

# Neurosymbolic Proof Search for Linguistics

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Logic for the AI Spring  
September 2022, Como



A composition calculus for  
vector-based semantic modelling  
with a localization for Dutch

NWO 360-89-070, 2017-2022



**Utrecht  
University**

# The Holy Trinity

Language	Logic	Computation
grammars	substructural logics	$\lambda$ -calculi
empirical linguistic rules	logical inference rules	computation steps
grammaticality	provability	type inhabitation
sentence	proof	program

**Lexicon:** a look-up table from words to *types* and *meanings*

The type-logical perspective:

- ▶ focus on *surface* syntax
- ▶ use LC/NL as core logic (ordered & linear types)
- ▶  $\diamond, \square$  modalities for structural control
- ▶ applicative derivational semantics via proof/term morphisms

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**Lexicon:** a look-up table from words to *types* and *meanings*

A **heretical** perspective:

- ▶ focus on *deep* (discontinuous) syntax
- ▶ use ILL as core logic (linear types)
- ▶  $\diamond, \square$  modalities for dependency annotation
- ▶ extra semantic dimension: adjunct/complement distinction

# Types

ILL<sub>—○</sub> plus  $\diamond, \square$  modalities for *dependency domain demarcation*.

$\mathbb{T}$  inductively defined as:

$$\mathbb{T} := A \mid T \multimap T \mid \diamond^d T \mid \square^d T \quad A \in \mathbb{A}, T \in \mathbb{T}$$

$\mathbb{A}$  – closed set of base types

$\multimap$  – linear function builder

$\diamond$  – reserved for “necessary arguments”, i.e. complements

$\square$  – reserved for “optional arguments”, i.e. subjects

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# Terms

$$\frac{}{c:T \vdash c:T} \text{Lex}$$

$$\frac{\Gamma \vdash s:T_1 \multimap T_2 \quad \Delta \vdash t:T_1}{\Gamma, \Delta \vdash st:T_2} \multimap E$$

$$\frac{\Gamma \vdash t:T}{\langle \Gamma \rangle^d \vdash \Delta^d t : \diamond^d T} \diamond^d I$$

$$\frac{\Gamma \vdash s:\Box^d T}{\langle \Gamma \rangle^d \vdash \nabla^d s:T} \Box^d E$$

$$\frac{}{x:T \vdash x:T} Ax$$

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$$\frac{\Gamma[\langle x:T_1 \rangle^d] \vdash t:T_2 \quad \Delta \vdash s:\diamond^d T_1}{\Gamma[\Delta] \vdash t[x \mapsto \nabla^d s]:T_2} \diamond^d E$$

$$\frac{\langle \Gamma, \langle x:T_1 \rangle^X, \Delta \rangle^d \vdash s:T_2}{\langle \Gamma, \Delta \rangle^d, \langle x:T_1 \rangle^X \vdash s:T_2} X$$

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# Example

“Meloni e Salvini in Piazzale Loreto”

$$\begin{array}{c}
 \frac{c_5 : \text{Loreto}}{c_5 \vdash \Box^{app} (np \multimap np)} \quad \frac{c_4 : \text{Piazzale}}{c_4 \vdash np} \\
 \frac{\langle c_5 \rangle^{app} \vdash np \multimap np}{\langle c_5 \rangle^{app}, c_4 \vdash np} \Box E \quad \frac{c_4 \vdash np}{c_4 \vdash np} \multimap E \\
 \frac{c_3 : \text{in} \quad \langle c_5 \rangle^{app}, c_4 \vdash np}{c_3 \vdash \Diamond^{obj} np \multimap \Box^{mod} (np \multimap np)} \quad \frac{\langle c_5 \rangle^{app}, c_4 \vdash np}{\langle \langle c_5 \rangle^{app}, c_4 \rangle^{obj} \vdash \Diamond^{obj} np} \Diamond I \\
 \frac{c_3, \langle \langle c_5 \rangle^{app}, c_4 \rangle^{obj} \vdash \Box^{mod} (np \multimap np)}{\langle c_3, \langle \langle c_5 \rangle^{app}, c_4 \rangle^{obj} \rangle^{mod} \vdash np \multimap np} \Box E \quad \frac{c_1 : e \quad \frac{c_0 : \text{Meloni}}{c_0 \vdash np} \quad \frac{c_2 : \text{Salvini}}{c_2 \vdash np}}{c_1 \vdash \Diamond^{cnj} np \multimap \Diamond^{cnj} np \multimap np} \Diamond I \quad \frac{c_0 \vdash np \quad c_2 \vdash np}{\langle c_2 \rangle^{cnj} \vdash \Diamond^{cnj} np} \Diamond I \\
 \frac{\langle c_3, \langle \langle c_5 \rangle^{app}, c_4 \rangle^{obj} \rangle^{mod} \vdash np \multimap np \quad c_1 \vdash \Diamond^{cnj} np \multimap \Diamond^{cnj} np \multimap np \quad \langle c_0 \rangle^{cnj} \vdash \Diamond^{cnj} np \quad \langle c_2 \rangle^{cnj} \vdash \Diamond^{cnj} np}{c_1, \langle c_0 \rangle^{cnj}, \langle c_2 \rangle^{cnj} \vdash np} \multimap E \\
 \frac{\langle c_3, \langle \langle c_5 \rangle^{app}, c_4 \rangle^{obj} \rangle^{mod}, c_1, \langle c_0 \rangle^{cnj}, \langle c_2 \rangle^{cnj} \vdash np}{\langle c_3, \langle \langle c_5 \rangle^{app}, c_4 \rangle^{obj} \rangle^{mod}, c_1, \langle c_0 \rangle^{cnj}, \langle c_2 \rangle^{cnj} \vdash np} \multimap E
 \end{array}$$

