



Haskell Monads



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Why Monads?

- ▶ Common solution to functional programming challenges
 - ▶ Exceptions
 - ▶ State (side effects)
 - ▶ I/O
- ▶ Advanced C++ TMP
 - ▶ Eric Niebler's (Joel Falcou's) Proto
 - ▶ Mixed compile-time runtime programming



Plan

- ▶ Computations vs. functions
- ▶ The Maybe monad
 - ▶ Type constructor
 - ▶ Bind—composing monadic functions
 - ▶ Return
 - ▶ The do notation
- ▶ The State Monad
 - ▶ Actions and functions returning actions
 - ▶ Evaluating expression trees



Teaser

```
template<class L, class R>
struct Compile<Plus<L, R>> : Prog {
    int operator()(Args args) {
        return Bind<Compile<L>, Bind<Compile<R>, Return>> (
            Compile<L>(),
            [](int left) -> Bind<Compile<R>, Return> {
                return Bind<Compile<R>, Return>(
                    Compile<R>(),
                    [left](int right) -> Return {
                        return Return(left + right);
                    }
                );
            }
        )(args);
    }
};
```



Computations vs. Functions

- ▶ Computations that are not functions
 - ▶ Not defined for all values of arguments (Errors, Exceptions)
 - ▶ Nondeterministic: returning a set of options (Parsers)
 - ▶ Side effects and state
 - ▶ Input/Output
- ▶ Functions that return “enriched” types like:
 - ▶ Maybe, Errors, Exceptions
 - ▶ Lists
 - ▶ Actions
 - ▶ I/O Actions



Monads

- ▶ Theoretical foundations: Category theory
- ▶ Elements of a monad (**Kleisli triple**)
 - ▶ Type constructor: a parameterized “enriched” type
 - ▶ Bind: composition of monadic functions
 - ▶ Return: Encapsulation of values into enriched types





The Maybe Monad

A toy example that introduces all the elements of a monad

The Maybe Type

- ▶ Motivation

```
size_t off = fileName.find('.');
string ext = fileName.substr(off, fileName.length() - off);
```

- ▶ Type constructor

- ▶ For all types a:
- ▶ Special value Nothing, or
- ▶ Just a

```
data Maybe a = Nothing | Just a
```



Haskell Data Types

- ▶ Data is immutable
 - ▶ Data “remembers” how it was created
 - ▶ Pattern matching used to extract this information

```
data Maybe a = Nothing | Just a
showMaybe m =
  case m of
    Nothing -> "Nothing"
    Just x -> "Something " ++ (show x)
```

```
> let x = Nothing
> showMaybe x
"Nothing"
> let y = Just 15
> showMaybe y
"Something 15"
```



C++ Maybe

```
data Maybe t = Nothing | Just t
```

```
enum MaybeTag { Nothing, Just };
```

```
template<class T>
struct Maybe {
    MaybeTag tag;
    T value; // valid if tag is Just
};
```

```
Maybe<size_t> off = safe_find(fileName, '.');
std::string ext;
if (off.tag == Just)
    ext = fileName.substr(off.value, fileName.length() - off.value);
```



Composing Maybe's in C++

```
Maybe<Foo> y = f(x);
if (y.tag == Just) {
    Maybe<Bar> v = g(y.value);
    if (v.tag == Just) {
        Maybe<Baz> z = h(v.value);
        if (z.tag == Just) {
            return z;
        }
    }
}
DO { // Ideally!
    auto y = f(x);
    auto v = g(y);
    auto z = h(v);
    return z;
}
```



Composing Maybe's in Haskell

```
compose n =  
  let m = f n in  
  case m of  
    Nothing -> Nothing  
    Just v1 ->  
      let m1 = g v1 in  
        case m1 of  
          Nothing -> Nothing  
          Just v2 ->  
            let m2 = h v2 in  
              case m2 of  
                Nothing -> Nothing  
                Just v3 -> v3
```

