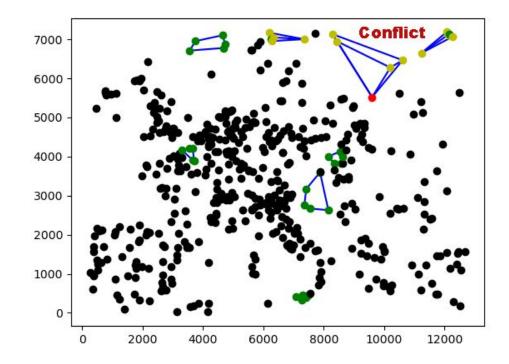




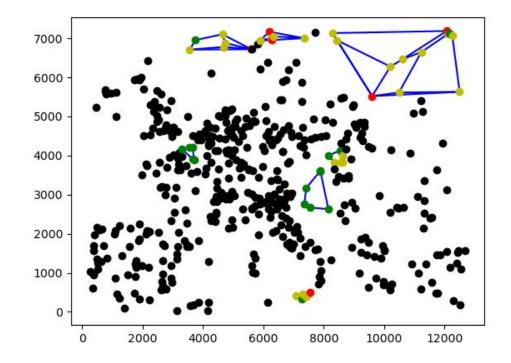




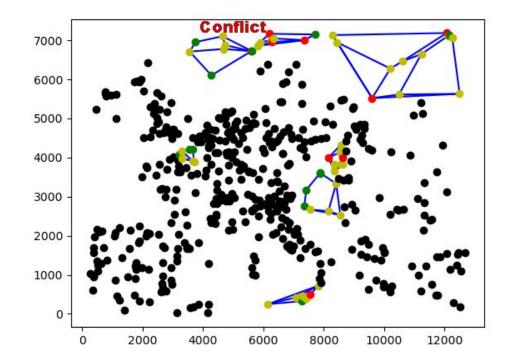




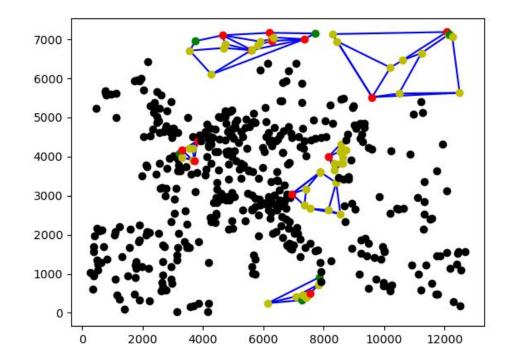




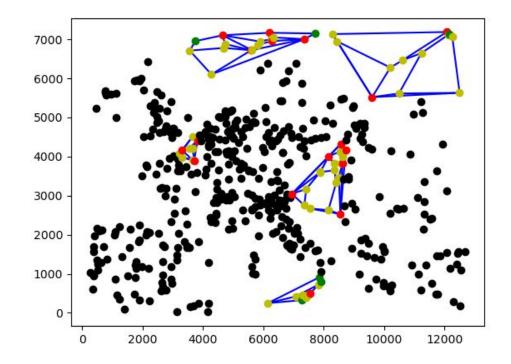




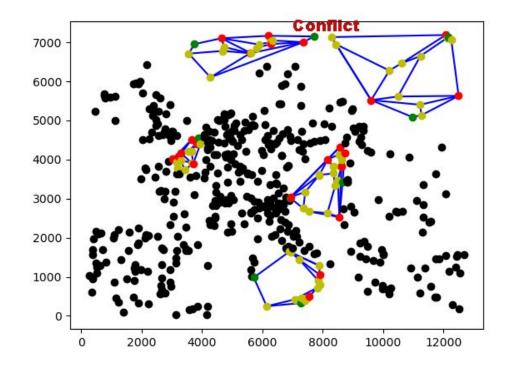




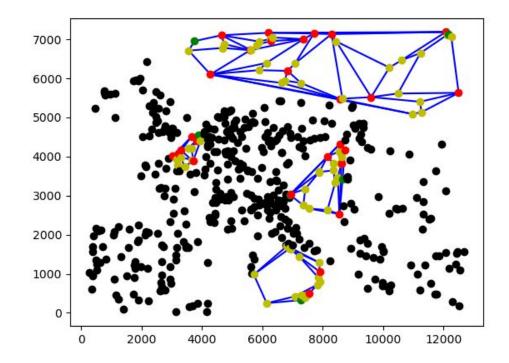




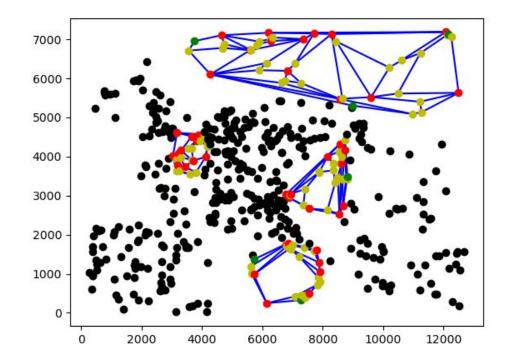




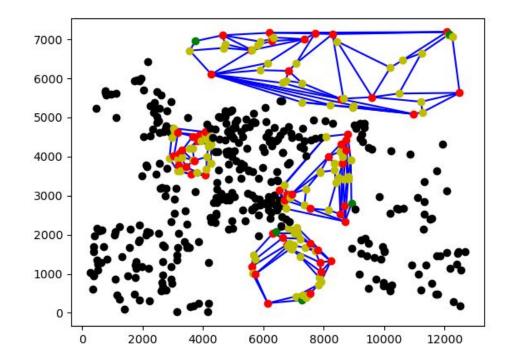




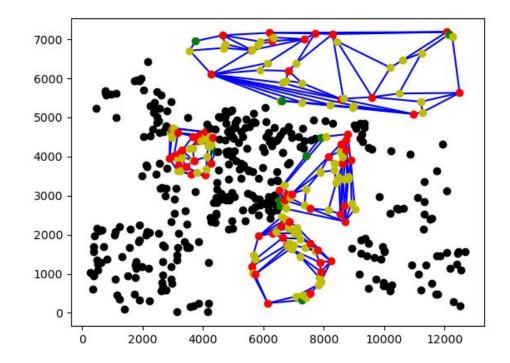




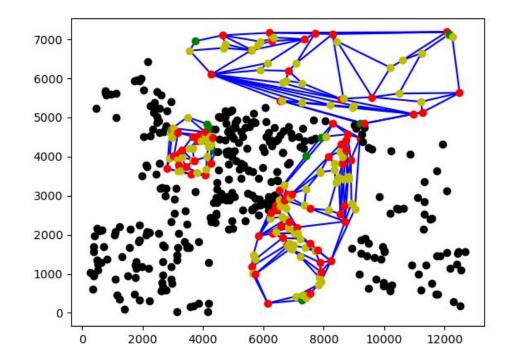




