



Abusing the Bitmask

A short story by Nicolas Waisman

nicolas@immunityinc.com

<http://twitter.com/nicowaisman>

Who am i?



- Regional Manager at Immunity, Inc.
- Research and Development of reliable Heap Overflow exploitation for CANVAS attack framework
- Big fan of Guarana.

Once upon a time



- There was a bug on **MS NTTP server ...**
- **MS04-036** was the starting point for finding new ways to exploit the heap.
- This is how this the bitmask heap technique born.

Public Exploits



V
S

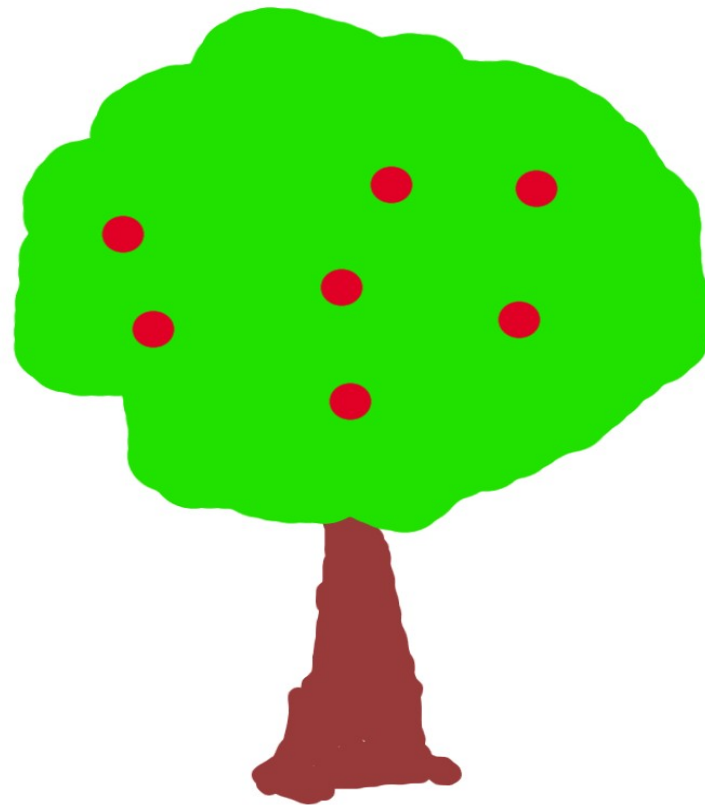
Commercial Exploits



Windows HEAP



The heap provides a mechanism for Allocating, Reallocation and freeing dynamic memory.



Windows HEAP



ntdll.dll

How was the heap exploited?

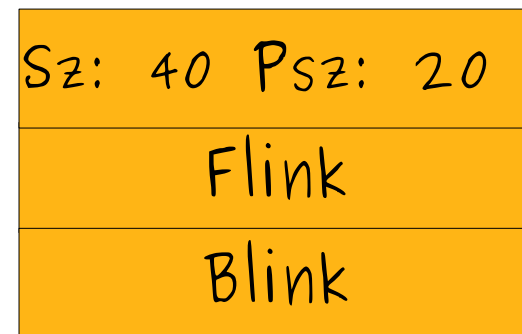
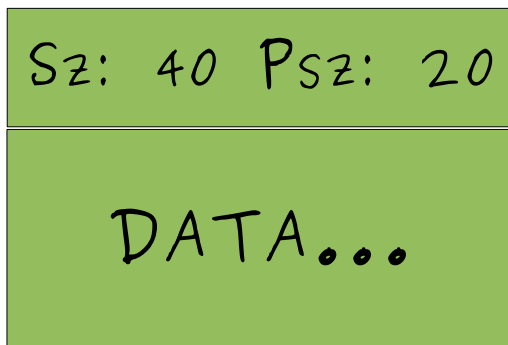


- In the past, research had taken advantage of the (in)famous “unlink” technique.
- This technique could allow an attacker to write 4 bytes, whenever they want.
- 4 bytes + some effort = Shellcode Execution

Learn heap exploitation in 3 minutes



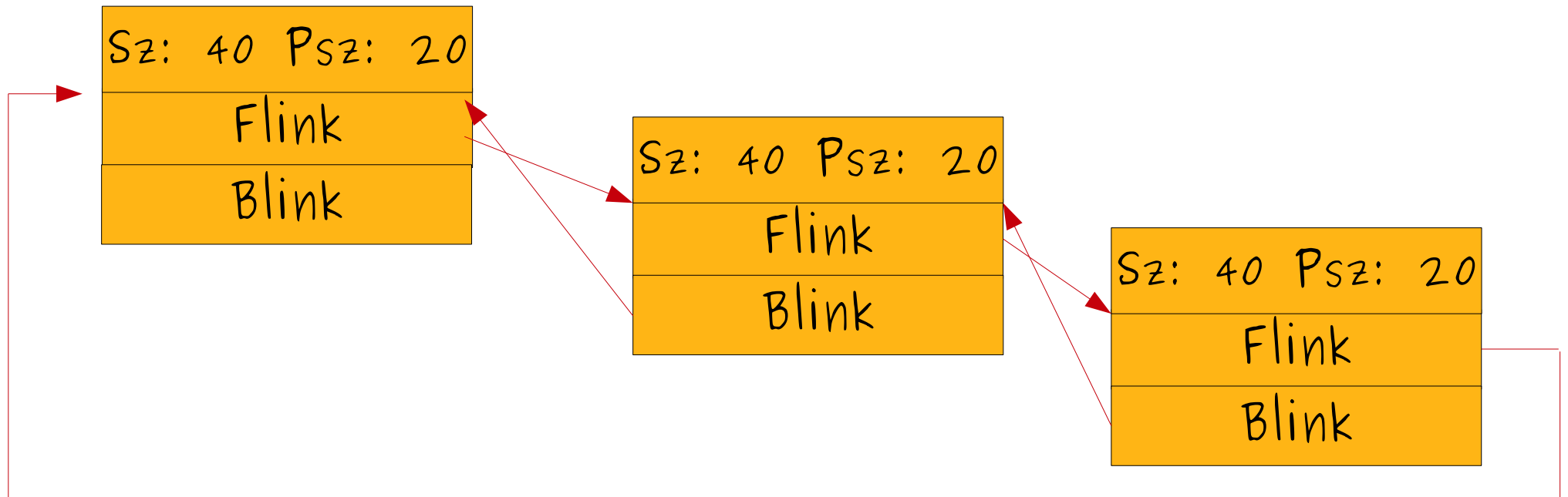
- Every piece of heap dynamic memory, are represented as chunks.
- Chunks has headers