Abstracting “the rest of the code” into a continuation

```
 $\lambda v1 \rightarrow$   
  let m1 = g v1 in  
  case m1 of  
    Nothing -> Nothing  
    Just v2 ->  
      let m2 = h v2 in  
        case m2 of  
          Nothing -> Nothing  
          Just v3 -> v3
```



Abstracting the Glue

```
compose n =  
let m = f n in  
case m of  
Nothing -> Nothing  
Just v1 ->  
    let m1 = g v1 in  
        case m1 of  
        Nothing -> Nothing  
        Just v2 ->  
            let m2 = h v2 in  
                case m2 of  
                Nothing-> Nothing  
                Just v3-> v3
```

```
compose n =  
let m = f n in  
--  
--  
bind m ( $\lambda$  v1 ->  
    let m1 = g v1 in  
        case m1 of  
        Nothing -> Nothing  
        Just v2 ->  
            let m2 = h v2 in  
                case m2 of  
                Nothing-> Nothing  
                Just v3-> v3)
```



Monadic Bind

- ▶ Takes a Maybe
- ▶ Takes a continuation
- ▶ Returns a Maybe

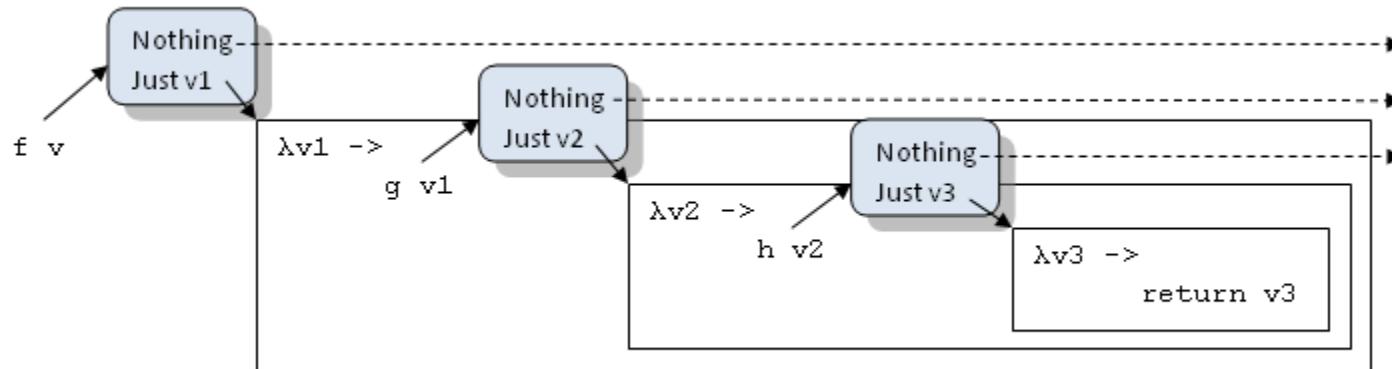
```
bind :: Maybe a -> (a -> Maybe b) -> Maybe b
```

```
bind m f =  
  case m of  
    Nothing -> Nothing  
    Just v -> f v
```

```
-- compact form  
bind (Just x) f = f x  
bind Nothing f = Nothing
```



Cascade of Continuations



Infix Notation

```
(>>=) :: Maybe a -> (a -> Maybe b) -> Maybe b
Nothing >>= cont = Nothing
(Just x) >>= cont = cont x
```

```
compose n =
  f n >>= λ v1 ->
    g n1 >>= λ v2 ->
      h n2 >>= λ v3 ->
        return v3
```



The *return* Function

- ▶ Does for values what the type constructor does for types
- ▶ Wraps any value into enriched type
- ▶ Trivial for the Maybe monad

```
return v = Just v
```



The *do* Notation

```
compose v =  
  f v >>= λ v1 ->  
    f v1 >>= λ v2 ->  
      f v2 >>= λ v3 ->  
        return v3
```

```
compose v = do  
  v1 <- f v  
  v2 <- g v1  
  v3 <- h v2  
  return v3
```