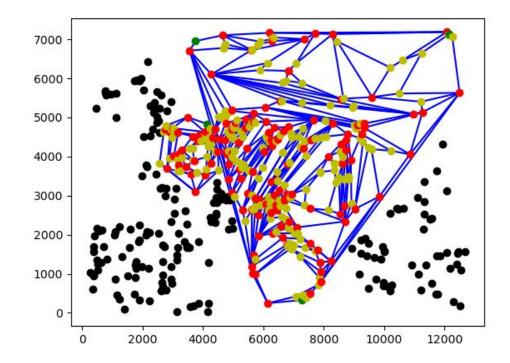




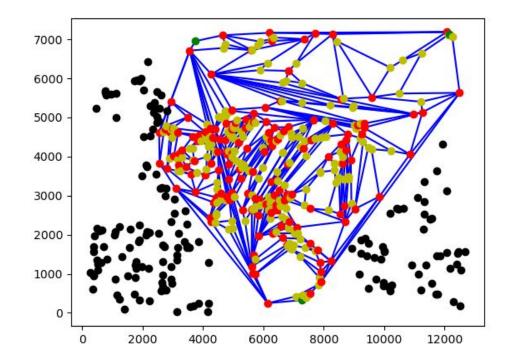




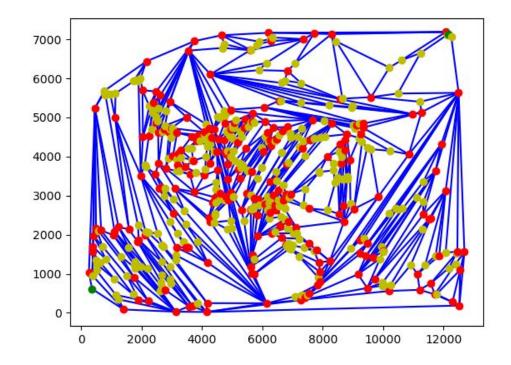














- Good results in starting areas, poor results everywhere else
- Merged instances lead to stretched polygons and long edges
- Convex hulls broken during merge need to be triangulated
- Produces more edges than other algorithms on almost all instances

Approach discarded

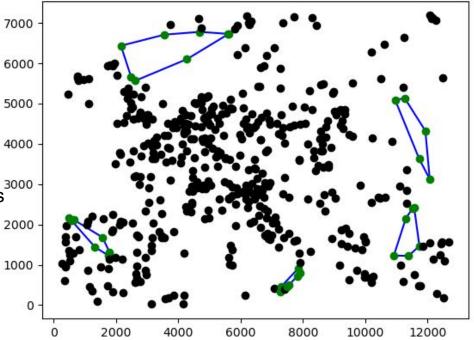




- Perform a set of independent passes
