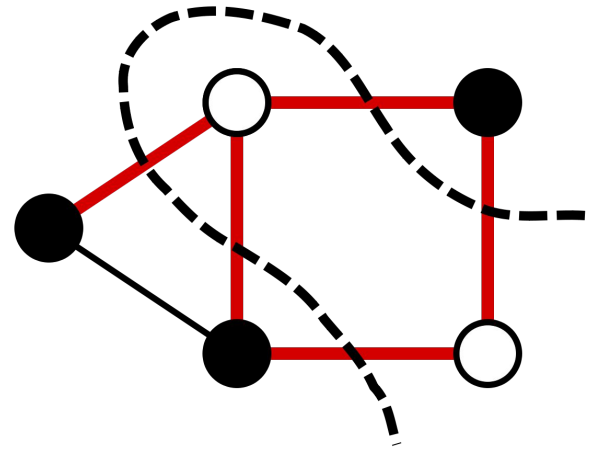


# Parallelized MAXCUT

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# What is MAXCUT?

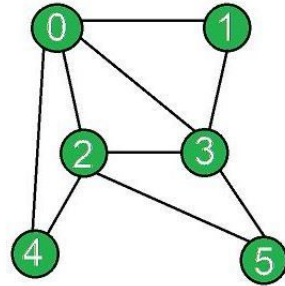
Given a graph  $G(V,E)$ , partition the vertices into  $V_1$  and  $V_2$  s.t. the number of edges between  $V_1$  and  $V_2$  is maximized i.e.



# MAXCUT as an optimization problem

$$\max_{y_i \in \{-1,1\}} \frac{1}{4} \sum_{i,j} w_{ij} (1 - y_i y_j),$$

$$\min_{y_i^2=1} \sum_{i,j} w_{ij} y_i y_j.$$



$w_{ij} = 0$  if  $(i,j)$  is not in  $E$ .

Otherwise it is the weight of edge  $(i,j)$

This is NP-hard.

|   | 0 | 1 | 2 | 3 | 4 | 5 |
|---|---|---|---|---|---|---|
| 0 | 0 | 1 | 1 | 1 | 1 | 0 |
| 1 | 1 | 0 | 0 | 1 | 0 | 0 |
| 2 | 1 | 0 | 0 | 1 | 1 | 1 |
| 3 | 1 | 1 | 1 | 0 | 0 | 1 |
| 4 | 1 | 0 | 1 | 0 | 0 | 0 |
| 5 | 0 | 0 | 1 | 1 | 0 | 0 |

# Relaxation of MAXCUT

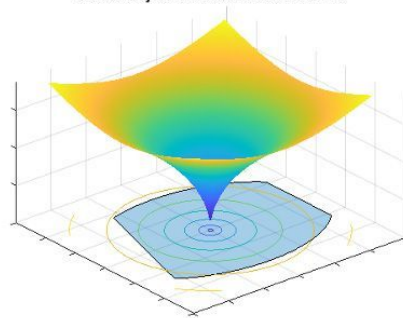
Convex optimization (an SDP):

$$\min_{Y \succeq 0 \in \mathbb{R}^{n \times n}} \text{Tr}(WY) \text{ s.t. } X_{ii} = 1, i = 1, \dots, n$$

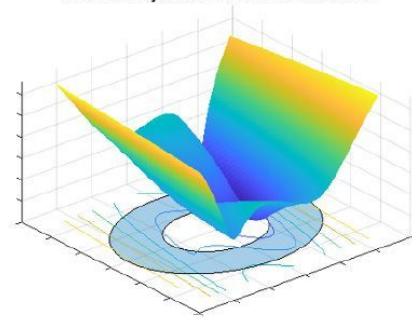
Rank constraint (no longer convex).

$$\min_{V \in \mathbb{R}^{k \times n}} \text{Tr}(WV^T V) \text{ s.t. } \|V_i\|_2 = 1, i = 1, \dots, n$$

Convex Objective and Convex Constraints



Nonconvex Objective and Nonconvex Constraints



# Mixing Method

If all other  $v_j$  are fixed, the last  $v_i$  is optimized with:

$$v_i = \text{normalize}\left(-\sum_j w_{ij}v_j\right)$$

Just keep looping through  $v_i$  until convergence.

# Parallelization

Initial hypothesis: We can loop over all  $v_i$  to find the one that minimizes wrt to the other  $v_j$  in parallel.

# Where I'm currently at

With  $n=30$ ,  $k=4$

SDP solver (way overkill, way too long): 64.722 ms (31442 allocations: 2.31 MiB)

Serial Mixing Method (a lot faster!): 1.756 ms (4762 allocations: 762.50 KiB)

Parallel Mixing Method is not working at the moment: doing the for loop in parallel does updates each vector independently, whereas when it is done serially the vectors actually converge. Instead it finds a  $v$  s.t.  $v^T v = J$ , the matrix of all 1s.

If this can be fixed, I think a  $n$ -times speedup could occur (where  $n$  is the number of processors)

Still working on this! Hoping to try to make even faster.

With  $n=100$ ,  $k=16$

SDP Solver: 1.125 s (311671 allocations: 25.15 MiB)

serial: 62.123 ms (47440 allocations: 19.51 MiB)

It seems the mixing method is still faster, but it may not scale as well with  $n$  as the generic solver? It is still better for smaller graphs.

Could also be an issue with implementation. Mixing method doesn't seem to scale well with  $k$ .