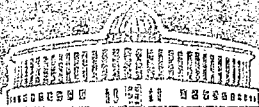


СООБЩЕНИЯ
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V.P.Shirikov, D.C.Marinescu

**THE PARAMETER PASSING
MECHANISM IMPLEMENTATION
IN DIFFERENT FORTRAN COMPILERS**

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MECHANISM IMPLEMENTATION
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* On leave of absence from the Institute
for Atomic Physics, Bucharest, Romania

Шриков В.П., Маринеску Д.К.

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Применение и реализация механизма передачи параметров
в различных трансляторах с языка ФОРТРАН

В работе содержится характеристика типов передачи параметров между различными программными модулями и способа реализации этих типов в различных вариантах трансляторов (в основном, трансляторов с языка ФОРТРАН). Изложение ведется применительно к используемым в настоящее время системам математического обеспечения ЭВМ серии ЕС, ЭВМ фирмы СДС и ИВМ, ЭВМ БЭСМ-6.

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Сообщение Объединенного института ядерных исследований. Дубна 1977

Shirikov V.P., Marinescu D.C.

E10 - 10543

The Parameter Passing Mechanism Implementation
in Different FORTRAN Compilers

This report contains some considerations, remarks and comparison for the types and methods of parameter transmissions between different programming modules and also for the methods of implementation for these transmissions in different compilers versions (first of all, FORTRAN compilers). The compilers under consideration are the compilers for the software systems which are used now with EC, CDC, IEM and BESM-6 computers.

The investigation has been performed at the Laboratory of Computing Techniques and Automation, JINR.

Communication of the Joint Institute for Nuclear Research. Dubna 1977

It is very important for any programmer to recognize in which way the parameters are communicated from one module to another for every implementation of the programming languages available.

For most programming languages this task is not so different since the designers of the compilers have no other choice but to implement that particular type of call as stated by people who have designed the language. For example:

- in ALGOL the default argument transmission mechanism is by name; it can be changed to transmission by value by specifying the parameter as such in the procedure header.

- In PL/1 the default argument transmission mechanism is by reference; it changes to transmission by value when the argument is enclosed in an extra set of parenthesis, in the CALL statement or in the function reference. Also transmission by value arises when the argument is a constant or involves operators.

- In COBOL the arguments are transmitted by reference.

As far as FORTRAN is concerned several approaches are used:

- the strategy announced in CDC manuals ^{/1,2/} is said to be the call by name with the call by value of expressions appearing as actual parameters; also the call by value is used by internal functions. This approach applies to both FTN and RUN compilers available on CYBER machines.

- the IBM approach ^{/3/} used in FORTRAN F , G , H compilers makes use of the call by value result, parameter passing mechanism, as default, with the user option to request the call by reference, by enclosing the dummy arguments between slashes.

The practical implications of such different points of view are: sometimes the same program compiled by different FORTRAN com-

compilers (and even at different optimization levels of the same compiler) gives different results when executed.

This might look queer for someone not aware of the fact that for FORTRAN the convention concerning the parameter passing mechanism is lax and somehow it is left to the compiler designer to decide about the parameter passing mechanism.

To illustrate this idea we present in table 1 the results obtained when the main program and the subroutine from example # 1 were translated by several compilers available around.

X = 1.0	SUBROUTINE SUB (A,B,C,D)
Y = 2.0	B = A + A
Z = 7.0	D = A + C
CALL SUB (X,X,X+Y,Z)	RETURN
PRINT 1,X,Y,Z	END
1 FORMAT(1H, 3F15.1)	
STOP	
END	

EXAMPLE 1.

COMPILER	TYPE OF CALL, ACCORDING TO THE MANUALS	VALUES			OBSERVATIONS
		X	Y	Z	
FTN OPT=0	Call by name	2	2	5	User request OPT=0
FTN OPT=1	Call by name	2	2	4	User request OPT=1
FFORTRAN	Call by value/ result	2	2	4	Used the default type of call.
FFORTRAN	Call by reference	2	2	5	The user requested call by reference.
RUN	Call by name	2	2	4	
BESM-6	Call by reference	2	2	5	

TABLE 1.

Let us briefly examine the four basic types of parameter communication mechanisms and some details of their practical implementation.

