# Approximation of Minimum Convex Partitioning 

software project and competition 2019/20

## Agenda

1. Introduction and Overview
2. DCEL
3. Nested Hulls
4. Single Convex Waves
5. Merged Convex Waves
6. Pass based
7. Start points
8. Solutions

## 1. Introduction and Overview

## 1. The CG Challenge 2020

- CG:SHOP = Computational Geometry - Solving Hard Optimization Problems
- Part of the CG Week in Zurich (June 22-26)
- Open Class contest
- Opened: September 30
- Closes: February 14


## 1. The Minimum Convex Partition Problem



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## 1. The Minimum Convex Partition Problem

- Complexity unknown
- At start: 247 instances
- Jan 21: 99 additional instances
- 4 types:
- uniform
- edge
- illumination
- orthogonally collinear points
- Tiebreaker: Time


## 1. Workflow

- Language: Python
- Communication: Slack
- Repository: GitHub
- Team meetings every Wednesday


## 1. Project Roadmap

23.10.19 Algorithm conception and proof of concept

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Common interface specification

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20.11.19 Baseline results

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### 8.12.19 Alternative algorithms

- Nested convex hulls


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- Removing edges from triangulation


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- Nested convex hulls
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31.1.20 Miscellaneous improvements / alternatives


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8.12.19 Alternative algorithms

- Nested convex hulls
- Removing edges from triangulation
- Linear integer programming
25.12.19 Result comparison
8.1.20 User Interface
31.1.20 Miscellaneous improvements / alternatives
15.2.20 Contingency buffer


## 2. Doubly Connected Edge List (DCEL)

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- Most commonly used representations for planar subdivisions
- It links together three sets of records:
$>$ Vertex
$>$ Edge
$>$ Face
- It provides the ability of traversing the faces of planar subdivision, visiting all the edges around a given vertex



## 2. Doubly Connected Edge List (DCEL)

- Edges are oriented counterclockwise inside each face
- Each edge is a border between two faces, and is therefore represented by two half-edges, one for each face



## 2. Doubly Connected Edge List (DCEL)

- Each vertex entry v has a pointer that point to an arbitrary outgoing edge called the IncidentEdge of $v$
- Each face entry $f$ has a pointer that point to an arbitrary half-edge on its border called the IncidentEdge of $f$
- Each half-edge entry e stores pointers to:
$>$ Its origin e.Origin
$>$ Its twin half-edge e.Twin
$>$ The face on its left e.IncidentFace
$>$ The next half-edge on its incident face e.Next
$>$ The previous hal-edge on its incident face e.Previous
$e_{5,4}$.previous
$e_{7,6}$. origin


## 2. Doubly Connected Edge List (DCEL)

| Vertex | Coordinates | IncidentEdge |
| :---: | :---: | :---: |
| $\mathbf{v}_{1}$ | $\left(\mathbf{x}_{1}, \mathbf{y}_{1}\right)$ | $\mathbf{e}_{1,2}$ |
| $\mathbf{v}_{2}$ | $\left(\mathbf{x}_{2}, \mathbf{y}_{2}\right)$ | $\mathbf{e}_{2,8}$ |
| $\ldots$ | $\ldots$ | $\ldots$ |


| Face | Edge |
| :---: | :---: |
| $\mathbf{f}_{1}$ | $\mathbf{e}_{8,7}$ |
| $\mathbf{f}_{2}$ | $\mathbf{e}_{4,5}$ |
| $\ldots$ | $\ldots$ |


| Half-edge | Origin | Twin | IncidentFace | Next | Previous |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{e}_{6,7}$ | $\mathbf{v}_{6}$ | $\mathbf{e}_{7,6}$ | $\mathbf{f}_{3}$ | $\mathbf{e}_{7,8}$ | $\mathbf{e}_{5,6}$ |
| $\mathbf{e}_{5,8}$ | $\mathbf{v}_{5}$ | $\mathbf{e}_{8,5}$ | $\mathbf{f}_{2}$ | $\mathbf{e}_{8,2}$ | $\mathbf{e}_{4,5}$ |
| $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ |

*In our implementation of DCEL we excluded the faces table, as we did not need it.

## 3. Nested Convex-Hulls Approach

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1. Iteratively keep computing c-hulls:
2. Compute the c-hull of all points in the data set
3. Subtract data points of the computed c-hull from the data set
4. Repeat 1.1 \& 1.2 until we get an empty data set
5. Connect each two sequential c-hulls in such a way that none of the added edges can be removed, unless we violate the convexity conditions
6. Except for the most outer c-hull, for each c-hull check for each edge if it can be removed

## 3. Nested Convex-Hulls, An Example

## 3. Example: Constructing C-Hulls



## 3. Example: Connecting Sequential Hulls



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## 3. Example: Removing Unneeded Edges

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## 3. Example: Removing Unneeded Edges

## 3. Example: Final Result



## 3. Nested Convex-Hulls, An Upper Bound

For V to be the number of vertices, we have:

