



Dr. Vishwanath Karad

**MIT WORLD PEACE
UNIVERSITY** | PUNE

TECHNOLOGY, RESEARCH, SOCIAL INNOVATION & PARTNERSHIPS

CSE2PM09A / AID2PM07A

Design and Analysis of Algorithm

DEPARTMENT OF COMPUTER SCIENCE AND ENGINEERING

Hashing

- Introduction to hashing
- Hash functions
- Collision resolution strategies- Open Addressing and Chaining
- Hash Table Overflow

Introduction to Hashing

Suppose that we want to store 10,000 students records (each with a 5-digit ID) in a given container.

- ❑ A linked list implementation would take $O(n)$ time.
- ❑ A height balanced tree would give $O(\log n)$ access time. Using an array of size 100,000 would give $O(1)$ access time but will lead to a lot of space wastage.

Is there some way that we could get $O(1)$ access without wasting a lot of space?

The answer is hashing.

Introduction to Hashing

□ Why Hashing?

- Sequential Search requires $O(n)$ Comparisons.
- It is not useful for large database.
- Binary Search requires less comparisons $O(n \log n)$
- But it requires data to be sorted.

Introduction to Hashing

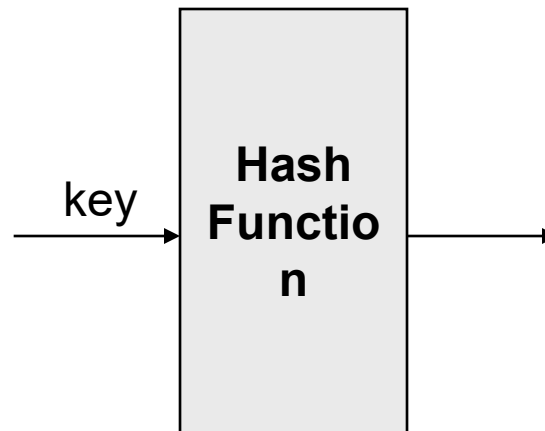
□ What is Hashing?

- Record for the key value is directly referred by calculating address from key value
- Address of element x is obtained by computing arithmetic function $f(x)$.
- Function $f(x)$ is called as **hash function**
- Table used for storing records is known as hash table
- **Best Case Time Complexity of hashing**= $O(1)$
- **Worst Case Time Complexity** = $O(n)$

Example

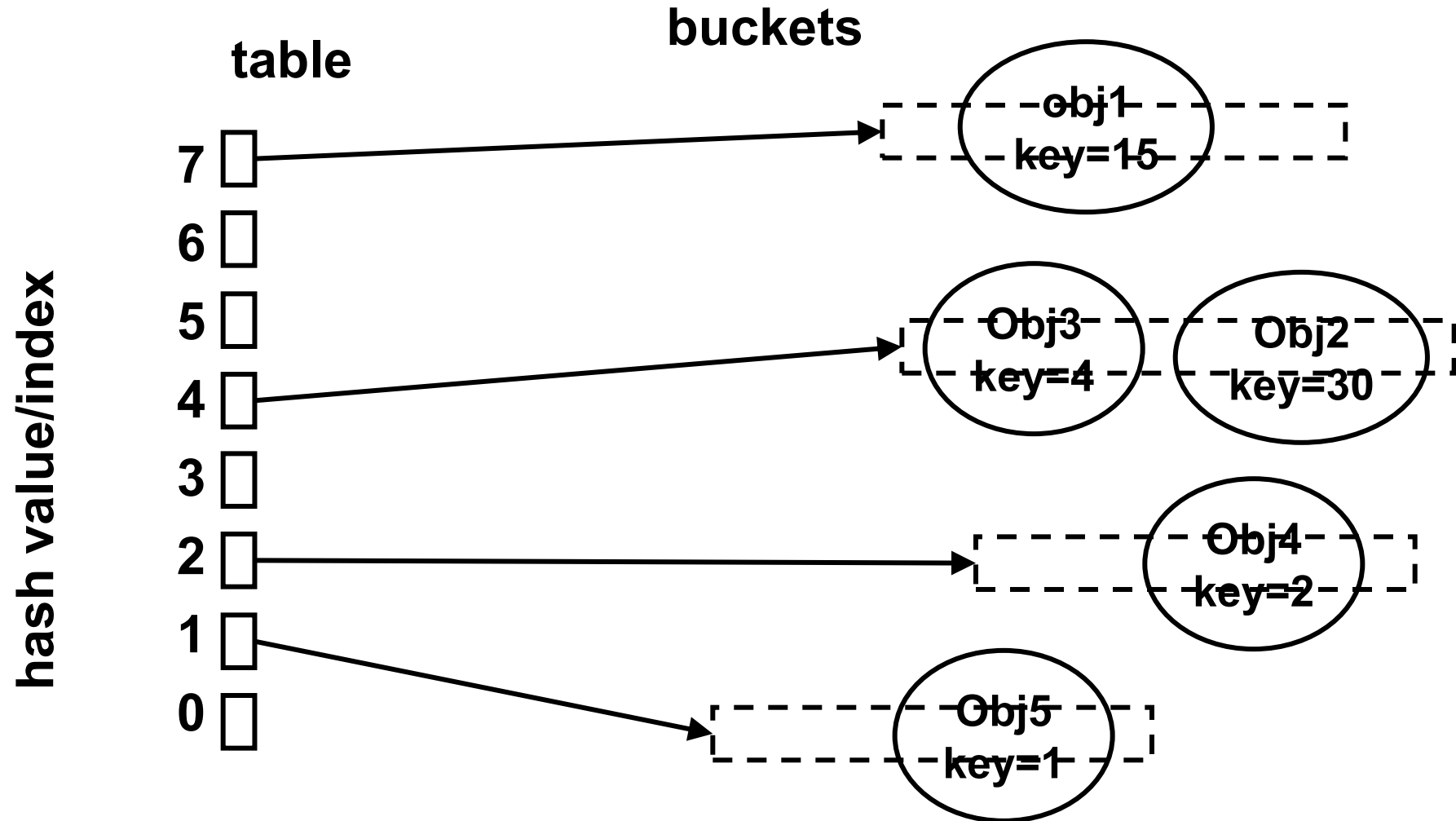
Items
john 25000
phil 31250
dave 27500
mary 28200