- A./ Call by reference

It is eventually the easiest to implement and practically consists of passing the addresses of the actual arguments to the called program. As far as our example 1 is concerned, the expression:

$$W = X + Y$$

is first evaluated and a stack consisting of the addresses of X, X, W and Z is constructed (see appendix 1.2 and figure #1)

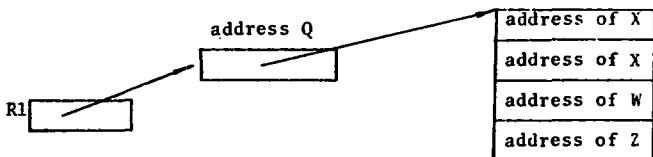


FIGURE #1. The call by reference implementation in FORTRAN F compiler.

The calling program loads in general purpose register R1 the address of a memory location Q where a pointer to the stack is to be found. It is the responsibility of the linkage editor to fill Q with the stack address.

When entered the subroutine copies the stack and throughout the computation operates with the locations in the main storage for X, X, W and Z, as defined by the calling program.

In this situation, side effects can occur; one of the most unpleasant ones occurs when a constant is passed as an actual argument and the called module attempts to modify the dummy argument to which the constant is associated. In case of example #2 presented below, most FORTRAN compilers would lead to the result J=11.

N = 8	SUBROUTINE S (I,J,K)
.....	K = K + J
M = 6	I = I + 2
J = 8	RETURN
CALL S (1,M,N)	END
J = J + 1	
PRINT 1,J	
1 FORMAT (1H , I5)	
STOP	
END	

EXAMPLE #2.

Sometimes to protect against unpleasant surprises, FORTRAN manuals ^{/3/} warn that "input parameters must not be changed inside the called module" (must not appear as the left-hand side of an assignment statement).

Another type of side effect is due to the fact that the call by reference is associated with the call by value of expressions appearing as actual parameters, as in the case of example #1; there, the expression is first evaluated using the old value for X. This type of side effect does not occur if the rule of not modifying the input parameters is observed and if every variable appearing in the expression is considered as an input parameter.

- B./ Call by value

Instead of passing the parameter addresses , in this case their values are communicated in a stack.

If the machine has enough registers the values of the parameters can be passed into the registers, thus saving time.

It should be noticed that there is no feedback from the called module to the calling one so that, though this strategy is most suitable for input parameter passing, it leaves no hope for output parameter communication. This is precisely the reason why this type of call has a limited use. As stated before the FTN compiler uses this type of call for internal functions since only input parameters are to be passed, their number is limited and anyway

the result is returned as a variable with the same name as the function name.

- C./ Call by value/result

As in the previous case the values of the actual parameters are passed but also a pointer to a stack of parameter addresses is communicated. As a result, the called module operates with the values passed, but before returning to the calling module the values of actual arguments are updated.

If we examine appendix 1.2 we see that when the subroutine SUB is entered, the general purpose register R1 contains the address Q, where a pointer to the stack of addresses is inserted by the linkage editor; a copy of parameter values is made (the sequence of instructions starting at label A20). Then the computation suggested by statements 2 and 3 of the subroutine is performed (the sequence of instructions starting at label A52).

Before returning, R1 is loaded with the address of a memory location, where a pointer to the stack of parameter addresses is to be found (the sequence of two instructions starting at label A36), then all four parameters X,X,W,Z get their values updated (the four pairs of instructions, load and move constant).

It results that in addition to the side effects previously encountered, in connection with the call by reference, here occurs an additional one; the computations performed by the called module are using the initial values of the input parameters.

- D./ Call by name

The call by name seems to be the most natural to communicate parameters since it is ment to implement the textural substitution and thus no side effects can possibly occur.

In case of our example #1, A is substituted by X, B by X, C by X+Y, D by Z and the executable statements of the subroutine are:

$X = X+X ; /X = 1+1 = 2/$

$Z = X+(X+Y); /Z = 2+2+2 = 6/.$

Obviously since the names of variables have a significance only within a module some sophisticated mechanism must be available so that when a reference is made to X in the called module, its address should be supplied. It means that at run-time, a routine must be entered every time a parameter is referred to, to supply

its address. This leads to considerable inefficiency and it is rather difficult to construct.