- ▶ Just syntactic sugar
- ▶ Don't confuse left arrow with assignment



C++ “do” Notation?

```
int compose(int v) {  
    auto v1 = f(v);  
    auto v2 = g(v1);  
    auto v3 = h(v2);  
    return v3;  
}  
  
try {  
    compose(x);  
} catch(...) {  
    // error handling  
}
```

```
compose v = do  
    v1 <- f v  
    v2 <- g v1  
    v3 <- h v2  
    return v3
```

- ▶ C++ exceptions
- ▶ Haskell Maybe, Error, and Exception monads
- ▶ Type safety (exception specification?)
- ▶ Not a general pattern in C++





Dealing with State

Short intro to the state monad

State and Side Effects

- ▶ Computation may access global/static variables
- ▶ Modeled by a function that
 - ▶ Takes state as argument
 - ▶ Returns (possibly modified) state
 - ▶ Together with regular return value
- ▶ Called “action”

$f :: \text{State} \rightarrow (\text{State}, t)$



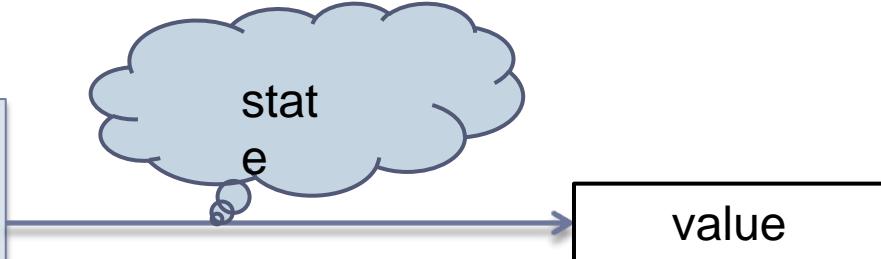
Delayed Execution

- ▶ Problem: Composing actions is messy
- ▶ Ideal:
 - ▶ Hide state (as if it were global) but
 - ▶ Take advantage of the strong typing of actions
 - ▶ Enforce proper sequencing of actions
- ▶ Solution
 - ▶ Separate action composition from action execution
 - ▶ Compose (higher-order) functions returning actions
 - ▶ Execute the final action by providing initial state
- ▶ Result: State Monad

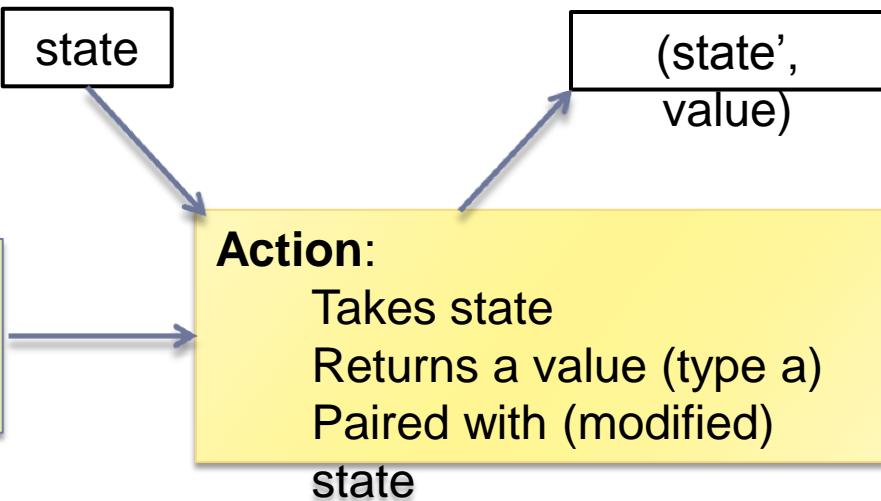


State Monad Pictorial

Computation:
Takes some arguments
Returns a value (type)
Accesses (mutable)
state