{
key



Hash Table	
0	
1	
2	
3	john 25000
4	phil 31250
5	
6	dave 27500
7	mary 28200
8	
9	

Hash Tables – Conceptual View



Key terms

- ❖ **Hash Table:** is an array of size $MAX[0 \text{ to } MAX-1]$
- ❖ **Hash Function:** that transforms a key into an address.
- ❖ **Bucket-** Bucket is an index position in hash table that can store more than one record
- ❖ When the same index is mapped with two keys, then both the records are stored in the same bucket
- ❖ **Probe-**Each action of address calculation and check for success is called probe
- ❖ **Collision-**The result of two keys hashing into the same address is called collision
- ❖ **Synonym-**Keys those hash to the same address are called synonyms

Key terms

- ❖ **Overflow**-The result of more keys hashing to the same address and if there is no room in the bucket, then it is said that overflow has occurred
- ❖ Collision and overflow are synonymous when the bucket is of size 1
- ❖ **Open or external hashing**-When we allow records to be stored in potentially unlimited space, it is called as open or external hashing
- ❖ **Closed or internal hashing**-When we use fixed space for storage eventually limiting the number of records to be stored, it is called as closed or internal hashing

Key Terms

Load density- The maximum storage capacity that is maximum number of records that can be accommodated is called as loading density

Full Table- Full table is the one in which all locations are occupied

Owing to the characteristics of hash functions, there are always empty locations

Load factor- the number of records stored in table divided by maximum capacity of table, expressed in terms of percentage

Hashing

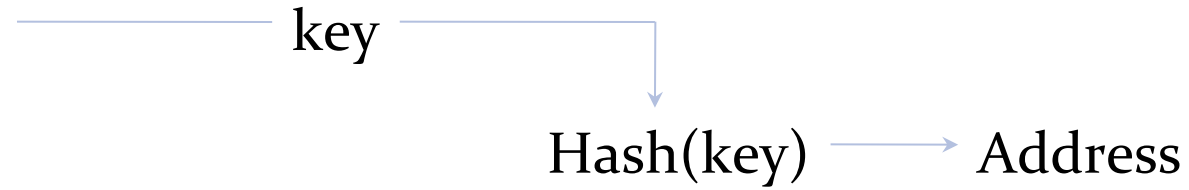
□ Example

31,42,35,67,24,19

$$F(x) = x \bmod 10$$

Location	X
1	31
2	42
3	
4	24
5	35
6	
7	67
8	
9	19

X	F(X)
31	1
42	2
35	5
67	7
24	4
19	9



-
- ❖ The resulting address is used as the basis for storing and retrieving records and this address is called as *home address of the record*
 - ❖ *For array to store a record in a hash table, hash function is applied to the key of the record being stored, returning an index within the range of the hash table*
 - ❖ The item is then stored in the table of that index position

Basic Idea

Use *hash function* to map keys into positions in a *hash table*

Ideally

If element e has key k and h is hash function, then e is stored in position $h(k)$ of table

To search for e , compute $h(k)$ to locate position. If no element, dictionary does not contain e .

Hash Functions

A **hash function** is any function that can be used to map a data set of an arbitrary size to a data set of a fixed size, which falls into the hash table.

The values returned by a hash function are called hash values, hash codes, hash sums, or simply hashes.

Example:

$$H(x) = x \bmod 10$$

Here, $H(x)$ is called as hash function while value returned by function is called as hash value.



Characteristics of a good hashing function

- The average performance of hashing depends on how the hash function distributes the set of keys among the slots
- Assumption is that any given record is equally likely to hash into any of the slots, independently of whether any other record has been already hashed to it or not
- This assumption is called as simple uniform hashing
- A good hash function is the one which satisfies the assumption of simple uniform hashing

Characteristics of a good hashing function

- ❖ Addresses generated from the key are uniformly and randomly distributed
- ❖ Small variations in the value of key will cause large variations in the record addresses to distribute records (with similar keys) evenly
- ❖ The hashing function must minimize the collision

Division Method

One of the required features of the hash function is that the resultant index must be within the table index range

One simple choice for a hash function is to use the modulus division indicated as MOD

The function returns an integer

If any parameter is NULL, the result is NULL

$$\text{Hash(Key)} = \text{Key} \% M$$

Hash Tables: Insert Example

For example, if we hash keys 0...1000 into a hash table with 5 entries and use $h(\text{key}) = \text{key} \bmod 5$, we get the following sequence of events:

Insert 2

	key	data
0		
1		
2	2	...
3		
4		

Insert 21

	key	data
0		
1	21	...
2	2	...
3		
4		

Insert 34

	key	data
0		
1	21	...
2	2	...
3		
4	34	...

Insert 54

There is a
collision at
array entry #4

???

Multiplication Method

The multiplication method works as:

1. Multiply the key 'k' by a constant A in the range $0 < A < 1$ and extract the fractional part of kA
2. Then multiply this value by M and take the floor of the result
 $\text{Hash}(k) = \text{M} (kA)$,

Donald Knuth suggested to use $A=0.61803398987$

Ex: key=107, assume $M=50$ key=123 $m=1000$

$$\begin{aligned} h(k) &= M * 107 * 0.61803398987 \\ &= 66.12 \end{aligned} \qquad \text{fractional part} = 0.12$$

$$\begin{aligned} h(k) &= 50 * 0.12 \\ &= 6 \\ h(k) &= 6 \end{aligned}$$

That means 107 will be placed at index 6 in hash table.

Multiplication Method

Disadvantage: Slower than the division method.

Advantage: Value of m is not critical.

- Typically chosen as a power of 2, i.e., $m = 2^p$, which makes implementation easy.