As far as the FTN compiler is concerned the type of call is by no means a call by name. Appendix 2.1 proves this statement when optimization level one is requested. Examining the COMPASS expansion we see that:

- registers X5,X4,X3,X2 are loaded with the addresses of X,X,W and Z using the pointer to the stack, provided in A0.
- the initial value of X (value 1), is loaded into X1.
- the new value for X (value 2) is computed in X7 and the memory location reserved for X is updated.
- when Z is computed, the old value of X, existing in X1, is used so that the result $4=1+3$ is obtained.

In appendix 2.2 we see that when optimization level zero is requested, the same procedure is followed, but for X it is used the value in memory address and not the one in register; so that the result $5=2+3$ is obtained.

A summary of values to be expected for the program in example #1 for the four possible types of parameter passing mechanisms is presented in table 2.

TYPE OF CALL	VALUES/EXAMPLE 1		
	X	Y	Z
Call by reference	2.0	2.0	5.0
Call by value	1.0	2.0	7.0
Call by value/ result	2.0	2.0	4.0
Call by name	2.0	2.0	6.0

TABLE 2

CONCLUDING REMARKS

- 1. The call by name is by far the best method of parameter passing since it gives no side effects; it is the most diffi-

cult to implement and eventually leads to inefficiency as far as execution time is concerned. It has not been encountered in any of the compilers under scrutiny.

- 2. As far as the unprejudiced programmer is concerned, he must not take for granted whatever manuals state, but he must try to understand what lies behind each type of call and he must be able to test the compiler he is using.

References

1. FORTRAN Extended Reference Manual. CDC publication #60329100.
2. FORTRAN Reference Manual. CDC publication #60174900.
3. FORTRAN Reference Manual. IBM publication.
4. G.J.Myers. IBM Systems Journal, vol. 15, # 3, 1976.

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on March 30, 1977.

APPENDIX 1. Tests made with the FORTRAN F compiler.

FORTRAN main program and its ASSEMBLER expansion.

The subroutine SUB in FORTRAN and its ASSEMBLER expansion

1. When call by value result is used (APPENDIX 1.1)
2. When call by reference is used (APPENDIX 1.2)

```
DOS FORTRAN IV 36CN-FO-479 3 8          MAINPGM          DATE 2'
C001          X=1,J
C002          Y=2,J
C003          Z=7,J
C004          CALL SUB(X,X,X+Y,Z)
C005          WRITE(3,1) X,Y,Z
C006          1  FORMAT(IH,3F9.1)
C007          STOP
C008          END
```

```
DOS FORTRAN IV 36CN-FO-479 3 8          MAINPGM          DATE 27

LOCATION    STA NUM    LABEL    OP    OPERAND
C00000          DC    15,0
C00002          LM    2,3,34(15)
C00005          L    15,30(0,15)
C0000A          LA    15,2(0,15)
C0000E          ST    15,4(0,13)
C00012          BCR  15,2
C00014          DC    00000000
C00018          DC    0704C1C9
C0001C          DC    0507C7D4
C00020          DC    01000000
C00024          DC    03000000
C00028          DC    08000000
C00108          L    15,4(0,13)
C0010C          A57  L    14,12(0,13)
C00110          LM    2,12,23(13)
C00114          MVI  12(13),255
C00118          BCR  15,14
C0011A          A36  L    15,112(0,13)
C0011E          LR    12,13
C00120          LR    12,4
C00122          BAL  14,64(0,15)
C00126          LR    13,12
C00128          1  LE    0,204(0,13)
C0012C          STE  0,96(0,13)
C00130          2  LE    0,208(0,13)
C00134          STE  0,100(0,13)
C00138          3  LE    0,212(0,13)
C0013C          STE  0,104(0,13)
C00140          4  LE    0,96(0,13)
C00144          AE    0,100(0,13)
C00148          STE  0,216(0,13)
C0014C          LA    1,116(0,13)
C00150          L    15,108(0,13)
C00154          BALR 14,15
C00156          BC    0,4(0,0)
C0015A          5  L    15,112(0,13)
C0015E          BCR  0,0
C00160          BAL  14,4(0,15)
C00164          DC    00000003
C00168          DC    00000080
C0016C          L    15,112(0,13)
C00170          BAL  14,8(0,15)
C00174          DC    0470D060
C00178          BAL  14,8(0,15)
C0017C          DC    0470D064
C00180          BAL  14,8(0,15)
C00184          DC    0470D068
C00188          BAL  14,16(0,15)
C0018C          7  L    15,112(0,13)
C00190          BAL  14,52(0,15)
C00194          DC    05404040
C00198          DC    0404
C0019A          END
```