Learn heap exploitation in 3 minutes



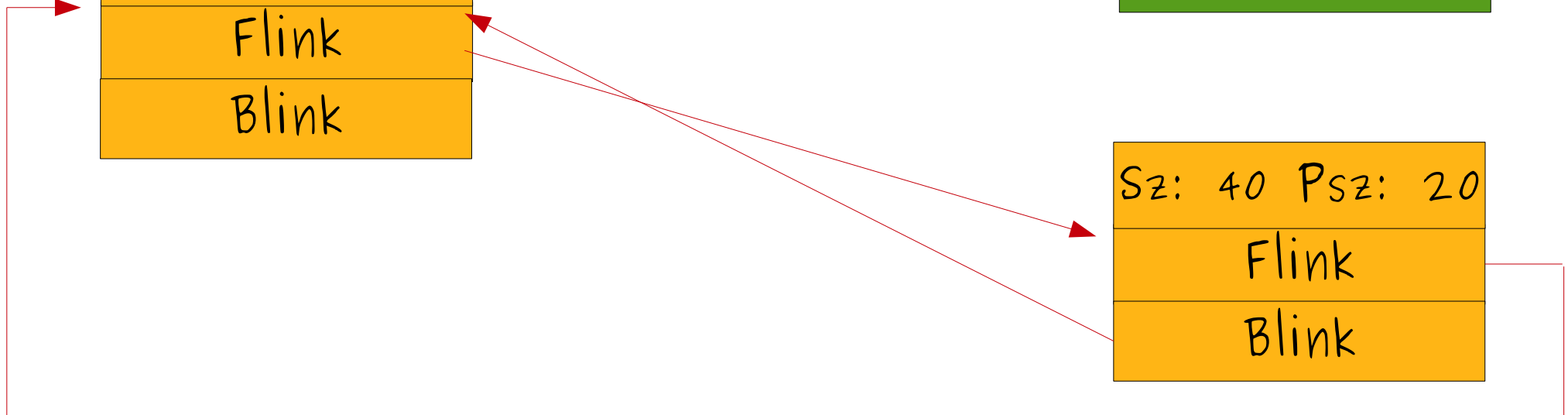
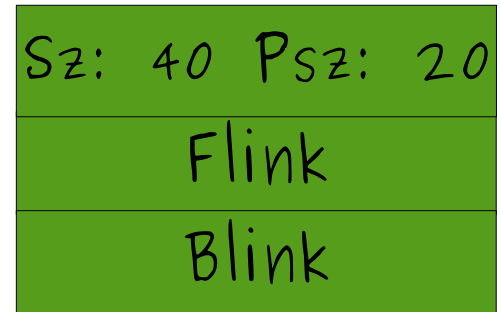
- Free chunks are connected forwardly and backwardly.



Learn heap exploitation in 3 minutes



- When memory is required, it will take one of the free chunk and unlink it.



What if we overflow the chunk before
being freed?

Demo



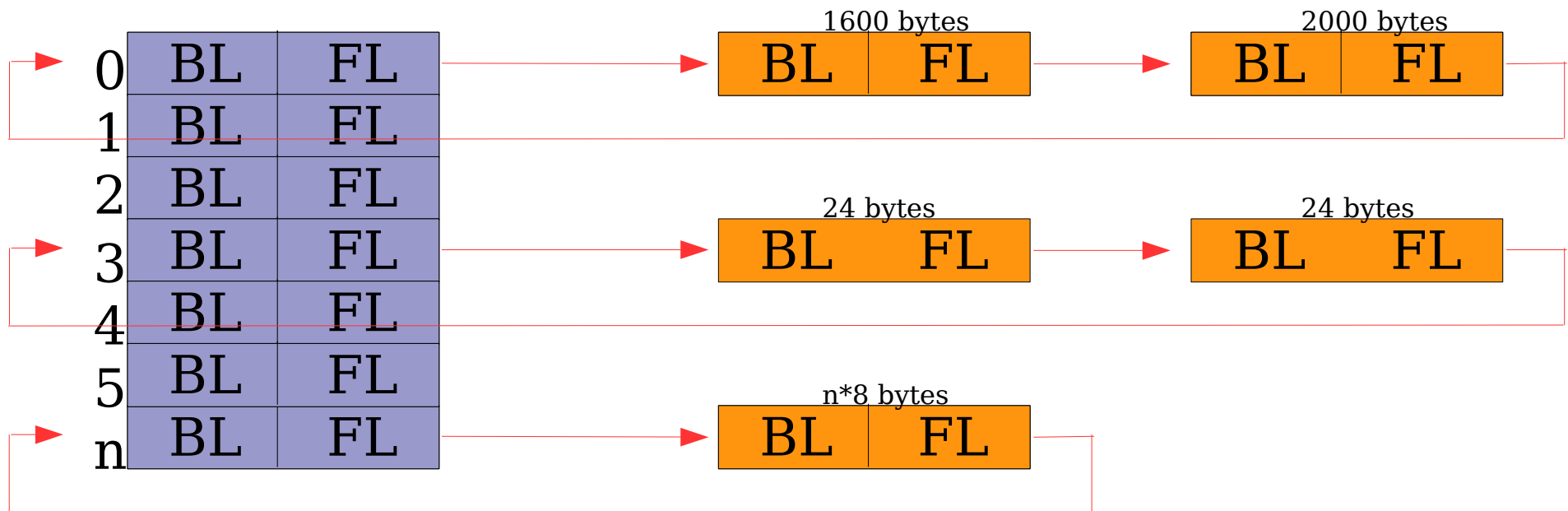
Need three volunteer from the audience

FreeList



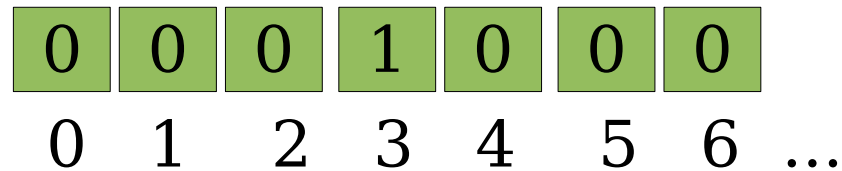
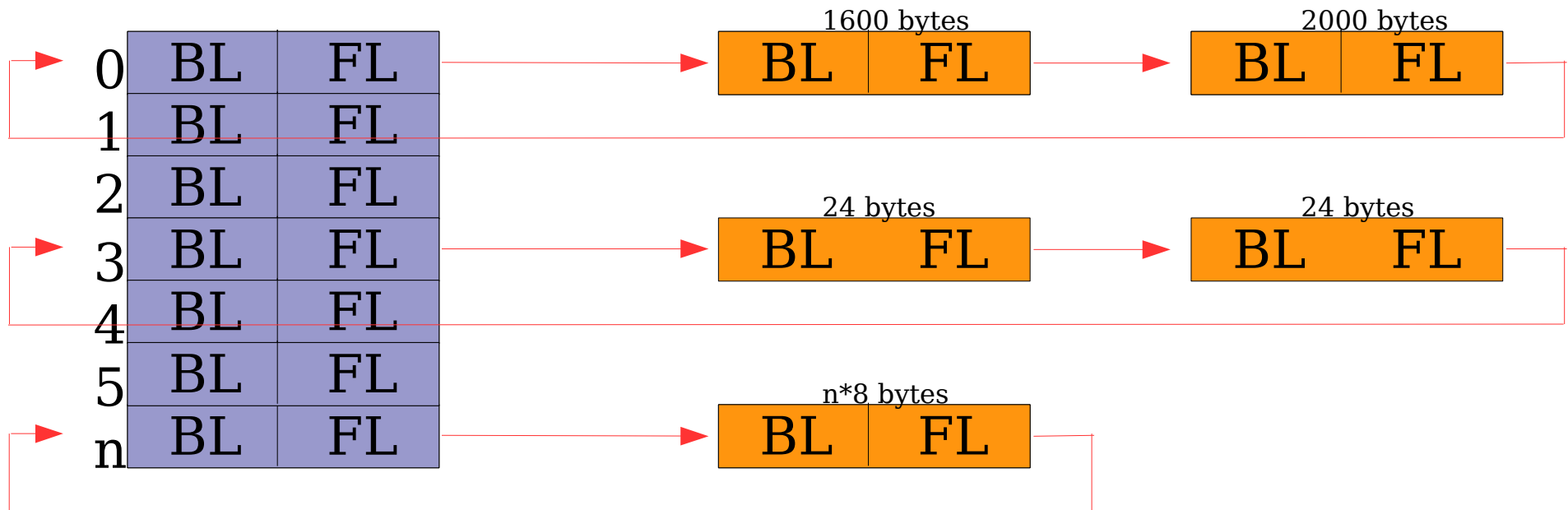
- For each “List” of connected chunks, there is a slot on the FreeList structure.
- This speed up the search process.

