

Intro to Rust Lang Lifetimes



Today: Lifetimes

We've used the term "lifetime" a few times before, and today we're going to explore what exactly it means.

- What is 'a' lifetime?
- How to think about lifetimes
- Other perspectives...

Lifetimes

Lifetimes are all about references, and **nothing** else.

- Informal definition:
Lifetimes provide a way for Rust to validate pointers at compile time
- Formal definition:
Lifetimes are named regions of code that a reference must be valid for
- Remember that references are just pointers with constraints!

Lifetimes vs Generics and Traits

Lifetimes are similar to trait bounds.

- Traits ensure that a generic type has the behavior we want
- Lifetimes ensure that references are valid for as long as we need them to be

Validating References

The main goal of lifetimes is to prevent *dangling references*.

```
fn main() {  
    let r;  
  
    {  
        let x = 5;  
        r = &x;  
    }  
  
    println!("r: {}", r);  
}
```



- What is the issue with this code?

Validating References

```
error[E0597]: `x` does not live long enough
--> src/main.rs:6:13
6 |         r = &x;
  |           ^^ borrowed value does not live long enough
7 |     }
  |     - `x` dropped here while still borrowed
8 |
9 |     println!("r: {}", r);
  |                   - borrow later used here
```

- The value that `r` refers to has gone out of scope before we could use it
- The scope of `r` is "larger" than the scope of `x`

The Borrow Checker

The Rust compiler's borrow checker will compare scopes to determine whether all borrows are valid.

Here is the same code, but with a lifetime diagram:

```
fn main() {  
    let r; // -----+--- 'a  
           // |  
    {     // |  
        let x = 5; // -+-- 'b  
        r = &x; // |  
    }     // -+  
           // |  
    println!("r: {}", r); // |  
}         // -----+  
           //
```



The Borrow Checker

```
fn main() {  
    let r; // -----+--- 'a  
           // |  
    {     // |  
        let x = 5; // -+--- 'b  
        r = &x; // |  
    }     // -+  
           // |  
    println!("r: {}", r); // |  
}         // -----+---
```

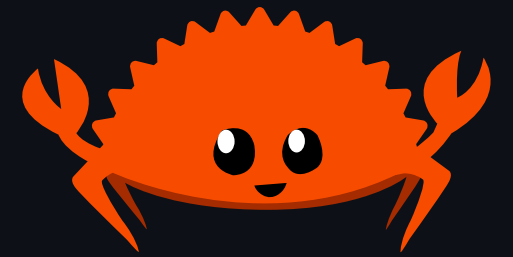


- The borrow checker will compare the "size" of the two lifetimes
 - `r` has a lifetime of `'a`
 - `r` refers to a variable with lifetime `'b`

Placating the Borrow Checker

We can fix this code by removing the scope.

```
fn main() {  
    let x = 5;           // -----+--- 'b  
                        // |  
    let r = &x;        // --+--- 'a |  
                        // | |  
    println!("r: {}", r); // | |  
                        // --+ |  
}                       // -----+---
```



- `x` now "outlives" `r`, so `r` can reference `x`

Generic Lifetimes

Let's try to write some string functions.

```
fn main() {  
    let string1 = String::from("abcd");  
    let string2 = "xyz";  
  
    let result = longest(string1.as_str(), string2);  
    println!("The longest string is {}", result);  
}
```

We want this output:

```
The longest string is abcd
```

- Let's implement `longest` !

longest

Here's a first attempt:

```
fn longest(x: &str, y: &str) -> &str {  
    if x.len() > y.len() {  
        x  
    } else {  
        y  
    }  
}
```

- We don't want to take ownership, so we take `&str` inputs



longest error

Unfortunately, our attempt will not compile:

```
error[E0106]: missing lifetime specifier
--> src/main.rs:9:33
   |
9  | fn longest(x: &str, y: &str) -> &str {
   |          ----      ----      ^ expected named lifetime parameter
   |
   = help: this function's return type contains a borrowed value,
         but the signature does not say whether it is borrowed from `x` or `y`
   help: consider introducing a named lifetime parameter
9  | fn longest<'a>(x: &'a str, y: &'a str) -> &'a str {
   |          +++++      ++          ++          ++
```

longest error

The help text from the compiler error reveals some useful information:

```
= help: this function's return type contains a borrowed value,  
but the signature does not say whether it is borrowed from `x` or `y`
```