TOTAL MEMORY REQUIREMENTS 00019A BYTES

APPENDIX 1.1. The subroutine SUB in FORTRAN and its ASSEMBLER expansion produced by the FORTRAN F compiler, when standard (call by value/result) parameter passing is requested.

```
0001      SUBROUTINE SUB(A,B,C,D)
0002      B=A+A
0003      D=A+C
0004      RETURN
0005      END
```

DOS FORTRAN IV 360N FC-479 3-R SUB DATE

SYMBOL B	LOCATION 90	SYMBOL A	SCALAR MAP LOCATION 94	SYMBOL D	LOCATI 93
-------------	----------------	-------------	------------------------------	-------------	--------------

DOS FORTRAN IV 360N FC-479 3-R SUB DATE

LOCATION	STA NUM	LABEL	OP	OPERAND
000000			BC	15,12(0,15)
000004			DC	07E2E4C2
000008			DC	40404040
00000C			STM	14,12,12(13)
000010			LM	2,3,40(15)
000014			LR	4,13
000018			L	13,26(0,15)
00001A			STM	3,6(0,4)
00001E			BCR	3,4,0(13)
000022			DC	15,2
000024			DC	00000000
000028			DC	00000000
00002C			DC	00000000
000028		A20	L	7,0(0,1)
00002C			MVC	100(3,13),0(2)
00002E			L	2,4(0,1)
000030			MVC	66(3,13),0(2)
000032			L	2,8(0,1)
000034			MVC	108(3,13),0(2)
000036			L	2,12(0,1)
000038			MVC	104(3,13),0(2)
00003A			BALR	2,0
00003C			L	3,6(0,2)
00003E			BCR	15,3
000040			DC	00000000
000042		A36	L	1,4(0,13)
000044			L	1,24(0,1)
000046			L	2,0(0,1)
000048			MVC	0(3,2),100(13)
00004A			L	2,4(0,1)
00004C			MVC	0(3,2),96(13)
00004E			L	2,8(0,1)
000050			MVC	0(3,2),108(13)
000052			L	2,12(0,1)
000054			MVC	0(3,2),104(13)
000056			L	13,4(0,13)
000058			L	14,12(0,13)
00005A			LM	3,1,28(13)
00005C			MVI	1,1,1,255
00005E			BCR	15,14
000060	2	A52	LE	0,100(0,13)
000062			AE	0,100(0,13)
000064			STE	0,96(0,13)
000066	3		LE	0,100(0,13)
000068			AE	0,108(0,13)
00006A			STE	0,104(0,13)
00006C	4		SR	15,15
00006E			L	1,0(0,13)
000070			BCR	15,1
000072			END	

TOTAL MEMORY REQUIREMENTS 00016E BYTES

The results obtained in this case:

X = 2.0 Y = 2.0 Z = 4.0

Indeed a call by value results is used.

APPENDIX 1.2. The subroutine SUB and its ASSEMBLER expansion when call by reference is requested by the user.

```

0001          SUBROUTINE SUB(/A/,/B/,/C/,/D/)
0002          B=A+A
0003          D=A+C
0004          RETURN
0005          END
  
```

DOS FORTRAN IV 360N FD 479 3 8

SUB

DATE

SYMBOL B	LOCATION 90	SYMBOL A	SCALAR MAP LOCATION 94	SYMBOL D	LOCATION 98
-------------	----------------	-------------	------------------------------	-------------	----------------