- Prioritize areas around startpoints
- No complex merging step required
- Waves constrained to a single convex polygon

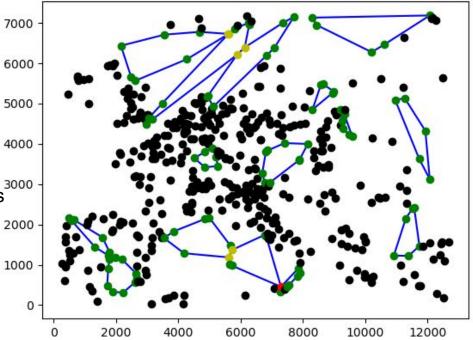


- Start a convex wave at each startpoint 6000
- Better startpoints have higher priority
- Only add a point if...
 - ...it can only see a single edge
 - ...it is not occluded by other points³⁰⁰⁰
 or edges²⁰⁰⁰



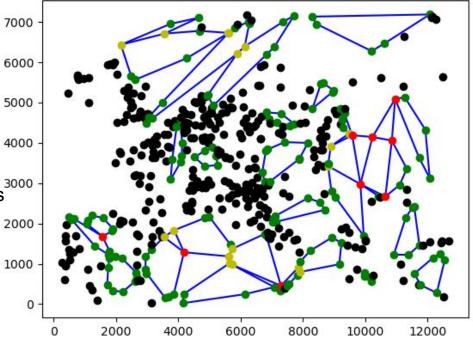


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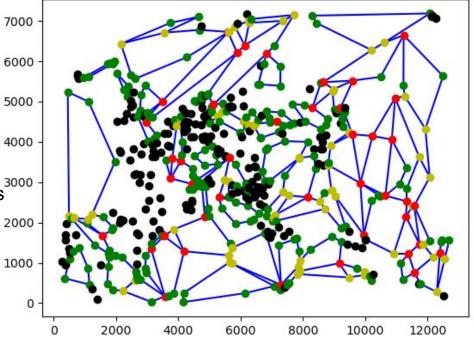


