

Too secure system

Bob worries about security too much and decided to change the original Pedersen commitment scheme adding there several steps. Find the vulnerability in Bob's system!

Initialization

- 1) Choose two big primes p and q such that $q|(p-1)$ and choose $g \in Z_p^*$ of order q .
- 2) Choose $x \in \mathbb{Z}_q^*$. x is our secret.
- 3) Compute $G = g^x \bmod p$
- 4) Transform G into G' : 1024-bit binary number with big endian.
- 5) Using hash-function SHA-512, compute $a = \text{SHA512}(G')$
- 6) Find a' : transform binary a into an integer number.
- 7) Find $\hat{a} = (a'^{a'}) \bmod \phi(p)$, where $\phi(p)$ is Euler function.
- 8) Finally, calculate $h = g^{\hat{a}} \bmod p$.

Parameters p and g are well-known and open for everyone.

Commitment

To make a commitment we choose a random r from $\{2, \dots, p-1\}$ and calculate $c = G \cdot h^r$. The commitment would be c .

Proof

After the decision is made by sender and verifier about the verification of the commitment, sender opens x and r values, and the verifier checks the equality $c = G \cdot h^r$.

Additional information

To get integer value from the string $S = S_1 S_2 \dots S_n$ of n ASCII characters, use the following formula:

$$\text{func}(S) = \sum_{i=0}^{n-1} \text{ASCII}(S_{i+1}) \cdot 2^{8 \cdot i}$$

where $\text{ASCII}(\dots)$ is an ASCII value of the character.

Task

Break the system with the following parameters. As a result, you should get the value r_2 which you disclose to the verifier with the string M_2 , and the verifier confirms the commitment as with (M_1, r_1) as with (M_2, r_2) .

The parameters are:

- $p = 12039102490128509125925019010000012423515617235219127649182470182570195018265927223$
- $g = 10729072579307052184848302322451332192456229619044181105063011741516558110216720725$
- $r_1 = 31245182471$
- $M_1 = \text{'Hi! I am Vadim Davydov from ITMO University'}$
- $M_2 = \text{'Transfer the points for easy task to this team'}$