FreeList



where $n < 128$

FreeListInUse





Address	Chunks
0x00000000	### Immunity's Heapdump ###
0x00150000	Dumping heap: 0x00150000
0x00150000	Flags: 0x00000002 ForceFlags: 0x00000000
0x00150000	Total Free Size: 0x00001d7b VirtualMemoryThreshold: 0x0000fe00
0x00000000	Segment[0]: 0x00150000
0x00000000	FreeListInUse 10011001011001110100110100000100 0000000000000000000001000000000010
0x00000000	01000000000000000000000000000000010 00000000000000000000000000000000
0x00150178	[000] 0x00150178 -> [0x001d07c0 0x00212c80]
0x00212c80	0x00212c80 -> [0x00150178 0x001f43c8] (00000471)
0x001f43c8	0x001f43c8 -> [0x00212c80 0x001bc008] (00000405)
0x001bc008	0x001bc008 -> [0x001f43c8 0x001d9dc8] (00000401)
0x001d9dc8	0x001d9dc8 -> [0x001bc008 0x001d5110] (000002a0)
0x001d5110	0x001d5110 -> [0x001d9dc8 0x001dd848] (00000202)
0x001dd848	0x001dd848 -> [0x001d5110 0x001e3008] (00000202)
0x001e3008	0x001e3008 -> [0x001dd848 0x001d0eb8] (00000150)
0x001d0eb8	0x001d0eb8 -> [0x001e3008 0x00204008] (00000141)
0x00204008	0x00204008 -> [0x001d0eb8 0x001a62a8] (0000012d)
0x001a62a8	0x001a62a8 -> [0x00204008 0x001c1878] (00000116)
0x001c1878	0x001c1878 -> [0x001a62a8 0x001d07c0] (0000009c)
0x001d07c0	0x001d07c0 -> [0x001c1878 0x00150178] (00000081)
0x00150180	[001] 0x00150180 -> [0x00150180 0x00150180]
0x00150188	[002] 0x00150188 -> [0x001becf8 0x001ab500]
0x001ab500	0x001ab500 -> [0x00150188 0x00188d40] (00000002)
0x00188d40	0x00188d40 -> [0x001ab500 0x001c4100] (00000002)
0x001c4100	0x001c4100 -> [0x00188d40 0x0018e868] (00000002)
0x0018e868	0x0018e868 -> [0x001c4100 0x001a5b48] (00000002)
0x001a5b48	0x001a5b48 -> [0x0018e868 0x0019b9a0] (00000002)
0x0019b9a0	0x0019b9a0 -> [0x001a5b48 0x00197da8] (00000002)
0x00197da8	0x00197da8 -> [0x0019b9a0 0x001d50b8] (00000002)
0x001d50b8	0x001d50b8 -> [0x00197da8 0x00201660] (00000002)
0x00201660	0x00201660 -> [0x001d50b8 0x001a9fd8] (00000002)
0x001a9fd8	0x001a9fd8 -> [0x00201660 0x001f70c8] (00000002)
0x001f70c8	0x001f70c8 -> [0x001a9fd8 0x001edde8] (00000002)
0x001edde8	0x001edde8 -> [0x001f70c8 0x001dcca0] (00000002)
0x001dcca0	0x001dcca0 -> [0x001edde8 0x001e0ff8] (00000002)
0x001e0ff8	0x001e0ff8 -> [0x001dcca0 0x001ffb58] (00000002)
0x001ffb58	0x001ffb58 -> [0x001e0ff8 0x001cb9b0] (00000002)
0x001cb9b0	0x001cb9b0 -> [0x001ffb58 0x001becf8] (00000002)
0x001becf8	0x001becf8 -> [0x001cb9b0 0x00150188] (00000002)
0x00150190	[003] 0x00150190 -> [0x00150190 0x00150190]
0x00150198	[004] 0x00150198 -> [0x00150198 0x00150198]
0x001501a0	[005] 0x001501a0 -> [0x001501a0 0x001501a0]
0x001501a8	[006] 0x001501a8 -> [0x001501a8 0x001501a8]



DISCLAIMER

The events and characters in this presentation are fictitious. Any similarity to actual persons, living or dead, is purely coincidental

FreeListInUse: Analogy



At least,
One chunk

Vulnerability in NNTP Could Allow Remote Code Execution

- Found in October 2004 by two Argentinian researchers: Lucas “Rompedor” Lavarello y Juliano Rizzo
- Various parsing errors of the XPAT command could lead into part of the heap being overwritten.
- The challenge was to exploit an **off-by-two**