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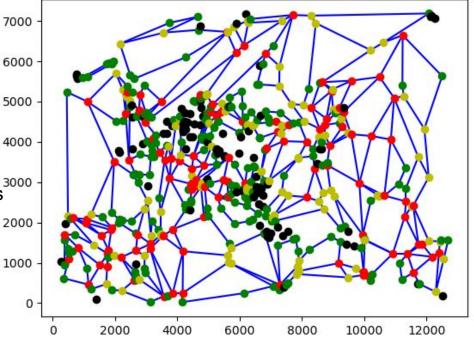


First Pass: Secure startpoints

- Start a convex wave at each startpoint • 6000 -
- Better startpoints have higher priority
- Only add a point if...
 - ...it can only see a single edge 0
 - 3000 ...it is not occluded by other points Ο 2000 or edges

1000

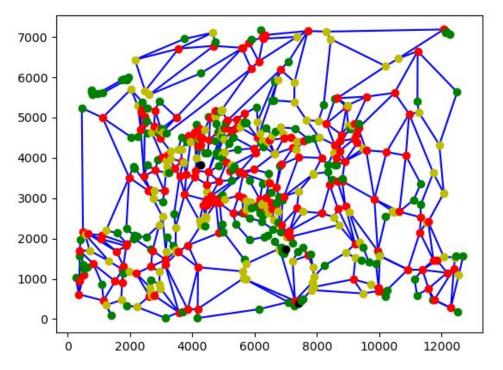
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Second Pass: Gather stray points

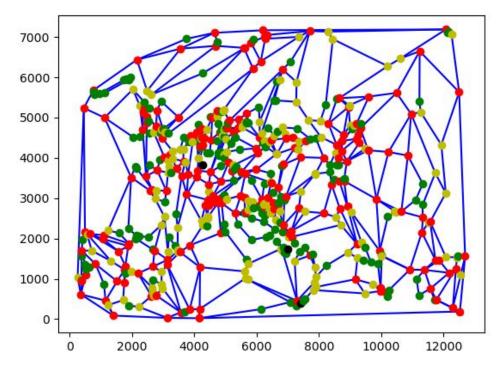
 Start a convex wave at each remaining stray vertex





Intermediate Pass: Convex Hull

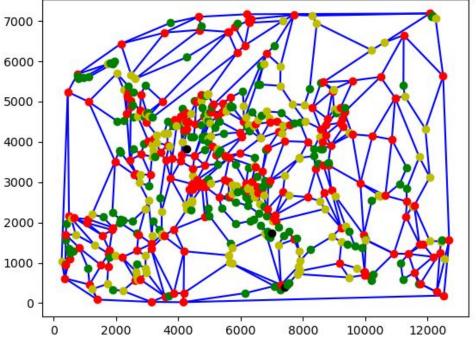
• Incorporate convex hull





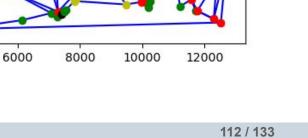
Intermediate Pass: Integrate islands

- Find islands via DFS on leftmost vertex 6000
- Connect each to their surrounding face 5000





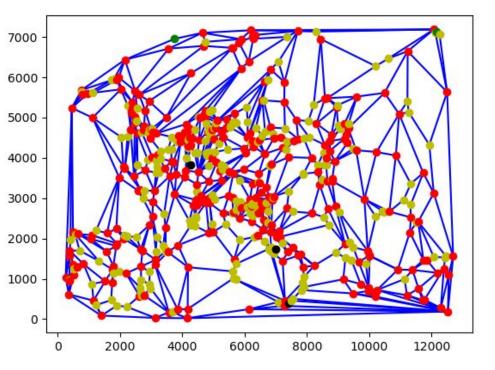
Intermediate Pass: Integrate islands Find islands via DFS on leftmost vertex 6000 Connect each to their surrounding face 5000 •





Third Pass: Resolve inflexes

- Find all inflex vertices
- Resolve these by connecting them to 1-2 opposing vertices