Example: $m = 1000$, $k = 123$, $A \approx 0.6180339887...$

$$\begin{aligned} h(k) &= \lfloor 1000(123 * 0.6180339887) \rfloor \\ &= \lfloor 1000 (76.018...) \rfloor \\ &= \lfloor 1000 (.018...) \rfloor = 18 \end{aligned}$$

Mid-Square Hashing

-
- ❖ The mid-square hashing suggests to take square of the key and extract the middle digits of the squared key as address
 - ❖ The difficulty is when the key is large. As the entire key participates in the address calculation, if the key is large, then it is very difficult to store the square of it as the square of key should not exceed the storage limit
 - ❖ So mid-square is used when the key size is less than or equal to 4 digits

Keys and addresses using mid-square

Key	Square	Hashed Address
2341	5480281	802
1671	2792241	922

The difficulty of storing larger numbers square can be overcome if for squaring we use few of digits of key instead of the whole key

Keys and addresses using mid-square

We can select a portion of key if key is larger in size and then square the portion of it

Key	Square	Hashed Address
234137	$234 * 234 = 54756$	475
567187	$567 * 567 = 321489$	148

Folding Technique

- ❖ In folding technique, the key is subdivided into subparts that are combined or folded and then combined to form the address
- ❖ For the key with digits, we can subdivide the digits in three parts, add them up, and use the result as an address.
- ❖ Here the size of subparts of key could be as that of the address

Folding Technique (contd...)

- ❖ There are two types of folding methods:
- ❖ **Fold shift** — Key value is divided into several parts of that of the size of the address. Left, right, and middle parts are added
- ❖ **Fold boundary** — Key value is divided into parts of that of the size of the address. Left and right parts are folded on fixed boundary between them and the centre part

Folding Technique

- ❖ For example, if the key is 987654321, it is understood as Left 987 Centre 654 Right 321
- ❖ For fold shift, addition is
- ❖
$$987 + 654 + 321 = 1962$$
- ❖ Now discard digit 1 and the address is 962
- ❖ For fold boundary, addition of reverse part is
- ❖
$$789 + 456 + 123 = 1368$$
- ❖ Discard digit 1 and the address is 368

Extraction Method

- ❖ When a portion of the key is used for the address calculation, the technique is called as the extraction method
- ❖ In digit extraction, few digits are selected and extracted from the key which are used as the address

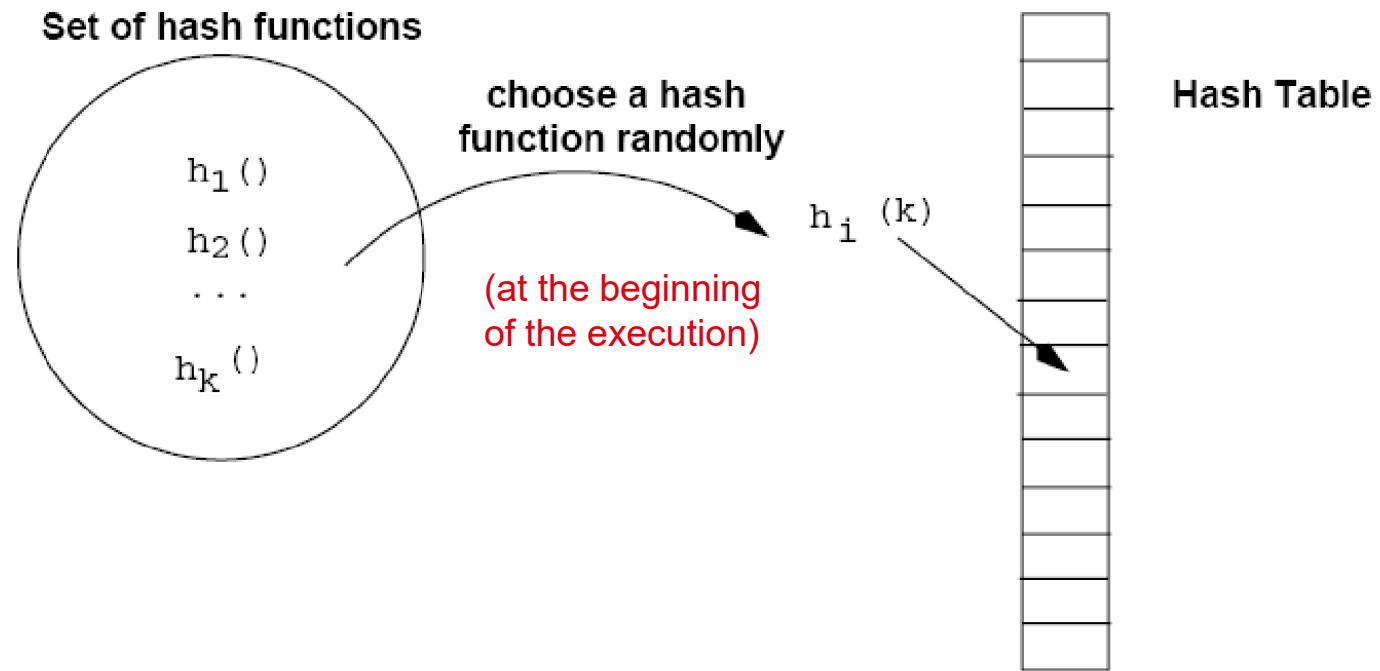
Keys and addresses using digit extraction

Key	Hashed Address
345678	357
234137	243
952671	927

Universal Hashing

-
- ❖ The main idea behind universal hashing is to select the hash function at random at run time from a carefully designed set of functions
 - ❖ Because of randomization, the algorithm can behave differently on each execution, even for the same input
 - ❖ This approach guarantees good average case performance, no matter what keys are provided as input