DOS FORTRAN IV 360N FD 479 3 8

SUB

DATE

LOCATION	STA NUM	LABEL	OP	OPERAND
000000			BC	15,12(0,15)
000004			DC	07E2E4C2
000008			DC	404C4040
00000C			STM	14,12,12(13)
000010			LM	2,3,40(15)
000014			LR	4,13
000018			L	13,36(0,15)
00001A			ST	13,8(0,4)
00001E			STM	3,4,0(13)
000022			BCR	15,2
000024			DC	00000000
000028			DC	00000000
00002C			DC	00000000
00002F		A20	L	2,0(0,1)
000030			LA	2,0(2,0)
000034			ST	2,160(0,13)
000038			L	2,4(0,1)
00003C			LA	2,0(2,0)
00003E			ST	2,96(0,13)
000040			L	2,8(0,1)
000044			LA	2,0(2,0)
000048			ST	2,108(0,13)
00004C			L	2,12(0,1)
00004E			LA	2,0(2,0)
000050			ST	2,104(0,13)
000054			BALR	2,0
000058			L	3,6(0,2)
00005C			BCR	15,3
000060			DC	00000000
000064		A36	L	1,4(0,13)
000068			L	1,24(0,1)
00006C			L	13,4(0,13)
000070			L	14,12(0,13)
000074			LM	2,12,28(13)
000078			MVI	12(13),255
00007C			BCR	15,14
000080	2	A52	L	10,100(0,13)
000084			LE	0,0(0,10)
000088			AE	0,0(0,10)
00008C			L	11,96(0,13)
000090			STF	0,0(0,11)
000094	3		LE	0,0(0,10)
000098			L	12,108(0,13)
00009C			AE	0,0(0,12)
0000A0			L	10,104(0,13)
0000A4			STE	0,0(0,10)
0000A8	4		SR	15,15
0000AC			L	13,0(0,13)
0000B0			BCR	15,1
0000B4			END	

TOTAL MEMORY REQUIREMENTS 00015E BYTES

The results obtained in this case:

X = 2.0 Y = 2.0 Z = 5.0

They indicate that it is really a call by reference.

APPENDIX 2. Tests with the FORTRAN Extended - FTN compiler.

Appendix 2.1. Main program and subroutine COMPASS expansion when OPT=1.

Appendix 2.2. The same but for the case OPT= 0

AM START 73/74 OPT=1 FIN 4.6+429

PROGRAM START (INPUT,OUTPUT)

X=1.0
Y=2.0
Z=7.0
CALL SUB(X,X+Y,Z)
PRINT 1,X,Y,Z
1 FORMAT(1H ,3F15.1)
STOP
END

004144 DATA.	17204060000000000000	CON. BSS LB	DATA 17204000000000000000
004144 DATA.	17214000000000000000		DATA 17214000000000000000
004145 DATA.	17227000000000000000		DATA 17227000000000000000
004146 DATA.	00000000000000000000		DATA 00000000000000000000
004147 DATA.	00000000000000000000		EXT STOP.
			EXT OUTC1.
			EXT SUB
			EXT QENTRY.
004150 DATA.		X	BSS 18
004151 DATA.		Y	BSS 18
004152 DATA.		Z	BSS 18
			USE CODE.
			LI
004113 CODE.	5150004144 DATA.	SA5	CON.
	5140004145 DATA.	SA4	CON.+13
004114 CODE.	5130004146 DATA.	SA3	CON.+23
	5110004125 CODE.	SA1	LAP1
004115 CODE.	10755	6X7	X5
	300+5	FX0	X6+X5
	10044	6X6	X4
004116 CODE.	5170004150 DATA.	SA7	X
	24700	6X7	80,X0
004117 CODE.	5160004151 DATA.	SA6	Y
	10633	6X6	X3
004120 CODE.	5170004132 CODE.	SA7	ST.
	5160004152 DATA.	SA6	Z
004121 CODE.	0100000000 <EXT>	RJT	SUB,54
	005004111		
004122 CODE.	5110004133 DATA.	SA1	LI01
004123 CODE.	0100000000 <EXT>	RJT	OUTC1.,60
	006004111		
004124 CODE.	5120004147 DATA.	SA1	CON.+78
	0400000000 <EXT>	EQ	STOP.
004125 CODE.	00000000000000000000 DATA.	LAP1	BSS
004125 CODE.	00000000000000000000 DATA.	APL	X
004126 CODE.	00000000000000000000 DATA.	APL	X
004127 CODE.	00000000000000000000 CODE.	APL	ST.
004130 CODE.	00000000000000000000 DATA.	APL	Z
004131 CODE.	00000000000000000000	APL	
004132 CODE.		ST.	BSS
		Z.	18
			END
			START

APPENDIX 2.1. The subroutine SUB in FORTRAN and its COMPASS expansion when the user requested optimization level one /OPT=1/ to the FTN.