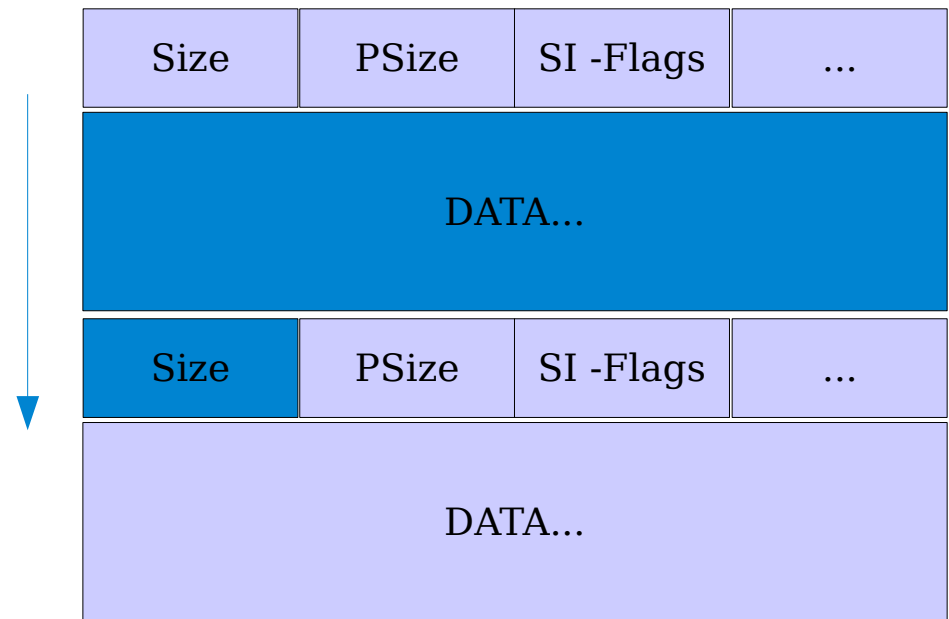
Moral: “Everything good start with a bug”



Off By Two

Based on a 2000 bytes chunks, we were able to overwrite the next chunk with only two bytes.

We are only able to control one of the two characters:



`\x41\x41\x41\x41\x41\x00`

How would you Exploit it?



- Lookaside: Nothing interesting
- Busy Chunk: Forward coalescation
- Free Chunk: Make the size bigger.
- Last Chunk: Make the size bigger.
- New Technique. (The whole point of this talk)

Size	PSize	SI -Flags	...
DATA...			

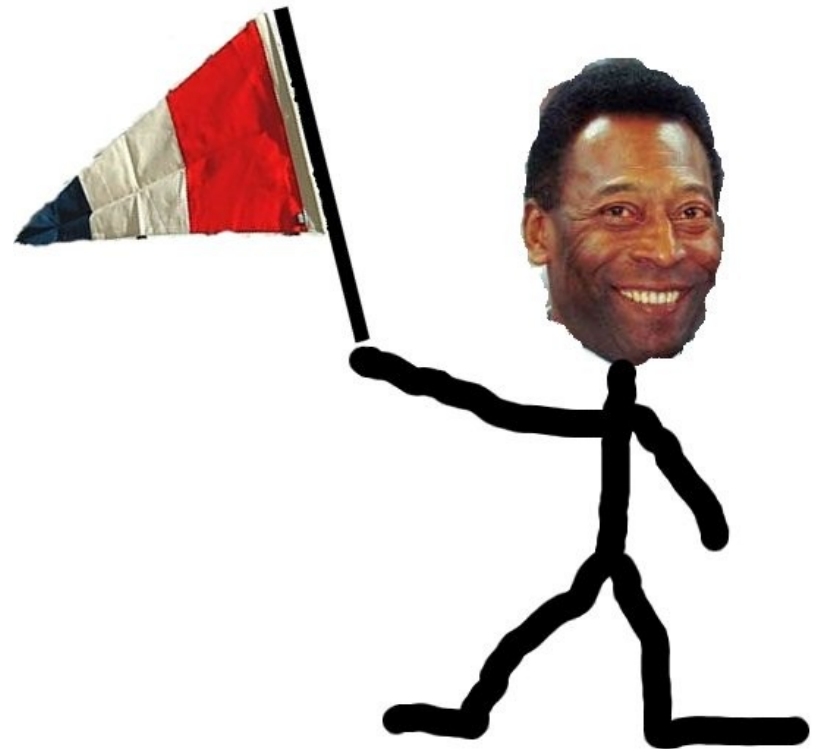
"A technique should take less requirements than a strawberry pudding recipe"

strawberry pudding law
- Sinan Eren

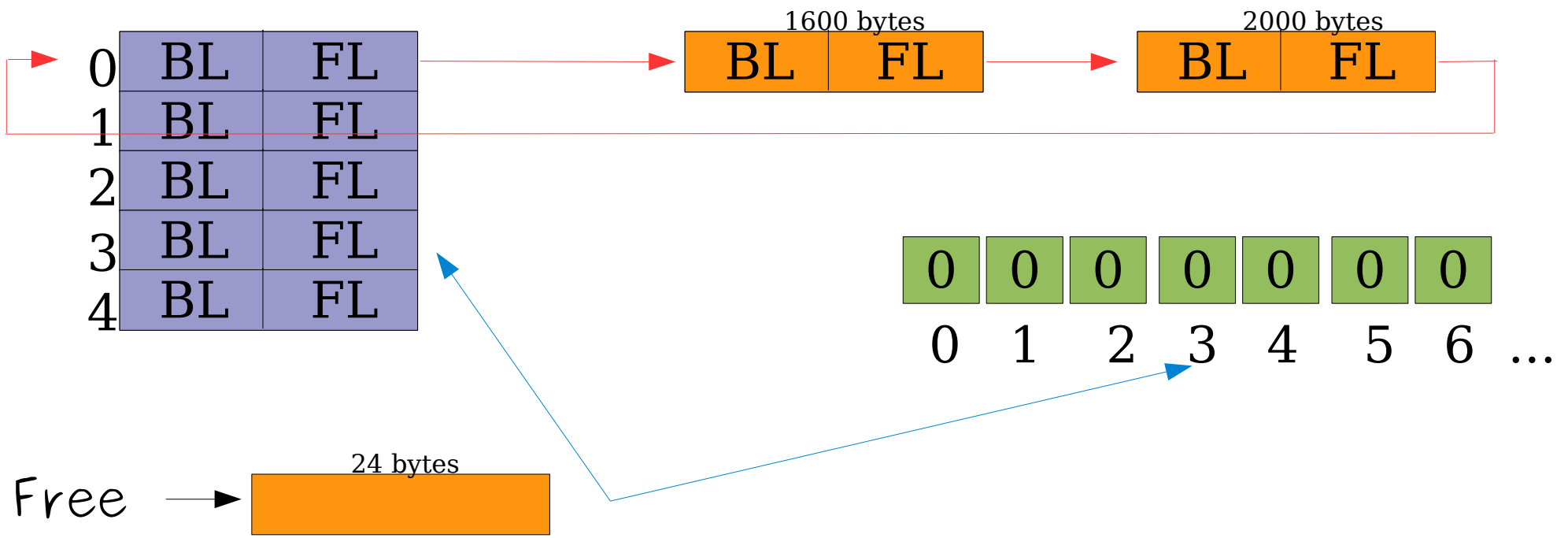


Interaction with the FreeListInUse

Set bit: When a FreeList slot of some size is **empty** and you want to free a chunk there.