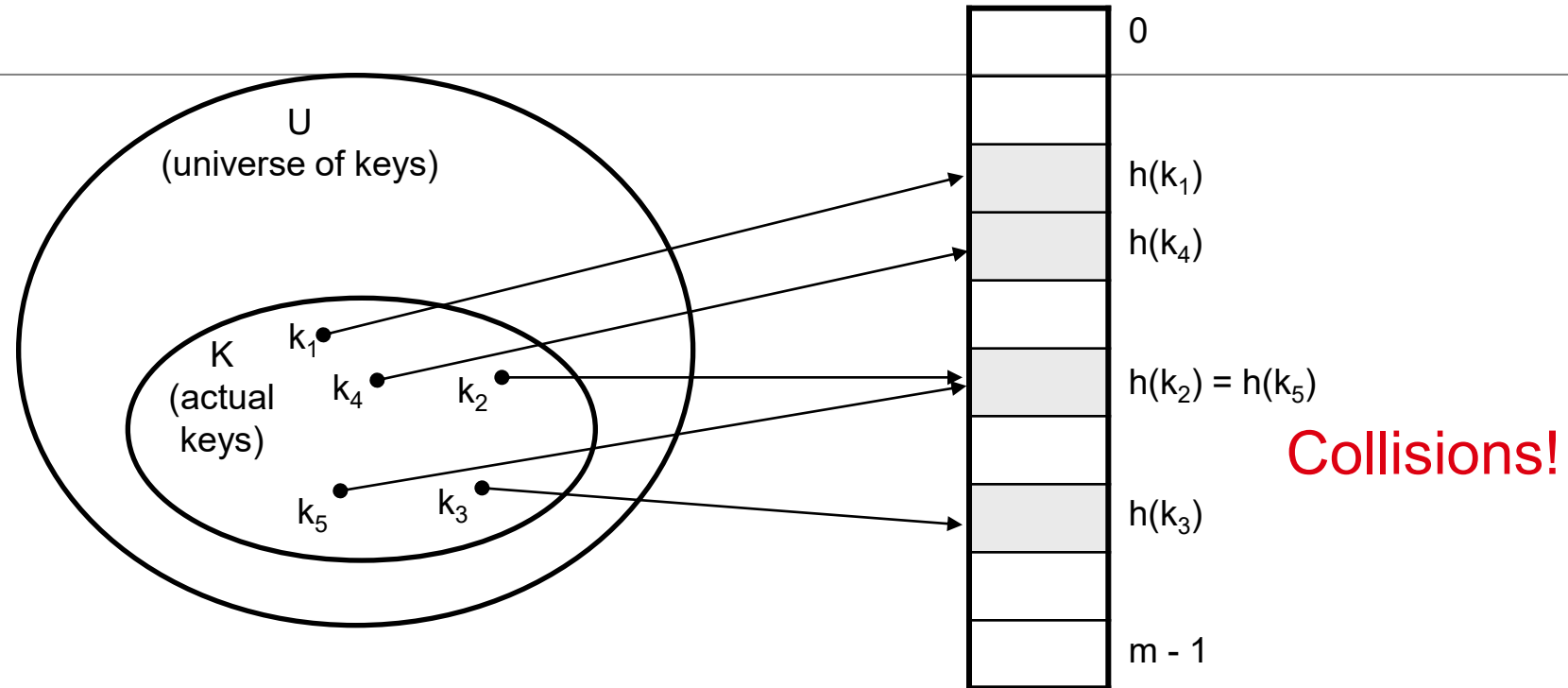
Universal Hashing



Issues

- ❖ A problem arises, however, when the hash function returns the same value when applied to two different keys
- ❖ To handle the situation, where two records need to be hashed to the same address we can implement a table structure, so as to have a room for two or more members at the same index positions

Do you see any problems with this approach?



Collision Resolution Strategies

Collision-Element to be inserted is mapped to same location

Example:

31,42,35,67,24,19,22

$F(x) = x \bmod 10$

Where to store 22?

Location	X
1	31
2	42
3	
4	24
5	35
6	
7	67
8	
9	19

Collision Resolution Strategies

Collision-Element to be inserted is mapped to same location

Example:

31,42,35,67,24,19,22

$F(x) = x \bmod 10$

Where to store 22?

Location	X
1	31
2	42
3	22
4	24
5	35
6	
7	67
8	
9	19

Collision Resolution Strategies

1. Open Addressing
 1. Linear probing
 2. Quadratic probing
 3. Double hashing
2. Separate chaining (or linked list)
3. Bucket hashing (defers collision but does not prevent it)

Open Addressing

When collision occurs, it is resolved by finding an available empty location other than the home address.

If Hash(key) is not empty, the positions are probed in the following sequence until an empty location is found.

When we reach the end of the table ,the search is wrapped around to start and search continues till the current collision location.

Closed hash tables use open addressing.

Linear Probing

- ❖ A hash table in which a collision is resolved by putting the item in the next empty place in following the occupied place is called linear probing
- ❖ This strategy looks for the next free location until it is found
- ❖ The function that we can use for probing linearly from the next location is as follows:
 - ❖ $(\text{Hash}(x) + p(i)) \text{ MOD Max}$
 - ❖ As $p(i) = i$ for linear probing, the function becomes
 - ❖ $(\text{Hash}(x) + i) \text{ MOD Max}$
 - ❖ Initially $i = 1$, if the location is not empty then it becomes 2, 3, 4, ..., and so on till empty location is found.

Linear Probing

- ❖ 1. With replacement 2. Without replacement

- ❖ **With replacement :**

- ❖ If the slot is already occupied by the key there are two possibilities, that is, either it is home address (collision) or not key's home address
- ❖ If the key's actual address is different, then the new key having the address at that slot is placed at that position and the key with other address is placed in the next empty position

Linear Probing

Example :

Given the input {4371, 1323, 6173, 4199, 4344, 9699, 1889} and hash function as $\text{Key} \% 10$, show the results for the following:

1. Open addressing using linear probing
2. Open addressing using quadratic probing
3. Open addressing using double hashing $h_2(x) = 7 - (x \text{ MOD } 7)$

Open addressing using linear probing with replacement

	Initial ly	Insert 4371	Insert 1323	Insert 6173	Insert 4199	Insert 4344	Insert 9699	Insert 1889
0							9699	9699
1		4371	4371	4371	4371	4371	4371	4371
2								1889
3			1323	1323	1323	1323	1323	1323
4				6173	6173	4344	4344	4344
5						6173	6173	6173
6								
7								
8								
9					4199	4199	4199	4199

Linear Probing

-
- ❖ **Without replacement :**
 - ❖ When some data is to be stored in hash table, and if the slot is already occupied by the key then another empty location is searched for a new record
 - ❖ There are two possibilities when location is occupied—it is its home address or not key's home address.
 - ❖ In both the cases, the without replacement strategy empty position is searched for the key that is to be stored