Monadic Function:
Takes some arguments
Returns an action





The Expression Monad

An example of a state monad

Expression Monad

- ▶ It's an example of a reader monad
 - ▶ It's a state monad with read-only state
- ▶ Expression trees ($\text{Arg1} * 3 + \text{Arg2}$) are constructed from
 - ▶ Constant (integer) nodes
 - ▶ Special placeholder nodes, Arg1 and Arg2
 - ▶ Plus and Times nodes
- ▶ State is a list of (two) arguments to an expression
- ▶ An action evaluates a corresponding expression given arguments
- ▶ Expression trees drive the composition of actions



Expression

- ▶ Recursive definition
- ▶ Tagged union

```
data Exp = Const Integer
          | Plus Exp Exp
          | Times Exp Exp
          | Arg1
          | Arg2
```

- ▶ State
 - ▶ Type alias, *type*, like C++ `typedef`

```
type Args = [Integer]
```



Action in the Reader Monad

Args → a

- ▶ Type constructor
 - ▶ *newtype* creates a new type
 - ▶ Haskell limitation: can't use a type alias “*type*”

newtype Prog a = PR (Args → a)

- ▶ Auxiliary function: runs a program given input arguments

run :: Prog t → Args → t

run (PR act) args = act args



Monadic Functions

- ▶ **getArg**
 - ▶ Creates an action that extracts n'th argument from input

```
getArg :: Int -> Prog Integer
getArg n = PR (λ args -> args !! n))
```

- ▶ **double**
 - ▶ For some n, creates an action that returns $2 * n$
 - ▶ It's a closure (captures n)

```
double n = PR (λ args -> 2 * n)
```

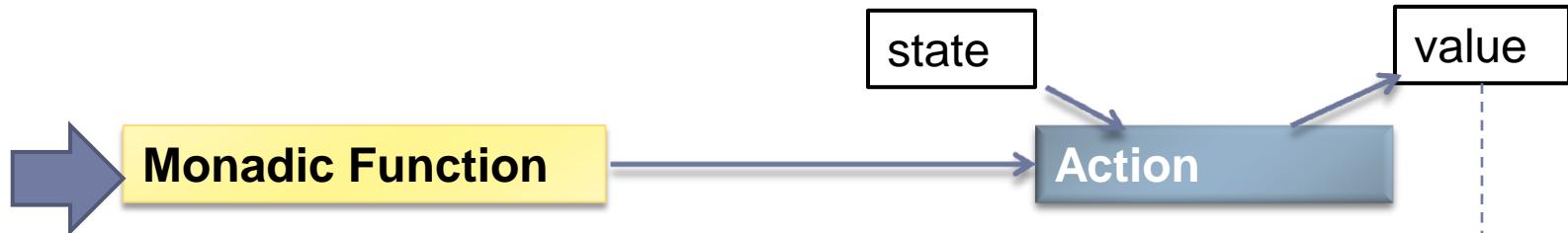


Bind

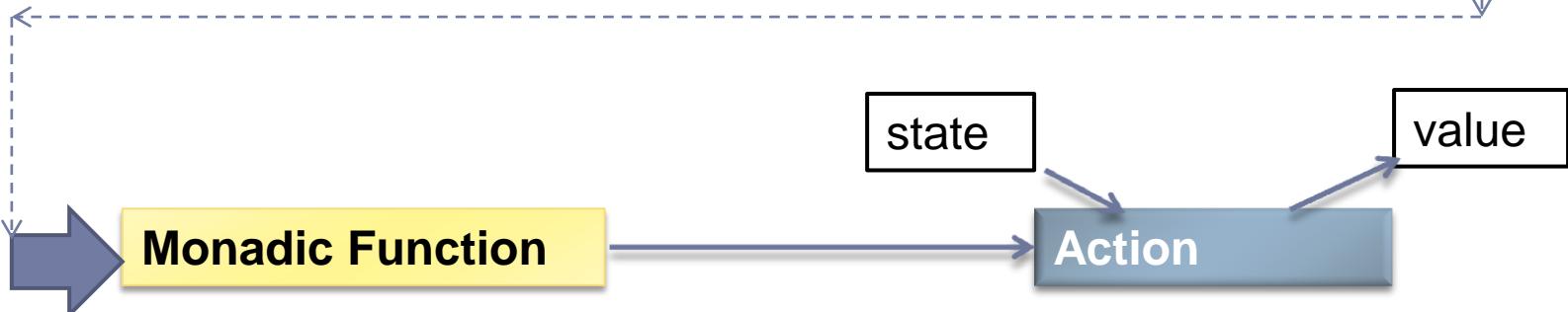
- ▶ Goal: create a new action combining, for instance
 - ▶ `getArg 0`
 - ▶ `double v`
- ▶ `getArg 0` creates an action
- ▶ `double v` forms the continuation (the rest)
- ▶ bind
 - ▶ takes an action: `Prog a`
 - ▶ Returned by `getArg 0`
 - ▶ a continuation: `a -> Prog b`
 - ▶ `$\lambda v \rightarrow \text{double } v$`
 - ▶ and returns action: `Prog b`