Interaction with the FreeListInUse

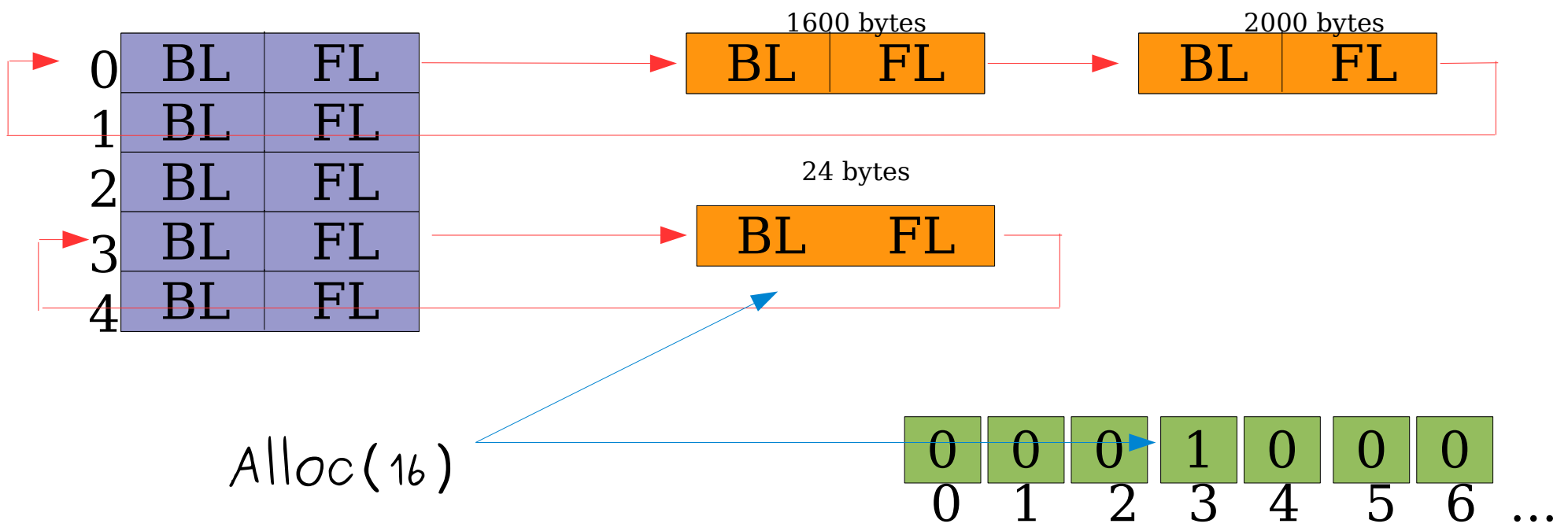


Interaction with the FreeListInUse

Unset bit: When you allocate a chunk which
is the **only member** of a FreeList Slot.



Interaction with the FreeListInUse



In Code...



CPU - thread 0000CE4, module ntdll

Address	Hex	Assembly	Comment
7C9111FE	0FB70E	MOVZX ECX,WORD PTR DS:[ESI]	ecx = Chunk->Size
7C911201	8BC1	MOV EAX,ECX	
7C911203	C1E8 03	SHR EAX,3	ByteInFreeList = Size / 8
7C911206	8985 28FFFFFF	MOV DWORD PTR SS:[EBP-D8],EAX	
7C91120C	83E1 07	AND ECX,7	EntryInByte = Size & 7
7C91120F	33D2	XOR EDX,EDX	
7C911211	42	INC EDX	
7C911212	D3E2	SHL EDX,CL	ByteToSet = 1 << EntryInByte
7C911214	8995 04FFFFFF	MOV DWORD PTR SS:[EBP-FC],EDX	
7C91121A	8D8418 58010000	LEA EAX,DWORD PTR DS:[EAX+EBX+158]	
7C911221	33C9	XOR ECX,ECX	
7C911223	8A08	MOV CL,BYTE PTR DS:[EAX]	tmpbyte = FreeListInUse[ByteInFreeList]
7C911225	33CA	XOR ECX,EDX	tmpbyte = xor(tmpbyte, ByteToSet)
7C911227	8808	MOV BYTE PTR DS:[EAX],CL	FreeListInUse[ByteInFreeList] = tmpbyte
7C911229	E9 18020000	JMP ntdll.7C911446	
7C91122E	8D93 78010000	LEA EDX,DWORD PTR DS:[EBX+178]	
7C911234	E9 5AFFFFFF	JMP ntdll.7C911193	
7C911239	81E1 FF000000	AND ECX,0FF	
7C91123F	75 40	JNZ SHORT ntdll.7C911250	

Tabla XOR

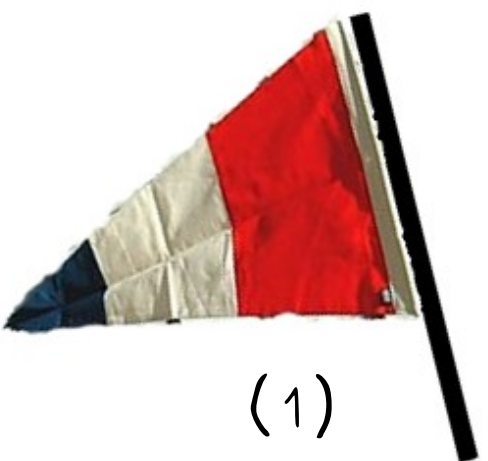
1 xor 1 = 0

1 xor 0 = 1

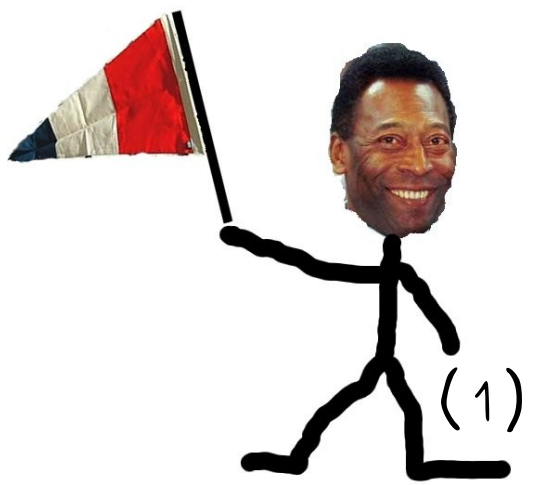
0 xor 1 = 1

0 xor 0 = 0

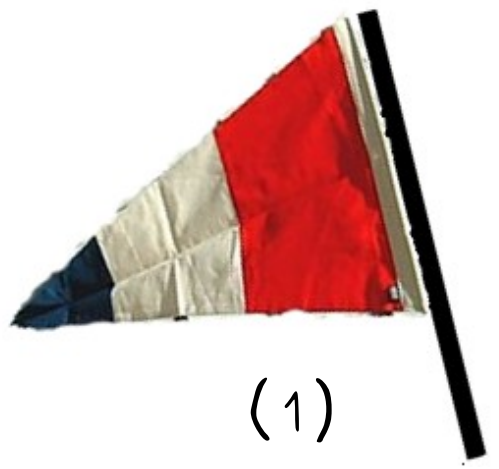
In Pseudo-Code...