Open addressing using linear probing without replacement

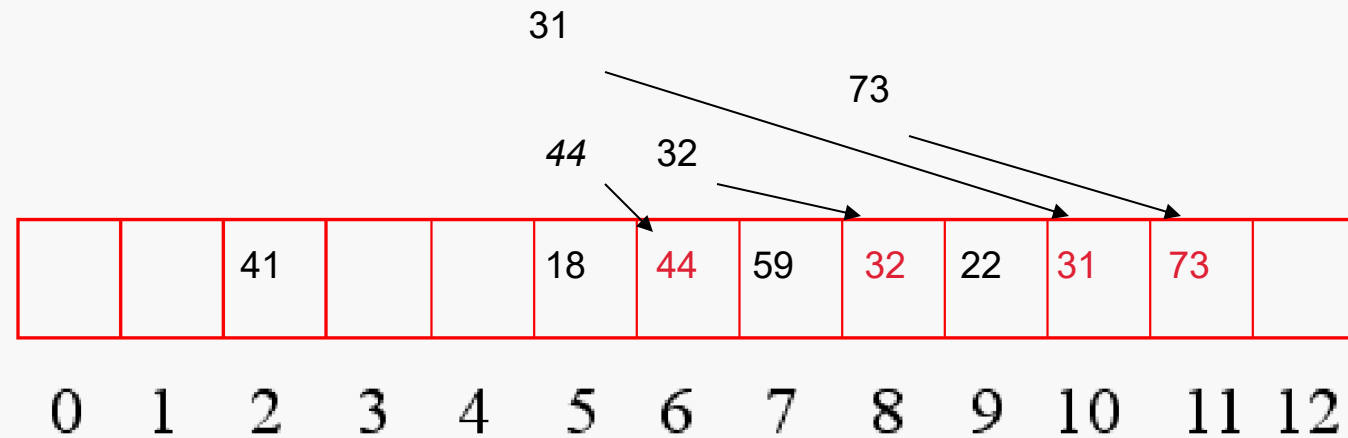
	Initial ly	Insert 4371	Insert 1323	Insert 6173	Insert 4199	Insert 4344	Insert 9699	Insert 1889
0							9699	9699
1		4371	4371	4371	4371	4371	4371	4371
2								1889
3			1323	1323	1323	1323	1323	1323
4				6173	6173	6173	6173	6173
5						4344	4344	4344
6								
7								
8								
9					4199	4199	4199	4199

Linear Probing without replacement Example

$$h(k) = k \bmod 13$$

Insert keys:

18 41 22 44 59 32 31 73



Linear Probing without replacement Example(cont.)

		41			18	44	59	32	22	31	73	
0	1	2	3	4	5	6	7	8	9	10	11	12



Linear probing (chaining) with and without replacement

12,01,04,03,07,08,10,02,05,14

Assume buckets from 0 to 9 and each bucket has 1 slot.

[illegible]

[illegible]



Functions for Linear Probing without replacement insert function

Algorithm int insert_linear_prob(int hashtable[],int key)

```
{
    loc=key % MAX;
    If(hashTable[loc]==-1)
    {
        hashtable[loc]=key;
        return(loc);
    }
    else {i=(loc+1)%MAX;
while(i!=loc)
{
    if(hashTable[i]==-1)
    {
        hashtable[i]=key;
        return(i);    }

    i=(i+1)%MAX;
}
```

```
if(i==loc) print("Hash is full")
}
```

Functions for Linear Probing without replacement- search function

```
Search_linear_probe(hashtable,key)
{
    int i,j;
    loc=key % MAX;
    if(hashtable[loc]==key) print "key found"
    else{ i=(loc+1)%MAX;
    while(i!=loc){
        if(hashtable[i]==key) {print "key found" break;}
        i=(i+1)%MAX;}
    }
    if(i==loc) print("key not found")
}
```

Functions for Linear Probing with replacement insert function

Algorithm `int insert_linear_prob(int hashtable[],int key)`

```
{
    loc=key % MAX;
    If(hashTable[loc]==-1)
    {
        hashtable[loc]=key;
        return(loc);
    }
    else{
temp=key
if(loc!=(hashtable[loc]%MAX){
temp=hashtable[loc];
Hashtable[loc]=key;}
i=(loc+1)%MAX;
```

```
while(i!=loc)
{
if(hashTable[i]==-1)
{
    hashtable[i]=temp;
    return(i);}

i=(i+1)%MAX;
}

if(i==loc) print("Hash is full")
}
```

Functions for Linear Probing with replacement- search function

```
Search_linear_probe(hashtable,key)
{
    int i,loc;
    loc=key % MAX;
    if(hashtable[loc]==key) print "key found"
    else{ i=(loc+1)%MAX;
    while(i!=loc){
        if(hashtable[i]==key) {print "key found" break;}
        i=(i+1)%MAX;}
    }
    if(i==loc) print("key not found")
}
```

Clustering

One problem with the above technique is the tendency to form “clusters”

A **cluster** is a group of items not containing any open slots

The bigger a cluster gets, the more likely it is that new values will hash into the cluster, and make it ever bigger

Clusters cause efficiency to degrade

Quadratic Probing

- ❖ In quadratic probing, we add offset by amount square of collision probe number
- ❖ In quadratic probing, the empty location is searched using the following formula
- ❖ $(\text{Hash}(\text{Key}) + i^2) \text{ MOD } \text{Max}$ where i lies between 1 to $(\text{Max} - 1)/2$
- ❖ Quadratic probing probes the array at location $\text{Hash}(\text{key})+1$, $\text{Hash}(\text{key})+4$, _____, $\text{Hash}(\text{key})+9$ etc.
- ❖ Quadratic probing works much better than linear probing, but to make full use of hash table, there are constraints on the values of i and Max so that the address lies in table boundaries

Open addressing using quadratic probing

- ❖ Let us insert these keys using quadratic probing
- ❖ For 6173, the hashed address $6173 \% 10$ gives 3 and it is not empty, hence using quadratic probing we get the address as follows:
Apply probing $\text{Hash}(6173) = (h(6173) + 1^2) \% 10 = 4$ and as it is empty, the key 6173 is stored there
- ❖ Now while inserting 4344, the location 4 is not empty and hence quadratic probing generates the address as $(4344 + 1^2) \% 10 = 5$ and as is empty 4344 is stored
- ❖ For key 9699, the address is $(h(9699) + 1^2) \% 10 = 0$ and is empty so store.
- ❖ While inserting 1889, the address $(1889 + 1^2) \% 10 = 0$ is not empty so probe again
- ❖ The address $(h(1889) + 2^2) \% 10 = 3$ is not empty so probe again.
- ❖ The address $(h(1889) + 3^2) \% 10 = 8$ is empty so store 1889 at location 8

Open addressing using quadratic probing

	Initially	Insert	Insert	Insert	Insert	Insert	Insert	Insert
		4371	1323	6173	4199	4344	9699	1889
0							9699	9699
1		4371	4371	4371	4371	4371	4371	4371
2								
3			1323	1323	1323	1323	1323	1323
4				6173	6173	6173	6173	6173
5						4344	4344	4344
6								
7								
8								1889
9					4199	4199	4199	4199

Quadratic Probing

If cell j is already occupied then location $j+1, j+4, j+9$ are checked.