Bind Pictorial



continuation



Bind

```
bind :: (Prog a) -> (a -> (Prog b)) -> (Prog b)
```

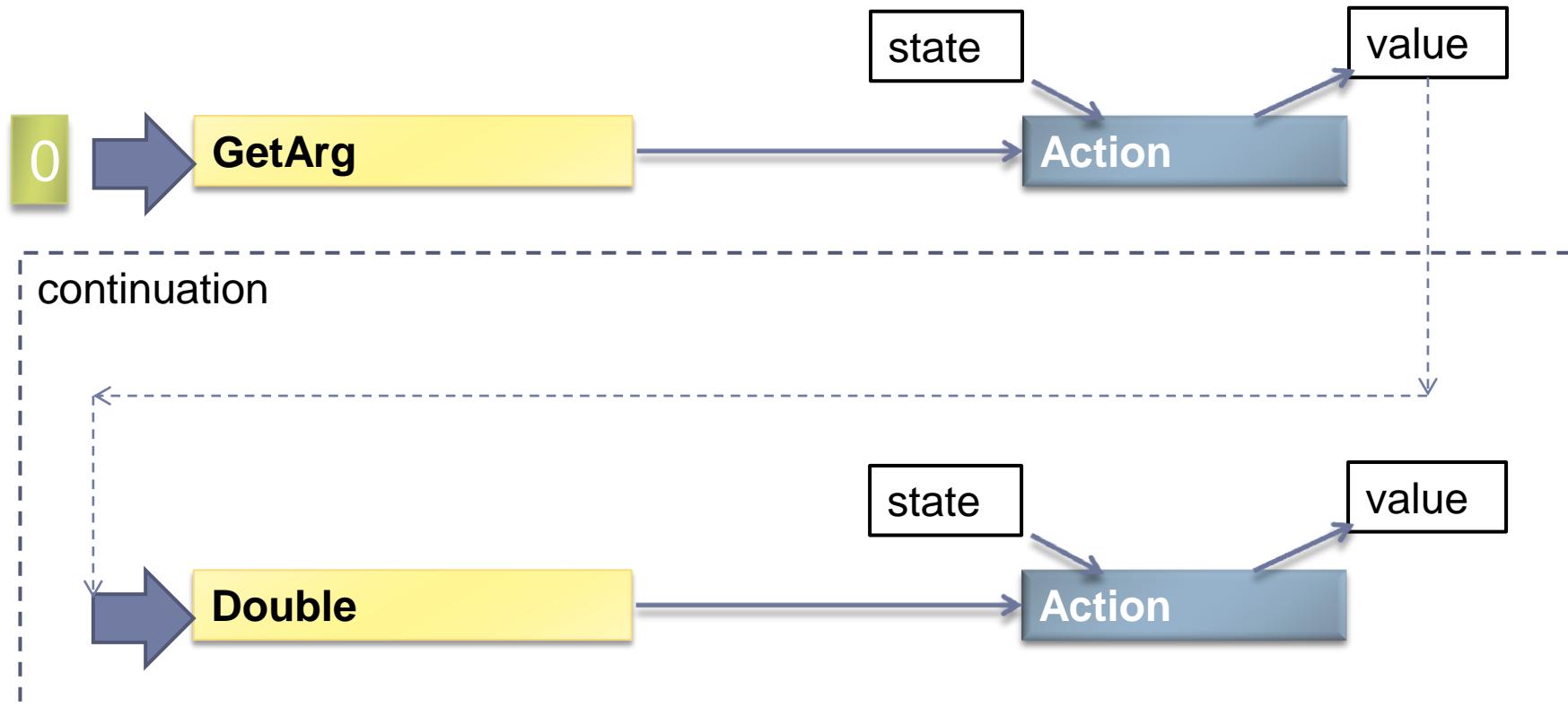
- ▶ Returns an action: a lambda of appropriate type

```
bind (PR act) cont =  
    PR ( $\lambda$  args -> ... produce value of type b ...)
```