XOR



=



XOR



=

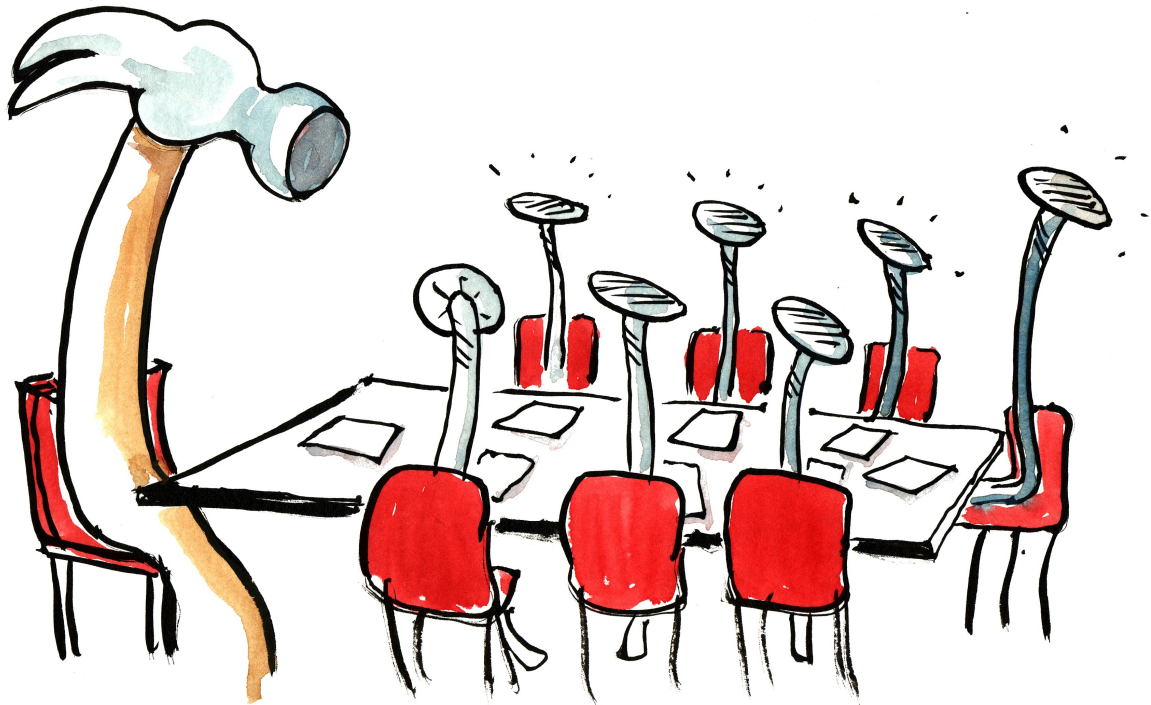


Did you get it?



Exploiting...

the Bitmask



When the only tool
you have is a
hammer...

Exploiting the Bitmask



```
Heap dump 0x00150000
Address Chunks
0x00150220 [015] 0x00150220 -> [ 0x00150220 | 0x00150220 ]
0x00150228 [016] 0x00150228 -> [ 0x0017dc60 | 0x001be040 ]
0x001be040 0x001be040 -> [ 0x00150228 | 0x001a9770 ] (00000016)
0x001a9770 0x001a9770 -> [ 0x001be040 | 0x0017dc60 ] (00000016)
0x0017dc60 0x0017dc60 -> [ 0x001a9770 | 0x00150228 ] (00000016)
0x00150230 [017] 0x00150230 -> [ 0x001caa28 | 0x001caa28 ]
0x001caa28 0x001caa28 -> [ 0x00150230 | 0x00150230 ] (00000017)
0x00150238 [018] 0x00150238 -> [ 0x00150238 | 0x00150238 ]
0x00150240
0x001c5fa0
Heap dump 0x00150000
Address Chunks
0x001c0668
0x00150248 0x001ca9b8 heap: *0x00150000* flags: 0x00000001 (B)
0x00150250 0x001ca9c8 0x001ca9c8> size: 0x00000058 (000b) prevsize: 0x00000010 (0002)
0x0015fd88 0x001ca9c8 heap: *0x00150000* flags: 0x00000001 (B)
0x00150258 0x001caa20 0x001caa20> size: 0x000000b8 (0017) prevsize: 0x00000058 (000b)
0x001bf960 0x001caa20 heap: *0x00150000* flags: 0x00000000 (F)
0x001c2990 0x001caa20 next: 0x00150230 prev: 0x00150230
0x001b85f0
```

Forcing the overflowed chunk to be:

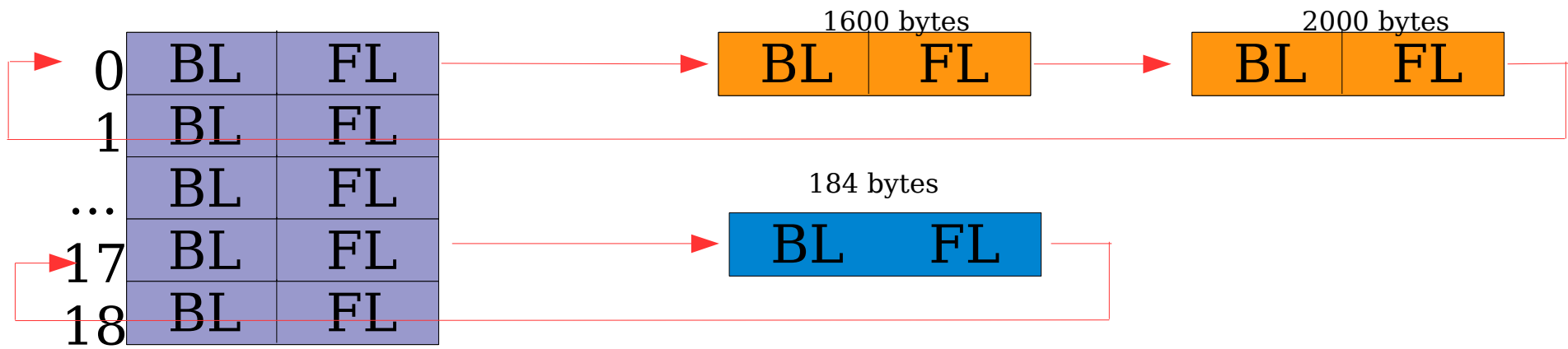
- Free
- Size < 1024
- Only chunk on a FreeList Slot

Exploiting the bitmask

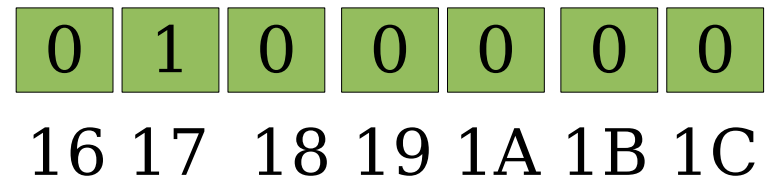


- Overflow **ONLY** the chunk's size with a size that correspond to an empty FreeList slot.