This reduces **primary clustering** of data.

Main Idea: Spread out the search for an empty slot – **Increment by i^2 instead of i**

$$h_i(X) = (\text{Hash}(X) + i^2) \% \text{TableSize}$$

$$h_0(X) = \text{Hash}(X) \% \text{TableSize}$$

$$h_1(X) = \text{Hash}(X) + 1 \% \text{TableSize}$$

$$h_2(X) = \text{Hash}(X) + 4 \% \text{TableSize}$$

$$h_3(X) = \text{Hash}(X) + 9 \% \text{TableSize}$$

Let the sequence of **keys** = 9 ,19 ,29 ,39 ,49 ,59,71 These keys are to be inserted into the hash table.

The hash function for indexing, $=K \bmod 10$, where k = key value.

index	keys
0	19
1	71
2	
3	29
4	
5	49
6	
7	
8	39
9	9

59

Quadratic Probing Operation

```
Algorithm Insert(int hashtable[],int key)
{
    start=key % MAX;
    j=start
    for(i=0;i<(MAX-1)/2,i++)
    {
        If(hashtable[j]==-1)add key
        j=(start+i*i)%MAX;
    }
}
```

Double Hashing

- ❖ Double hashing uses two hash functions, one for accessing the home address of a Key and the other for resolving the conflict. The sequence for probing is generated in the following sequence:
-

- ❖ The formula to be used for double hashing is

$$H1(key) = key \% \text{tablesize}$$

$$H2(key) = R - (key \% R)$$

where R is the prime no. smaller than the size of the table.

The formula is : $H1(key) + i * H2(key)$ $i = 0, 1, 2, \dots$

- ❖ The resultant address is divided by modulo Max
- ❖ Every key in K that shares a common factor with the number of buckets R will be hashed to a bucket that is a multiple of this factor.
- ❖ By choosing R to be a number that has very few factors: a prime number

Open addressing using double hash function

- ❖ While inserting 6173, the address is $\text{Hash1}(6173) = 6173 \% 10 = 3$ and 3 is not empty
- ❖ Let us use double hashing. Hence the address is as follows:
- ❖ $\text{Hash}(6173) = [\text{Hash1}(6173) + \text{Hash2}(6173)] \% 10 = 3 + (R - 6173 \% R)$ (let R be 7) $= 3 + (7 - 6) = 4$ Since 4 is empty, we store 6173 at location 4
- ❖ Now let us store 4344. The address $4344 \% 10 = 4$ and as location 4 is not empty, we use double hashing and we get $\text{Hash}(4344) = 7$
- ❖ Now for 9699 double hashing generates address 2 and as it is empty, we store it there.
- ❖ For key 1889, double hashing generates address 0 and as it is empty, we store 1889 at location 0

Open addressing using double hash

	Initially	Insert 4371	Insert 1323	Insert 6173	Insert 4199	Insert 4344	Insert 9699	Insert 1889
0								1889
1		4371	4371	4371	4371	4371	4371	4371
2							9699	9699
3			1323	1323	1323	1323	1323	1323
4				6173	6173	6173	6173	6173
5								
6								
7						4344	4344	4344
8								
9					4199	4199	4199	4199

Double Hashing Example

insert(14)

$$14\%7 = 0$$

insert(8)

$$8\%7 = 1$$

insert(21)

$$21\%7 = 0$$

$$5 - (21\%5) = 4$$

insert(2)

$$2\%7 = 2$$

insert(7)

$$7\%7 = 0$$

$$5 - (7\%5) = 3$$

0	14
1	
2	
3	
4	
5	
6	

1

0	14
1	8
2	
3	
4	
5	
6	

1

0	14
1	8
2	
3	
4	21
5	
6	

2

0	14
1	8
2	2
3	
4	21
5	
6	

1

0	14
1	8
2	2
3	7
4	21
5	
6	

??

probes:

89,18,49,58,69,23
Prime number $R = 7$
Table size $M = 10$

0	69
1	
2	
3	58
4	
5	
6	49
7	
8	18
9	89

Rehashing

- ❖ If table gets full, insertion using open addressing with quadratic probing might fail or it might take too much time.
- ❖ Rehashing tells us what to do when the hash table gets full.
- ❖ Instead of waiting for the hash table to get completely full, it is more efficient to rehash when the table is about 70% or 80 % full.
- ❖ To find the solution for this is to build another table that is about twice as big and scan down the entire original hash table, compute the new hash value for each record, and Insert them in a new table.
- ❖ In rehashing we can also choose series of hash functions while increasing the size of TABLE.

Rehashing (contd...)

For example, if table is of size 7 and hash function is $\text{key} \% 7$ then,

Insert 7,15,13,74,73

0	7
1	15
2	
3	73
4	74
5	
6	13

Rehashing (contd...)