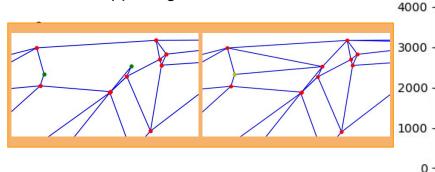


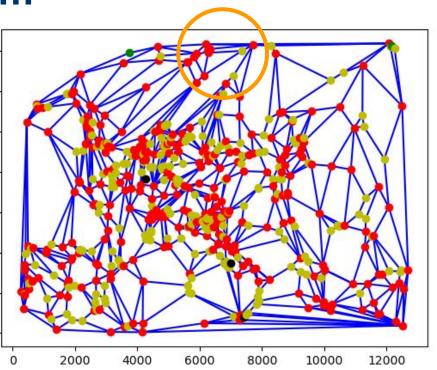
7000

6000

Third Pass: Resolve inflexes

- Find all inflex vertices
- Resolve these by connecting them to 1-2 opposing vertices

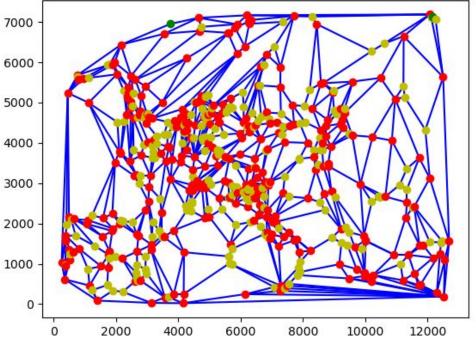






Intermediate Pass: Integrate stray points

- Find any remaining stray points as well 6000
 as their respective convex face 5000
- Incorporate them by iterating around the surrounding face



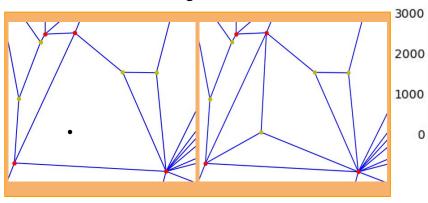


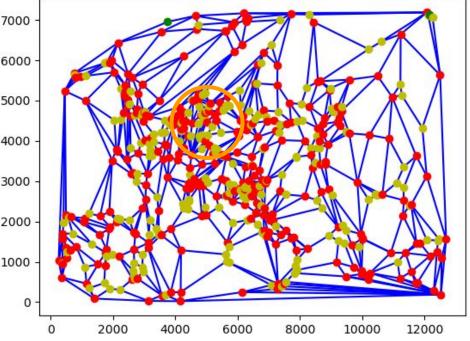
Intermediate Pass: Integrate stray points

Find any remaining stray points as well 6000 • as their respective convex face 5000

4000

Incorporate them by iterating around • the surrounding face

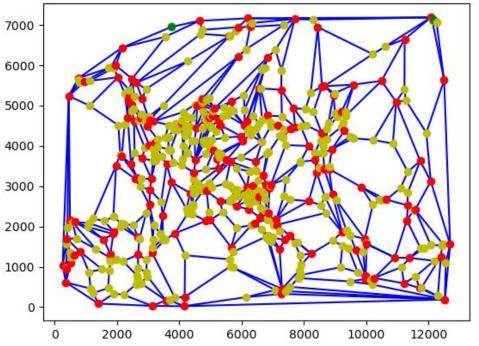






Fourth Pass: Clean

- Iterate over all edges and verify that they are required
- Remove the ones that are not

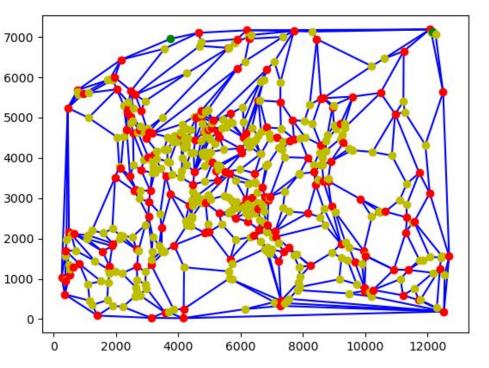




Fourth Pass: Clean

- Iterate over all edges and verify that they are required 50
- Remove the ones that are not

How do we chose adequate startpoints?

