Introduction to Distributed Databases

- Basic Concepts and Notions
- **II.** <u>Distributed DB Design</u>
- Distributed Query Processing

II. Distributed Database Design

II. 1. Design Approach

■ Top-Down Design Process

- Design a Global Schema
- Distribute the Global Schema objects to obtain Local Schemas
- → Definition of Homogeneous Distributed Database : "DDB in which each local DB is managed by the same DBMS"; Gardarin et al. 1989.
- Advantages
- The growing ease of incremental business: By the facility of adding a new site
- Increases performance: Exploit the capacities of parallel treatments/processing

Bottom-Up Design Process

- · Consists of integrating existing LS in one or more Global Schemas
- Allows the integration of existing local Database into a Federated Database
- Needs a semantic reconciliation of schemas (data type,
- Ensures the continuity of services
 - Fragmentation & allocation are imposed
- → Definition of Heterogeneous Distributed DB: "Distr. DB in which the local databases are managed by different DBMS"; Gardarin et al. 1989

II.2. Fragmentation Strategies

Fragmentation (Definition):

Sub-relation obtained by selection of tuples and attributs from a global relation. Each fragment has a subset of the tuples of the relation.

- → Two possible ways to divide a relation : *Horizontally* (selection of tuples), *Vertically* (selection of attributs). In addition, "there is possibility of nesting a fragments in hybrid fashion"
- Fragmentation Rules: the main objective is to preserve the semantic consistency of the DDB
 - 1.: Decomposition without loss of information (R1, R2, ... Rn)
 - 2.: No duplication: restriction predicates define disjoint sets of tuples
 - 3.: Reconstruction rule (global relation): R = Function (R1, R2, ...Rn)

II.2. Fragmentation Strategies

Primary Horizontal Fragmentation PHF

- PHF partitions a relation along its tuples. It is defined by a selection operation on the the relation : Ri= $\sigma_{Fi}(R)$, 1 <= i <= k where Fi is the selection formula used to obtain fragment Ri (also called fragmentation predicate)
- Wines (w_Id, vineyard, vintage, Area) ∈ GS
 - Wines1= $\sigma_{Area=Bourgogne}$ (Wines)
 - − Wines2= σ_{Area≠Bourgogne}(Wines)
- 3: Reconstruction rule (global relation) : Wines = Wines1 \cup Wines2

Benefits

 Favors the selection Queries by limiting the number of accessed fragments

II.2. Fragmentation Strategies: Derived Horizontal Fragment

- Objective: The PHF can also be useful for processing of join queries. We need to modify the HF definition By applying Semi-join predicates
- DHF Def.: Partitioning function of a relation into sub-sets of tuples, each being defined by a semi-join operation of the relation with a fragment of another relation.
- **Example**: Wines (w_Id, vineyard, vintage, area), Drink (w_Id, drinker_Id, date, quantity)

 - Reconstruction Rule : Drink = Drink1 ∪ Drink2
- Exercice: Select * from W, D where W.w_Id= D.w_Id and area= «Bourgogne »
- DHF Benefits?
 - Improves joins between pairs of fragments Wines 1 ∞ Drink 1 and Wines 2 ∞ Drink 2

II.2. Derived Horizontal Fragmentation (2)

Wines 1	w_ld	vineyard	vintage	aera
	1	Chablis	2010	Bourgogne
	2	Pommard	2009	Bourgogne
Wines2	w_ld	wineyard	vintage	aera
Willesz	w_iu	willeyalu	viiitage	aera
	3	Juliennas	2010	Beaujolais
	4	Chinon	2011	Loire
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Drink1	drinker_ld	wineld	date	quantity
	1	1	14/02/2012	1
	1	2	1/1/2013	2
Drink2	drinker_ld	wineld	date	quantity
	1	3	14/02/2013	3
	3	4	1/1/2013	2

II.2. Fragmentation Strategies: Vertical fragmentation

- VF Defintion: «The VF function distributes a relation on projection attributes: fragments are obtained by projections applyed to a logical relation ».
- The reconstruction operator for VF: « is the join of the fragments on the common attributes »
- **Example**: Wines (w_Id, vineyard, vintage, area, degree, price) ∈ GS
 - -Wines1= π_{w_Id} , vineyard, area, vintage(Wines)
 - Wines2= $\pi_{\text{W_Id, degree, price}}$ (Wines)
 - Reconstruction rule: Wines = Wines1 ∞ Wines2
- VF Benefits:
 - Favors project queries by limiting the number of accessed fragments.

II.2. Fragmentation Strategies: Hybrid fragmentation

- Wines (w_Id, vineyard, vintage, aera, degree, price) ∈ GS
 - Wines1= $\pi_{w_{Id}, \text{ degree, price}}$ (Wines)
 - Wines2 = $\sigma_{\text{area}=\text{bourgogne}}(\pi_{\text{w_Id, wineyard, vintage, aera}}(\text{Wines}))$
 - − Wines3 = $\sigma_{\text{aera} \neq \text{bourgogne}} (\pi_{\text{w_Id, wineyard, vintage, aera}} (\text{Wines}))$
 - Reconstruction rule: Wines = Wines1 ∞ (Wines2 ∪ Wines3)
 - Benefits?
 - Benefits of horizontal & vertical fragmentations!

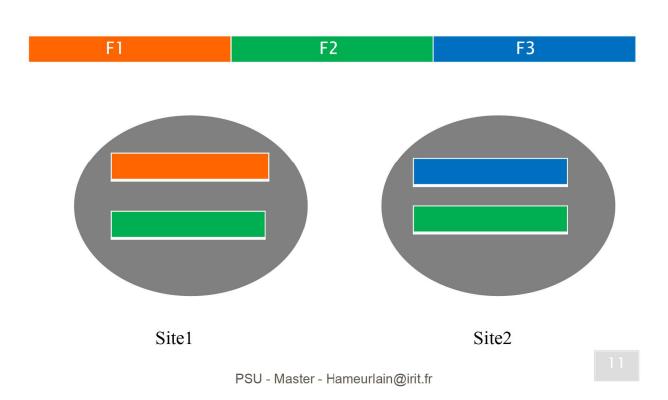
II.3. Choice of a fragmentation Strategy

- Complex issue because it depends on access requirements of queries
 - Classification of Queries:
 - Repetitive Queries
 - Ad-hoc Queries
 - Query Types: < Simple, Medium, Complex>
 - Query Natures: <Selection, Projection, Join>

II.4. Allocation of fragments (1/2)

- ■Logically group data in accordance with the access objectives of applications ⇒ performance
 - Ex: Wines 1 & Drink 1 will be grouped on the same site
- Data Placement
 - to allocate fragments on sites

II.4. Fragment allocation (2/2)



Exercice (will be done in the next tutorial)

- We consider the following Global Relational Schema GRS:

 Cars (regNum, Brand(=marque), type, power, renting_agency, price)

 Customers (customerId, name, firstname, address, phone)

 Emp (empId, name, firstname, address, renting_agency, salary, commi.)

 Rent (regNum, customerId, date, discount).
- Hypothesis: We consider 3 agencies : <Pau, Agen, Toulouse>
- Constraints:
 - Each car is managed by a single agency.
 - Each agency manages its sales performance and issues invoices to its customers.
 - Each employee is assigned to an agency, his salary and his commissions managed only by Toulouse.

Questions :

- Fragments?
- Fragments' allocation?
- Reconstruction rules?