0	
1	
2	
3	
4	
5	73
6	74
7	7

8	
9	
10	
11	
12	
13	13
14	
15	15
16	

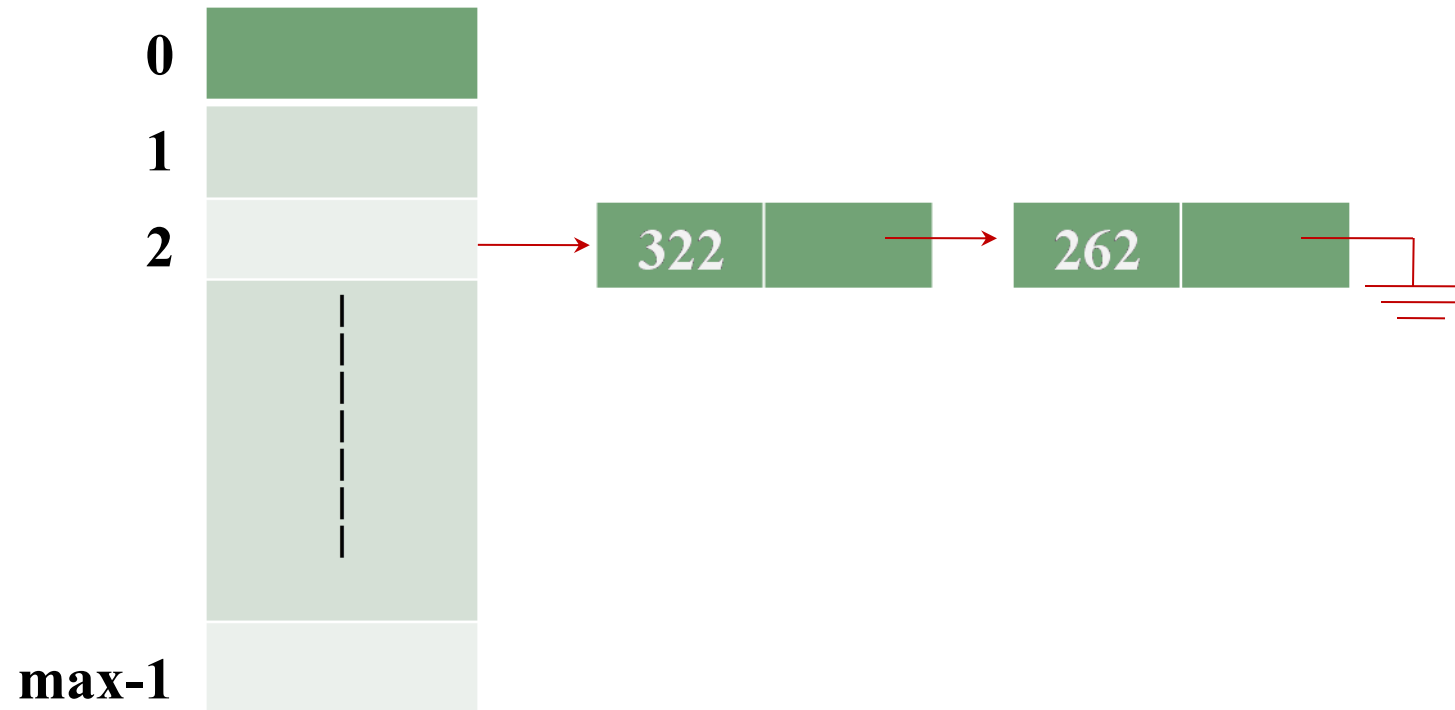
As the table is more than 70 % full, new table is created and the values are inserted in the new table

The size of the new table is 17, that is next prime of double of 7 that is 14
Rehashing is very expensive, as its running time is $O(N)$.

Chaining

-
- ❖ This technique used to handle synonym is chaining that chains together all the records that hash to the same address. Instead of relocating synonyms, a linked list of synonyms is created whose head is home address of synonyms
 - ❖ However, we need to handle pointers to form a chain of synonyms
 - ❖ The extra memory is needed for storing pointers

Chaining



Comparison : Rehashing and chaining

Chaining

- ❖ Unlimited number of synonyms can be handled in chaining
- ❖ Additional cost to be paid is overhead of multiple linked lists
- ❖ Sequential search through chain takes more time

Rehashing

- ❖ Limited but still a good number of synonyms are taken care of
- ❖ The table size is doubled but no additional field of link is to be maintained
- ❖ Searching is faster when compared to chaining

Hash Table Overflow

-
- ❖ An overflow is said to occur when a new identifier is mapped or hashed into a full bucket
 - ❖ When the bucket size is one, collision and overflow occur simultaneously

Open Addressing for Overflow Handling

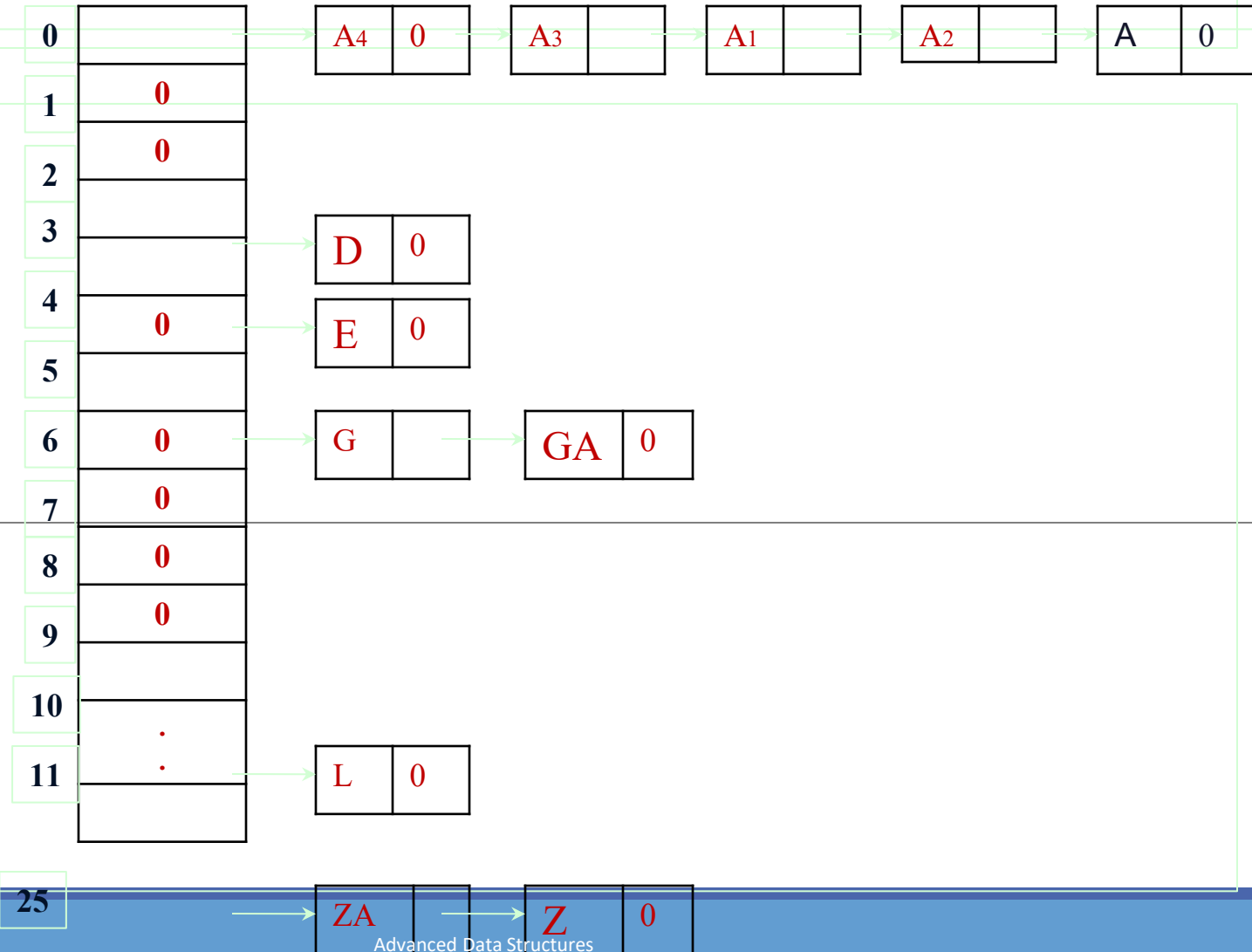
-
- ❖ When a new identifier is hashed into a full bucket, we need to find another bucket for this identifier
 - ❖ The simplest solution is to find the closest unfilled bucket.
 - ❖ This is called as linear probing or linear open addressing