- When constructing nested c-hulls, we add at most $V$ edges
- When connecting two c-hulls, we connect each vertex of the inner hull to at most 2 vertices of the outer hull, except when only one vertex left as last c-hull, which need to be connected to at most 3 vertices of the outer hull, so the worst case would be:
- If the most outer hull is of size 3 and the most inner hull is of size 1 , then for connecting hulls we add at most $2 *(\mathrm{~V}-3)+1$ edges
- Suppose in the deletion step no edge was eligible to be deleted Then the max number of edges that can be added is $\mathrm{V}+2^{*}(\mathrm{~V}-3)+1=3 \mathrm{~V}-5$ edges

[^0]
## 3. Nested Convex-Hulls, Run-Time

Sorting:
$O(n \log (n))$
Constructing Convex-Hulls :
Connecting Convex-Hulls : $O(n)$

Deleting Edges:

Overall Run-Time:

$$
O\left(n^{2}\right)
$$

## 4. Convex Waves

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Sort by distance to startpoint $\rightarrow Q$
First three points in $Q \rightarrow H$
For each point $p$ in $Q$ :
Get visible bounds of $p$ to $H$
Connect $p$ to all points in-between
Remove redundant edges
Update $H$ to the convex hull
[A,D,E,B,F,G,C,H]

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## 4. Convex Waves



## 5. Merged Convex Waves

- Perform a convex wave for each startpoint
- Merge two waves on collision


## 5. Merged Convex Waves



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## 5. Merged Convex Waves

- Good results in starting areas, poor results everywhere else
- Merged instances lead to stretched polygons and long edges
- Convex hulls broken during merge need to be triangulated
- Produces more edges than other algorithms on almost all instances

Approach discarded

## 6. Pass-Based Algorithm

## 6. Pass-Based Algorithm

- Perform a set of independent passes
- Prioritize areas around startpoints
- No complex merging step required
- Waves constrained to a single convex polygon


## 6. Pass-Based Algorithm

First Pass: Secure startpoints

- Start a convex wave at each startpoint
- Better startpoints have higher priority
- Only add a point if...
- ...it can only see a single edge
- ...it is not occluded by other points
or edges



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## 6. Pass-Based Algorithm

First Pass: Secure startpoints

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## 6. Pass-Based Algorithm

Second Pass: Gather stray points

- Start a convex wave at each remaining stray vertex


## 6. Pass-Based Algorithm

Intermediate Pass: Convex Hull

- Incorporate convex hull



## 6. Pass-Based Algorithm

Intermediate Pass: Integrate islands

- Find islands via DFS on leftmost vertex 6000
- Connect each to their surrounding face



## 6. Pass-Based Algorithm

Intermediate Pass: Integrate islands


## 6. Pass-Based Algorithm

Third Pass: Resolve inflexes

- Find all inflex vertices
- Resolve these by connecting them to 1-2 opposing vertices



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## 6. Pass-Based Algorithm

Intermediate Pass: Integrate stray points

- Find any remaining stray points as well 6000 as their respective convex face
- Incorporate them by iterating around the surrounding face



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## 6. Pass-Based Algorithm

## Fourth Pass: Clean

- Iterate over all edges and verify that they are required
- Remove the ones that are not



## 6. Pass-Based Algorithm

## Fourth Pass: Clean

- Iterate over all edges and verify that they are required
- Remove the ones that are not
$\Rightarrow$ How do we chose adequate startpoints?



## 7. Startpoints

## 7. Startpoints - What are Good Start Points?


starting within clusters

starting in empty spaces

## 7. Startpoints - Dropped Concepts


clustering with kmeans

empty spaces with fixed-distance grid

## 7. Startpoints - Promising Concepts


triangulation edge length

$$
O\left(n^{2}\right)
$$


max area triangle
$O(n \log n)$

max spanning triangle
$O(n \log n)$

## 7. Startpoints - Distrilbution


triangulation edge length

$$
O\left(n^{2}\right)
$$


max area triangle
$O(n \log n)$

max spanning triangle
$O(n \log n)$

## 7. Startpoints - Was It Worth It?


triangulation edge length
1254 edges in solution 2082 start points

$$
O\left(n^{2}\right)
$$


max area triangle 1250 edges in solution 1383 start points
$O(n \log n)$

max spanning triangle 1253 edges in solution 1383 start points
$O(n \log n)$

random 1242 edges in solution 4 start points
$O(1)$

## 8. Solutions

## 8. Solutions - Comparing the Algorithms



Solved Instances by Algorithm


Quantity of Points solved by Algorithm

## 8. Solutions - Score

$$
\text { score }=\frac{T-A}{T}
$$

- $\boldsymbol{T} \xlongequal[=]{\#}$ edges in triangulation
- $\boldsymbol{A} \xlongequal[=]{=}$ edges in solution
- $0<$ score $<1$
- \% of deleted edges from triangulation
- bigger score is better


## 8. Solutions - Plotting the Algorithms

\#points: 500, Triangulation \#edges: 1480


## 8. Solutions - Many Collinear Points

## \#points: 326, Triangulation \#edges: 932


single convex waves (463 edges)
score: 0.5

nested hulls (403 edges)
score: 0.57

pass based (438 edges)
score: 0.53

## 8. Solution



Final score: 155.432 - Instances: 346 - deleted Edges: 44,9\%

## 8. Solutions - Score Distribution



## 8. Solutions - Score (Many Collinear Points)



## Thank You For Your Attention




[^0]:    *In practice: $2 \mathrm{~V}-\sim 20 \%$ " $20 \%$ of 2 V "