Exploiting the bitmask



Overflow the chunk of size 184 for 992. Chunk continue being connected to slot 17.



Exploiting the bitmask



Exploiting the bitmask



Exploiting the bitmask

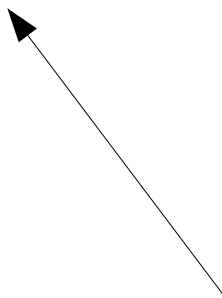


- Force an Allocation of the chunk of the original size.

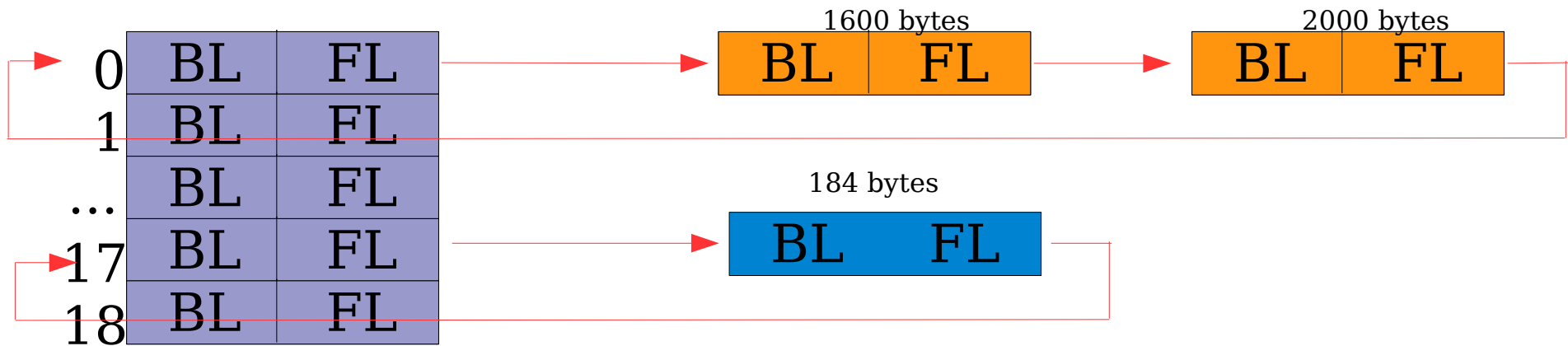
Exploiting the bitmask



Allocate(184)



Exploiting the bitmask



An allocation of 184 bytes will

0	0	0	1	0	0	0
---	---	---	---	---	---	---

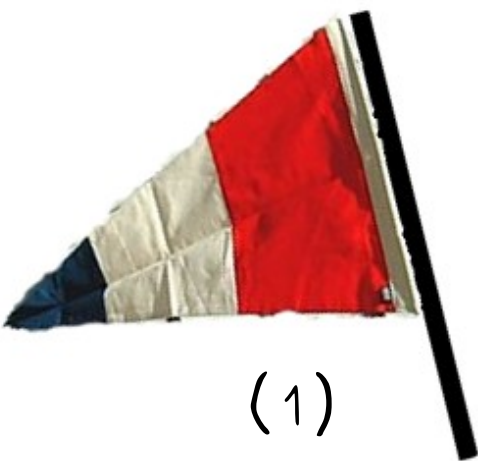
try to take that chunk out

79	7A	7B	7C	7D	7E	7F
----	----	----	----	----	----	----

of the list. But since it's the only (and last) chunk it need to **Unset** the corresponding FreeListInUse, but the opposite happens...

Unset pseudo code: **XOR[7c]=1**

Exploiting the bitmask



XOR



=



Exploiting the bitmask



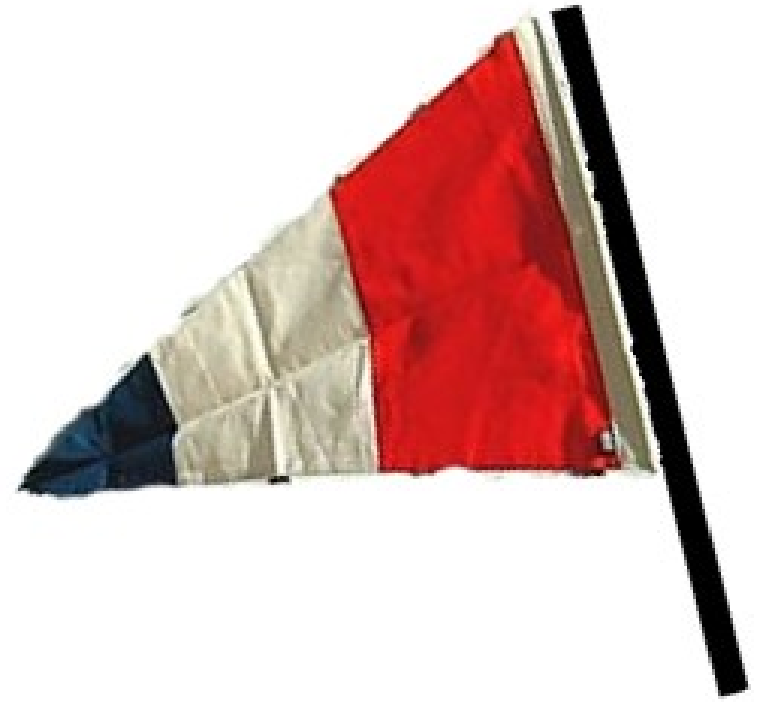
As a result, we had a FreeListInUse slot looking like it has chunks, but it doesn't

Fallum ergo sum



Exploiting the bitmask

- Next time **allocation** of fake size (992), the algorithm will find out that 0x7C seems available (but it's empty).
- An empty slot points to itself (This means, its point to the FreeList)



Exploiting the bitmask

FreeList[7c] points to itself.

An allocation will return a pointer to itself.

What to write?

HEAP + 0x57C =
RtlCommitRoutine

Address	Hex dump	ASCII	
00150558	58 05 15 00	58 05 15 00	X...X..
00150560	60 05 15 00	60 05 15 00	`...`..
00150568	C8 47 18 00	C8 47 18 00	ÈG...ÈG..
00150570	E0 81 15 00	10 FA 17 00	à...ú..
00150578	08 06 15 00	00 00 00 00
00150580	F0 9F 1E 00	00 00 02 05	ø...ø..
00150588	00 00 00 00	00 50 D9 02PÙ..