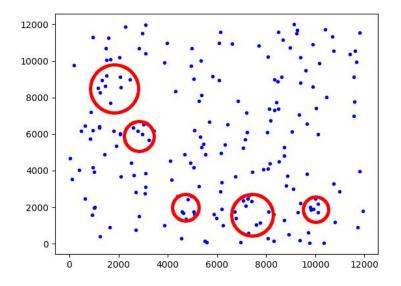




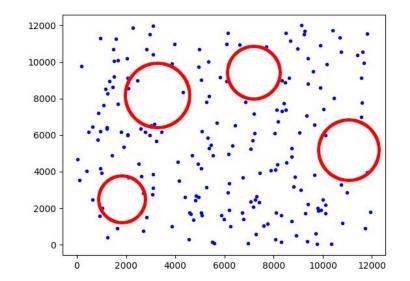
7. Startpoints



7. Startpoints - What are Good Start Points?



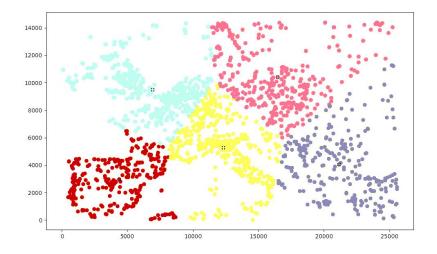
starting within clusters



starting in empty spaces



7. Startpoints - Dropped Concepts



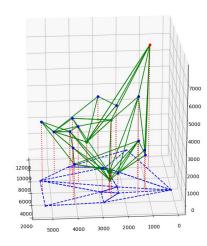
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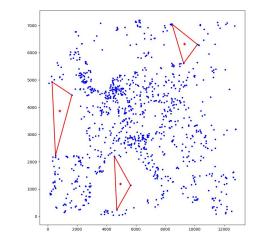
clustering with kmeans

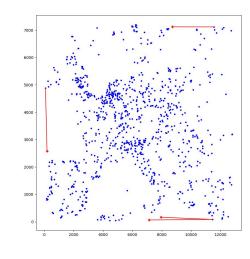
empty spaces with fixed-distance grid



7. Startpoints - Promising Concepts







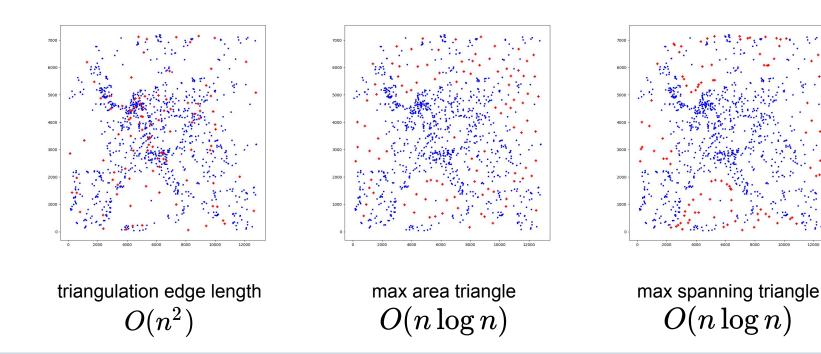
triangulation edge length $O(n^2)$

max area triangle $O(n \log n)$

max spanning triangle $O(n \log n)$



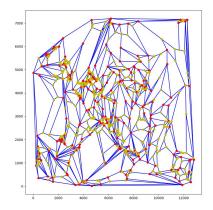
7. Startpoints - Distribution



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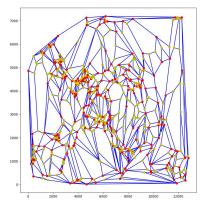