ITIME SUB 73774 OPT=1 FTN 4.6+42^s

```
SUBROUTINE SUB(A,R,C,D)
  B=A+A
  D=A+C
  RETURN
END
```

COMPASS code as generated by the FTN compiler :

			USE DATA.	
			USE START.	
			TRACE	SUB,SUB,4B
000000 START.	23250215555555000003			
000001 START.	00050000000000000000			
000002 START.	51300010010203000000		FENTRY	SUB,FENTRY.,0,C
000003 START.	J4034000036100045500			
000004 START.	74600540105160000000			
			FORPAR	A
			FORPAR	B
			FORPAR	C
			FORPAR	D
			USE DATA..	
			USE DATA.	
			USE	CODE.
000005 CODE.	54500		SA5	A0
	5046000002		SA4	A0+2B
000006 CODE.	5030000001		SA3	A0+1B
	5120000003		SA2	A0+3B
000007 CODE.	53150		SA1	X5
	30711		FX7	X1+X1
	53540		SA5	X4
	53730		SA7	X3
000010 CODE.	30051		FX0	X5+Y1
	24000		HX6	B0,X0
	53020		SA6	X2
000011 CODE.	J400000002	START.	LQ	EXIT.
			Z.	END

The results obtained in this case (they indicate a call by value/result parameter passing mechanism):

X = 2.0 Y = 2.0 Z = 4.0

The FORTRAN Extended Reference manual^{1/1} states that the call by name constitutes the standard parameter passing mechanism, associated with a call by value for expressions appearing as actual parameters. Differences between different optimization levels in this respect are not mentioned.

PROGRAM START (INPUT,OUTPUT)

X=1.0

Y=2.0

Z=7.0

CALL SUB(X,X,X+Y,Z)

PRINT 1,X,Y,Z

1 FORMAT(1H,3F15.1)

STOP

END

004151 DATA. 17004000300000000000
 004151 DATA. 17214010000000000000
 004152 DATA. 17227000000000000000
 004153 DATA. 00000000000000000000
 004154 DATA. 00000000000000000000

004155 DATA.
 004156 DATA.
 004157 DATA.

004114 CODE. 6102000002
 5150004151 DATA.

004115 CODE. 10755
 5170004155 DATA.

004116 CODE. 6102000003
 5150004152 DATA.

004117 CODE. 10755
 5170004156 DATA.

004120 CODE. 6112001004
 5150004153 DATA.

004121 CODE. 10755
 5170004157 DATA.

004122 CODE. 6112001005
 5150004155 DATA.

004123 CODE. 5140004156
 5110004132 CODE.

004124 CODE. 30645
 24700
 5170004137 CODE.

004125 CODE. 0100001000
 0050004111 <EXT>

004126 CODE. 6102001006
 5110004100 DATA.

004127 CODE. 0100000000
 00060004111 <EXT>

004130 CODE. 6102000010
 5110004154 DATA.

004131 CODE. 0400001000 <EXT>

004132 CODE. 00000000000000000000

004132 CODE. 00000000000000000000

004133 CODE. 00000000000000000000

004134 CODE. 00000000000000000000

004135 CODE. 00000000000000000000

004136 CODE. 00000000000000000000

004137 CODE. 00000000000000000000

CON. BSS 08
 DATA 17204000000000000000
 DATA 17214000000000000000
 DATA 17227000000000000000
 DATA 00000000000000000000

EXT STOP.
 EXT OUTCI.
 EXT SUB
 EXT FTNRPV.

X BSS 10
 Y BSS 10
 Z BSS 10
 * USE CODE.

S60 B2+2B
 SA5 CON.
 BX7 X5
 SA7 X

S80 B2+3B
 SA5 CON.+10
 BX7 X5
 SA7 Y

S80 B2+1B
 SA5 CON.+2B
 BX7 X5
 SA7 Z

S80 B2+5B
 SA5 X
 SA4 Y
 SA1 (AP1

FX0 XL+X5
 NX7 BB,X0
 SA7 ST.
 RJT SUR,5B

S80 B2+6B
 SA1 1101
 RJT OUTCI.,5B

S80 B2+109
 SA1 CON.+3B
 EQ STOP.

(AP1 BSS 08
 APL X
 APL X
 APL ST.
 APL Z
 APL
 ST. BSS 10
 Z. END START

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