$\lambda$  read-off:

$$\nabla^{mod} (\text{in} (\nabla^{app} (\text{Loreto}) \text{Piazzale})) (e \Delta^{cnj} (\text{Meloni}) \Delta^{cnj} (\text{Salvini}))$$

**Resource:** æthel ~70 000 proofs (in Dutch)  
see: [abs/1912.12635](https://abs/1912.12635)

# Towards Wide Coverage

## Lexical Types on demand?

Word  $\mapsto$  Type is not 1-1  $\implies$  how to choose right type?

Corpus  $\subset$  language  $\implies$  new words?

Inductive type universe  $\implies$  new types?

## Proof Search?

Non-directional types  $\implies$  all permutations must be considered

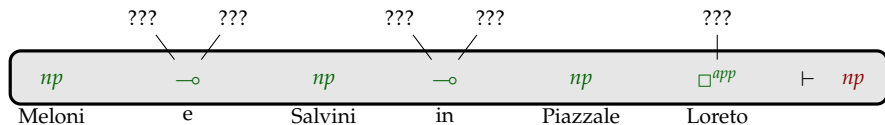
Hypothetical reasoning  $\implies$  Prawitz problem

# Contextual Dynamic Types

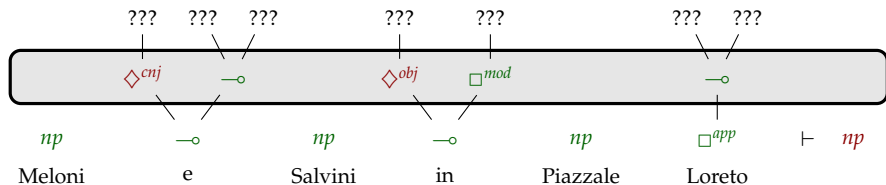
???      ???      ???      ???      ???      ???      ⊢    ???  
Meloni      e      Salvini      in      Piazzale      Loreto