7. Startpoints - Was It Worth It?



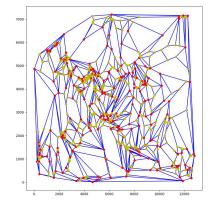
triangulation edge length 1254 edges in solution 2082 start points

 $O(n^2)$



max area triangle 1250 edges in solution 1383 start points

 $O(n \log n)$



max spanning triangle 1253 edges in solution 1383 start points

 $O(n \log n)$

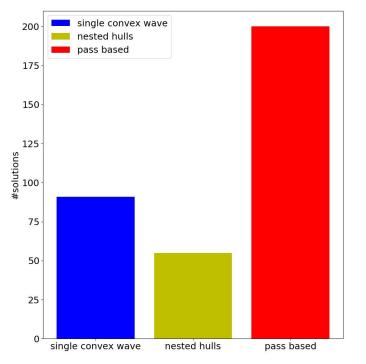
random 1242 edges in solution 4 start points

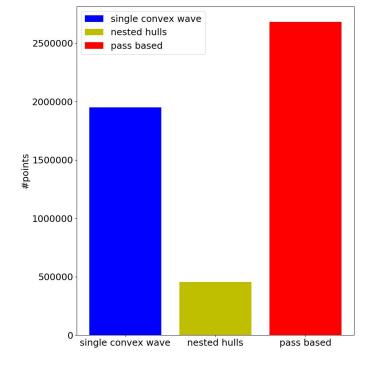


8. Solutions



8. Solutions - Comparing the Algorithms



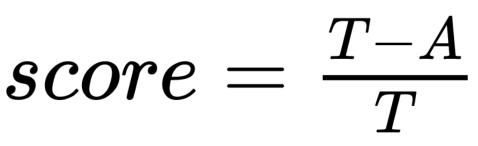


Solved Instances by Algorithm

Quantity of Points solved by Algorithm



8. Solutions - Score

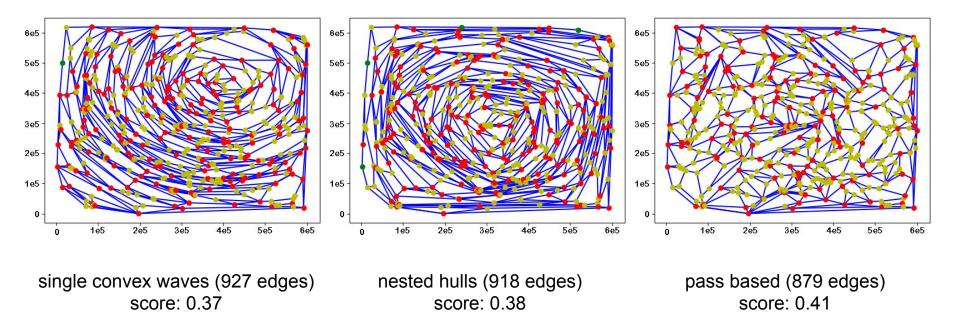


- T = #edges in triangulation
- *A* [≙] #edges in solution
- 0 < *score* < 1
- % of deleted edges from triangulation
- bigger score is better



8. Solutions - Plotting the Algorithms

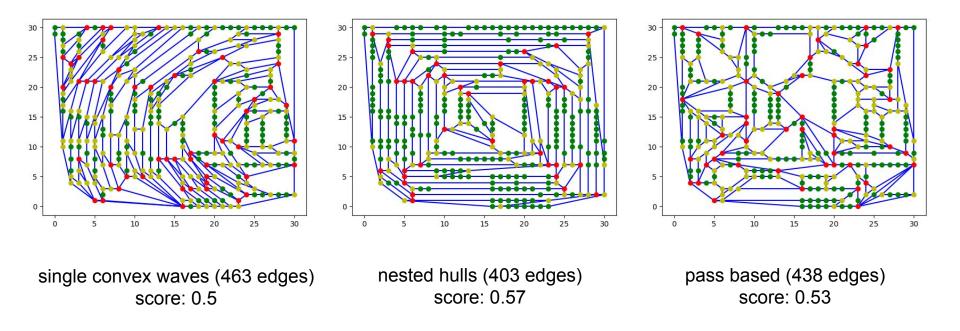
#points: 500, Triangulation #edges: 1480