Overflow Handling by Chaining

- ❖ Since the sizes of these lists are not known in advance, the best way to maintain them is as linked chains
- ❖ In each slot, additional space is required for a link
- ❖ Each chain has a head node.
- ❖ The head node, however, usually is much smaller than the other nodes, since it has to retain only a link
- ❖ As the list is accessed at random, the head nodes should be sequential

0	1	2	3	4	5	6	7	8	9	10	11		25
A	A2	A1	D	A3	A4	G _A	G	Z _A	E		L		Z

Chaining



Separate Chaining-Used with Open Hashing

Data Structure:

```
#define MAX 1
```

```
class node
```

```
{
```

```
    int data;
```

```
    node * next;
```

```
};
```

```
node *hashtable[MAX];
```

```
Algorithm Initialize(node *hashtable[])
```

```
{
```

```
    for (i=0;i<MAX;i++)
```

```
{
```

```
    allocate memory for head[i];
```

```
    Hashtable[i]->data=i
```

```
    Hashtable[i]->next=NULL;
```

```
}
```

```
}
```

Separate Chaining-Used with Open Hashing

Operations on Hash Table:

Algorithm Insert(node *hashtable[],int x)

```
{  
    int loc;  
    node *p,*q;  
    loc=x % MAX;  
    q=new node;  
    q->data=x;  
    q->next=NULL;  
    If(hashtable[loc]->next==NULL)  
        Hashtable[loc]->next=q;  
    else  
    {  
        for(p=hashtable[loc];p->next!=NULL;p=p->next);  
        p->next=q;  
    }  
}
```

Separate Chaining-Used with Open Hashing

Operations on Hash Table:

Algorithm `node *find(node *hashtable[],int x)`

{

int loc,

node *p;

loc=x % MAX;

p=hashtable[loc];

While(p!=NULL && x!=p->data)

 p=p->next;

return (p);

}

FAQ

What is hashing?

Why there is need of hashing?

What is the time complexity of hashing?

What is collision?

Define hash function, hash key

What are the collision resolution techniques?

MCQs

A hash table of length 10 uses open addressing with hash function $h(k) = k \bmod 10$, and linear probing. After inserting 6 values into an empty hash table, the table is as shown below. Which one of the following choices gives a possible order in which the key values could have been inserted in the table?

0	
1	
2	42
3	23
4	34
5	52
6	46
7	33
8	
9	

- A) 46, 42, 34, 52, 23, 33
- B) 34, 42, 23, 52, 33, 46
- C) 46, 34, 42, 23, 52, 33
- D) 42, 46, 33, 23, 34, 52

How many different insertion sequences of the key values using the same hash function and linear probing will result in the hash table shown above?

- A.10
- B.20
- C.30
- D.40

MCQs

The keys 12, 18, 13, 2, 3, 23, 5 and 15 are inserted into an initially empty hash table of length 10 using open addressing with hash function $h(k) = k \bmod 10$ and linear probing. What is the resultant hash table?

0	
1	
2	2
3	23
4	
5	15
6	
7	
8	18
9	

(A)

0	
1	
2	12
3	13
4	
5	5
6	
7	
8	18
9	

(B)

0	
1	
2	12
3	13
4	2
5	3
6	23
7	5
8	18
9	15

(C)

0	
1	
2	12, 2
3	13, 3, 23
4	
5	5, 15
6	
7	
8	18
9	

(D)

MCQs

Consider a hash table of size seven, with starting index zero, and a hash function $(3x + 4) \bmod 7$. Assuming the hash table is initially empty, which of the following is the contents of the table when the sequence 1, 3, 8, 10 is inserted into the table using closed hashing? Note that ‘_’ denotes an empty location in the table.

A 8, _, _, _, _, _, 10

B 1, 8, 10, _, _, _, 3

C 1, _, _, _, _, _, 3

D 1, 10, 8, _, _, _, 3

Given the following input (4322, 1334, 1471, 9679, 1989, 6171, 6173, 4199) and the hash function $x \bmod 10$, which of the following statements are true? i. 9679, 1989, 4199 hash to the same value ii. 1471, 6171 has to the same value iii. All elements hash to the same value iv. Each element hashes to a different value (GATE CS 2004)

A i only

B ii only

C i and ii only

D iii or iv

Practice Assignments

1. Write a pseudo code to implement quadratic probing
2. Store roll numbers of the students in a database and implement double hashing onto it.
3. Store roll numbers of the students in a database and implement linear probing with and without replacement onto it.
4. Implement chaining method by identifying one suitable application.

Takeaway

- Hashing is one of efficient searching technique.
- Hashing's best case time complexity is $O(1)$
- Hashing collision is resolved by using different collision resolution techniques.

References

1. Horowitz, Sahani, Dinesh Mehta, “Fundamentals of Data Structures in C++”, Galgotia Publisher, ISBN: 8175152788, 9788175152786.
2. Peter Brass, “Advanced Data Structures”, Cambridge University Press, ISBN: 978-1-107-43982