- ▶ This lambda will be executed when args available

```
bind (PR act) cont =  
    PR ( $\lambda$  args ->  
        let v = act args  
        (PR act') = cont v  
    in  
        act' args) --- produces the final value
```

Composing Actions Pictorial



Composing Actions

- ▶ Goal: create a new action combining
 - ▶ `getArg 0`
 - ▶ `double v`

```
test0 :: Prog Integer
test0 =
    bind (getArg 0) ( $\lambda v \rightarrow$  double v)
```

```
> let prog = test0
> run prog [3,4]
6
> run prog [11,0]
22
```



The Reader Monad

- ▶ Reusing Haskell type-class Monad (the *instance* declaration)
- ▶ *return* and *bind* applicable to any reader monad
- ▶ Bind as an infix operator $\gg=$

```
instance Monad Prog where
    -- return :: a -> Prog a
    return v = PR (λ args -> v)
    -- (gg=) :: Prog a -> (a -> Prog b) -> Prog b
    (PR act) gg= cont = PR (λ arg -> let v = act arg
                                (PR act') = cont v
                                in act' arg)
```

```
test0 =
    bind (getArg 0)
        (λ v -> double v)
```

```
test1 = do
    v <- getArg 0
    double v
```



Monadic Programming

- ▶ Define the *compile* monadic function
 - ▶ Given an expression produces a program (action) to calculate this expression
 - ▶ Will be “specialized” for various expression patterns
 - ▶ Composes smaller monadic functions into larger ones

```
compile :: Exp -> Prog Int
```



Definition of the compile function

- ▶ Matching Const node

```
compile (Const c) = return c
```

- ▶ Matching plus node: recursive calls

```
compile (Plus e1 e2) =
  do
    v1 <- compile e1
    v2 <- compile e2
    return (v1 + v2)
```

- ▶ Matching Arg1 node

```
compile Arg1 = getArg 0
```

Testing

- ▶ Compile an expression: $x * y + 13$

```
testExp =  
  let exp = (Plus (Times Arg1 Arg2) (Const 13))  
  in compile exp
```

- ▶ Run compiled expression with input [3, 4]

```
> let args = [3, 4]  
> Let prog = testExp  
> run prog args
```

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Conclusion

- ▶ Similar patterns in other languages
 - ▶ Command pattern: creates command objects (actions?), combined using Composite pattern (bind?)
 - ▶ Lambdas and closures may be returned from functions in C++0x
 - ▶ Help with inversion of control
- ▶ EDSLs as monads
- ▶ Further study
 - ▶ Type classes
 - ▶ IO monad



Bibliography

- ▶ Mike Vanier's blog (Haskell monads):
<http://mvanier.livejournal.com/3917.html>
- ▶ My blog (extended treatment of current presentation):
<http://bartoszsmilewski.wordpress.com/2011/01/09/monads-for-the-curious-programmer-part-1/>
- ▶ Brian McNamara, Yannis Smaragdakis, [Syntax sugar for FC++: lambda, infix, monads, and more.](#)