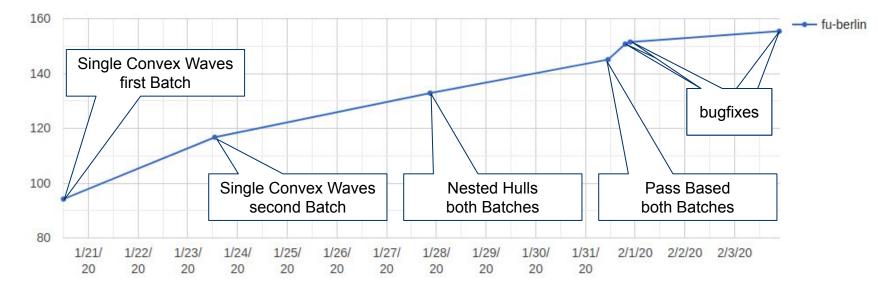
8. Solutions - Many Collinear Points

#points: 326, Triangulation #edges: 932





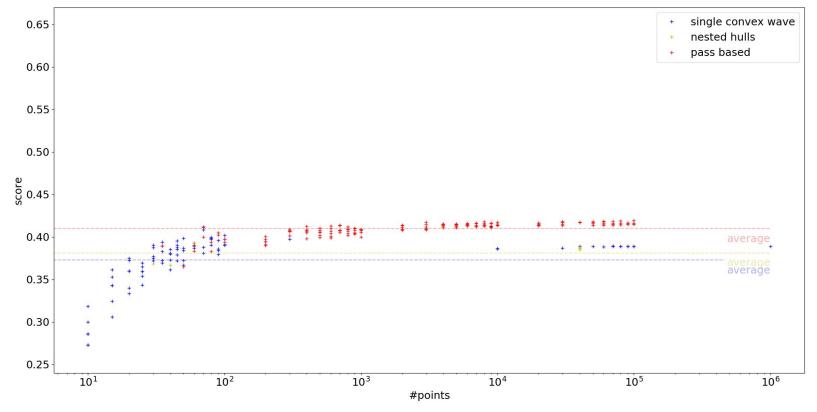
8. Solution



Final score: 155.432 - Instances: 346 - deleted Edges: 44,9%



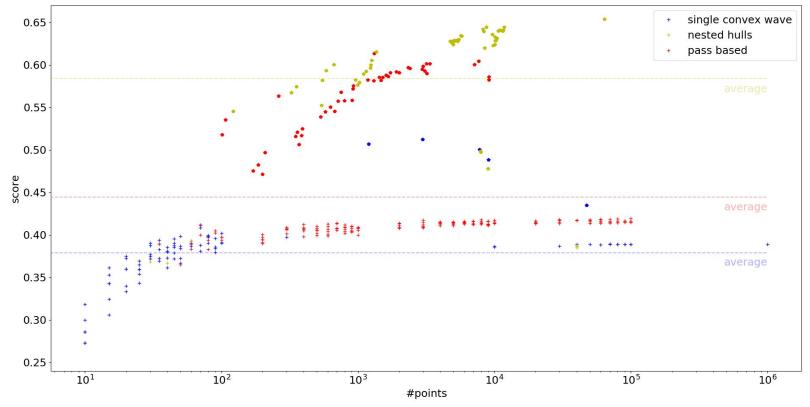
8. Solutions - Score Distribution



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8. Solutions - Score (Many Collinear Points)



132 / 133



Thank You For Your Attention