Address	Chunks
0x00150540	[079] 0x00150540 -> [0x001d6168 0x0015d660]
0x0015d660	0x0015d660 -> [0x00150540 0x001b40a0] (00000079)
0x001b40a0	0x001b40a0 -> [0x0015d660 0x001d6168] (00000079)
0x001d6168	0x001d6168 -> [0x001b40a0 0x00150540] (00000079)
0x00150548	[07a] 0x00150548 -> [0x00150548 0x00150548]
0x00150550	[07b] 0x00150550 -> [0x001bdad8 0x001b3638]
0x001b3638	0x001b3638 -> [0x00150550 0x001bdad8] (0000007b)
0x001bdad8	0x001bdad8 -> [0x001b3638 0x00150550] (0000007b)
0x00150558	[07c] 0x00150558 -> [0x00150558 0x00150558]
0x00150560	[07d] 0x00150560 -> [0x00150560 0x00150560]
0x00150568	[07e] 0x00150568 -> [0x001847c8 0x001847c8]

Neither the Heap
Cookie or ASLR
Affect this technique




Cookie Diet...



At the moment of Allocation, the algorithm doesn't check FreeList chunk's for cookies!

(FTW!)

Size	PSize		...
FLink		Blink	
DATA...			

Two bytes overwrite

+

Two controlled allocations

=

One reliable exploit



Bitmask Alternatives:

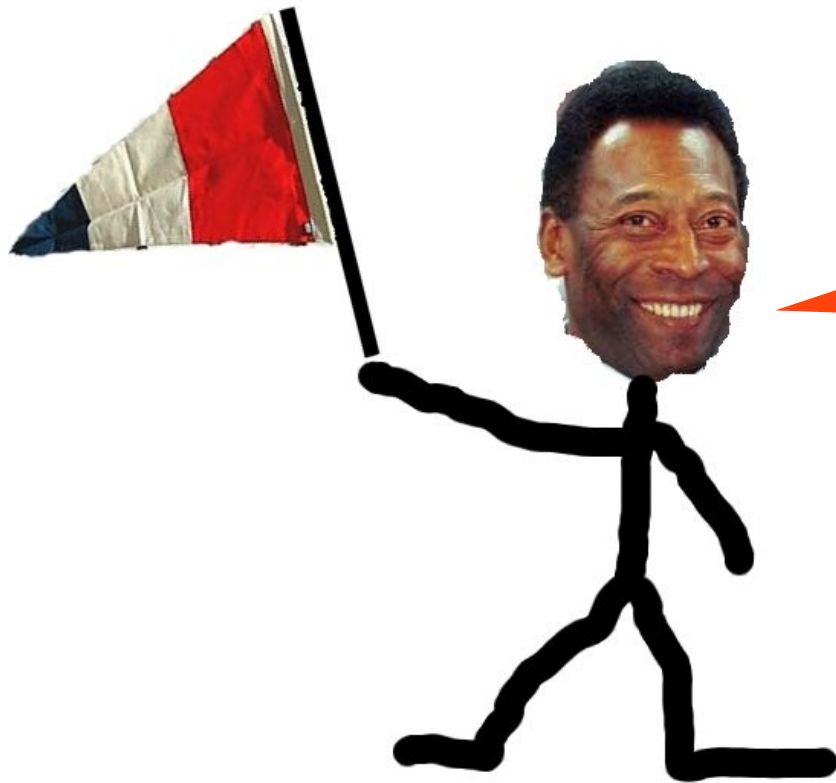
XOR another bit



What if we make the fake size be bigger than 1018 (FreeList size)?

John McDonald y Chris Valasek technique

We will XOR 1 bit after the FreeList.



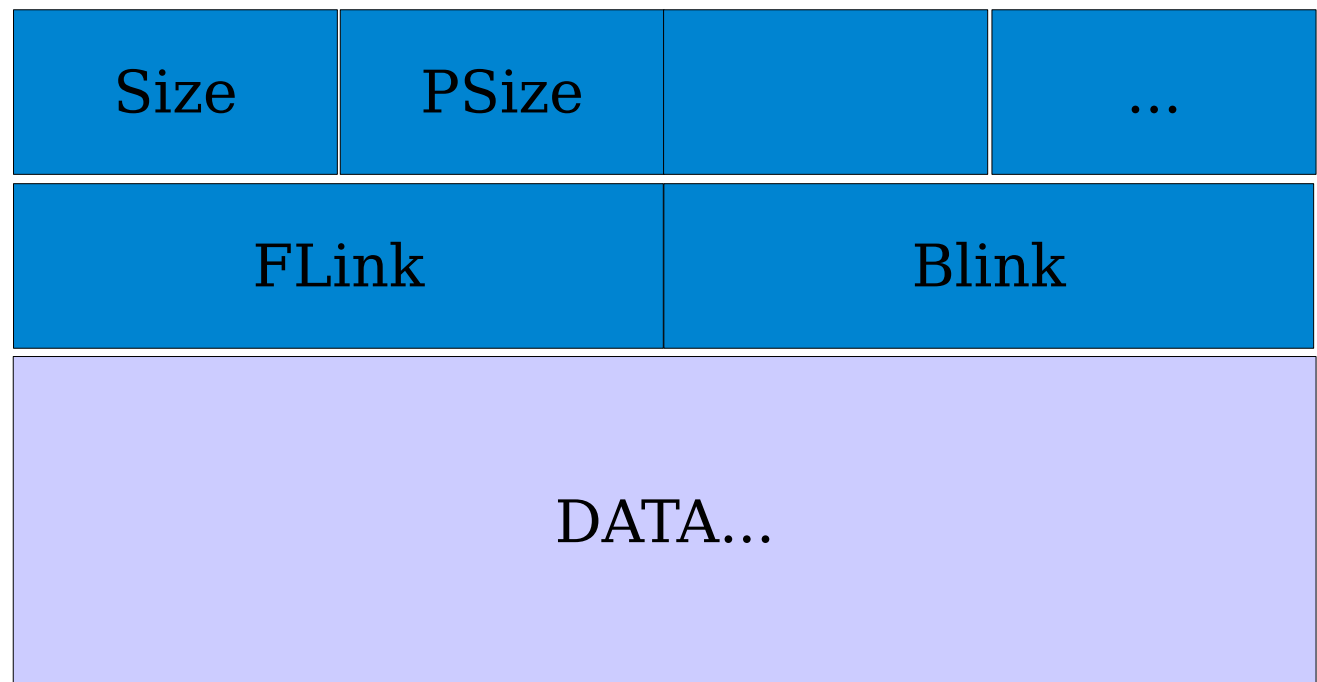
I recommend the
RtlCommitRoutine

Bitmask Alternatives:

16 bytes overwrite



What if instead of two, you can overwrite 16 bytes?



Make:

Forward link = Backward link

Trick:

Bitmask to think you
are the only chunk.



Exploiting...

Demo

Considerations

- Vista and Windows7 protect the chunk's header (per-heap key encryption and checksum).
- As it was explained the technique won't work on Vista/Win7, but its highly possible that one the variations work.



Homework



- Read “*Practical Windows XP / 2003 Heap Exploitation*” by John McDonald y Chris Valasek for more techniques and fun.

A close-up photograph of a small, round, yellow pill with a small indentation on its surface. The pill is resting on a piece of white paper with faint, cursive handwriting in blue ink. The handwriting is out of focus, but some words like "center" and "est" are visible. The lighting is soft, creating a slight shadow to the right of the pill.

Questions?