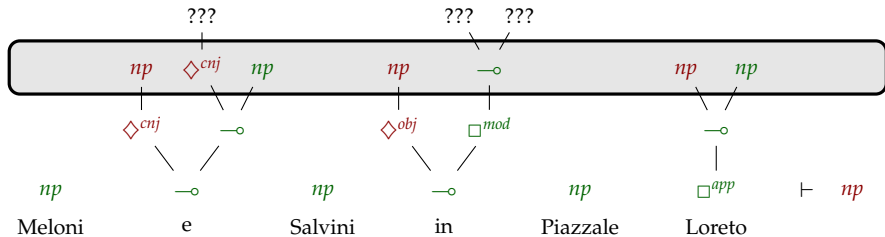
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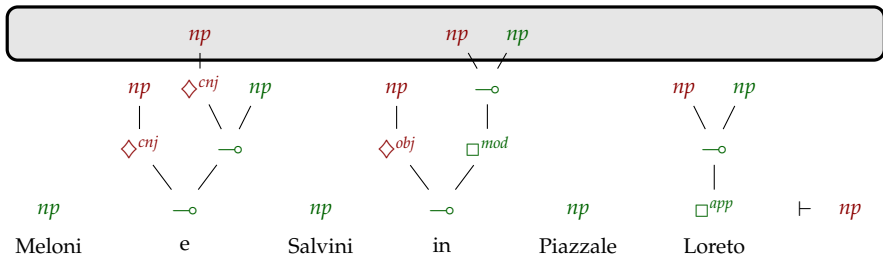
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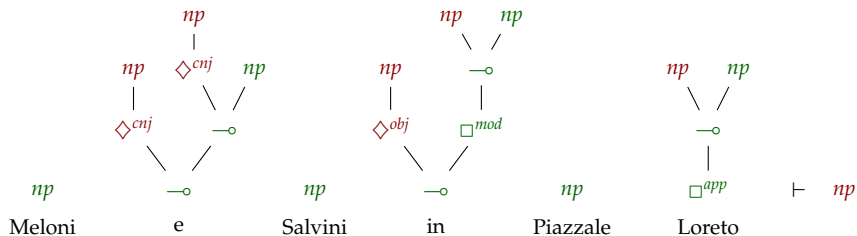
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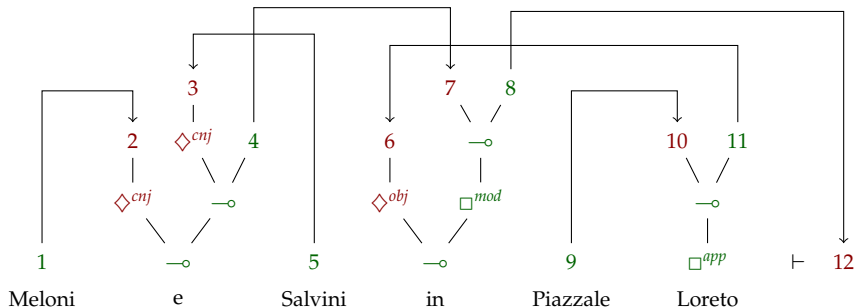


gray box := heterogeneous graph convolution kernel  
see: [abs/2203.12235](https://arxiv.org/abs/2203.12235)

# Proofs as Graphs

## From N.D. to Proof Nets

- ▶ proof structure : tree  $\rightarrow$  graph
- ▶ proof construction : depth induction  $\rightarrow$  parallel assignment
- ▶ proof search : backward/forward chaining  $\rightarrow$  factorial combinatorics

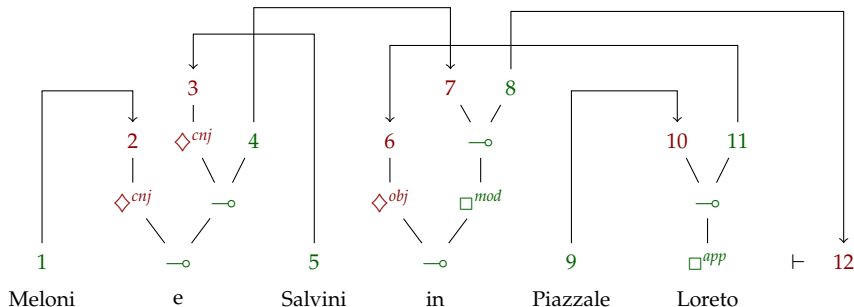


# Proofs as Neural Graphs

## The problem

Find bijection  $\{p : A^+\} \leftrightarrow \{n : A^-\} \forall A \in \mathbb{A}$  in sentence

here  $\{1 \mapsto 2, 4 \mapsto 7, 5 \mapsto 3, 8 \mapsto 12, 9 \mapsto 10, 11 \mapsto 6\}$



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## The solution

compute & normalize  $A^+ \times A^-$ , train against underlying permutation matrix

	2	3	6	7	10	12
1						
4						
5						
8						
9						
11						

see: [abs/2009.12702](https://arxiv.org/abs/2009.12702)



# Final words

## done

- ▶ a logic for dependency control
- ▶ a corpus of linguistically grounded theorems
- ▶ a hyper-efficient neurosymbolic parser

## todo

- ▶ derivational ambiguity?
- ▶ neural integration of proof verification?
- ▶ neural proof generation?

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- ▶ <https://arxiv.org/abs/1808.07447>

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