- Rust can't figure out if the reference returned refers to `x` or `y`
- In fact, neither do we!

longest error

```
fn longest(x: &str, y: &str) -> &str {  
    if x.len() > y.len() {  
        x  
    } else {  
        y  
    }  
}
```



- We don't know which execution path this code will take
- We also don't know the lifetimes of the input references
- Thus we cannot determine the lifetime we return!
- We will need to *annotate* these references

Lifetime Annotation Syntax

We can annotate lifetimes with generic parameters that start with a `'`, like `'a`.

```
&i32           // a reference
&'a i32       // a reference with an explicit annotated lifetime
&'a mut i32   // a mutable reference with an explicit lifetime

&'hello usize // annotations can be any word or character,
&'world bool  // as long as it starts with a tick (')
```

- Annotations do not change the how long references live, they only describe the relationship between lifetimes of references
 - An annotation by itself has little meaning

longest Lifetimes

Let's return back to our `longest` function.

```
fn longest(x: &str, y: &str) -> &str {  
    if x.len() > y.len() {  
        x  
    } else {  
        y  
    }  
}
```



- What do we want the function signature to express?
- What should the relationship be between the lifetimes of the references?

Longest Lifetimes

What exactly are we returning?

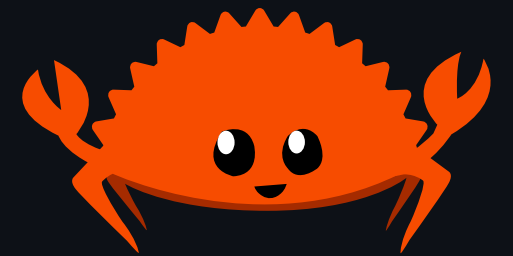
```
if x.len() > y.len() {  
    x  
} else {  
    y  
}
```

- We return either `x` or `y`, which each have their own lifetimes
- We want the returned reference to be valid as long as *both* input references `x` and `y` are valid
- So we want lifetimes of `x` and `y` to *outlive* the returned lifetime

Longest Lifetimes

Since lifetimes are a kind of generic parameter, we must declare them like normal generic type parameters.

```
fn longest<'a>(x: &'a str, y: &'a str) -> &'a str {  
    if x.len() > y.len() {  
        x  
    } else {  
        y  
    }  
}
```



- This will compile now!
- *Remember that these lifetime annotations don't change the lifetimes of any values*

Lifetime Annotations in Functions

We can extrapolate a lot from a function's signature, even without the body.

```
fn longest<'a>(x: &'a str, y: &'a str) -> &'a str;
```

- This function takes two string slices (`x` and `y`) that live at least as long as the lifetime `'a`
- The string slice returned (the longer of `x` or `y`) will also live at least as long as `'a`

Lifetime Annotations in Functions

```
fn longest<'a>(x: &'a str, y: &'a str) -> &'a str;
```

- When calling `longest`, the lifetime that is substituted for `'a` is the intersection of the lifetimes of `x` and `y`
- In practice, this means the lifetime returned by `longest` is the same as the smaller of the two input lifetimes

Borrow Checker Example 1

Let's look at some examples where the borrow checker is and isn't happy.

```
fn main() {  
    let string1 = String::from("long string is long");  
  
    {  
        let string2 = String::from("xyz");  
        let result = longest(string1.as_str(), string2.as_str());  
        println!("The longest string is {}", result);  
    }  
}
```

- `string1` is valid in the outer scope
- `string2` is valid in the inner scope
- `result` should only be valid in the smaller scope (by our lifetime annotations)
 - Since `println!` is in the smaller (inner) scope, this works!

Borrow Checker Example 2

Let's reorder some things around.

```
fn main() {  
    let string1 = String::from("xyz");  
    let result;  
    {  
        let string2 = String::from("long string is long");  
        result = longest(string1.as_str(), string2.as_str());  
    }  
    println!("The longest string is {}", result);  
}
```



- `result` should only be valid in the smaller (inner) scope, but we try to reference it in the outer scope

Borrow Checker Example 2

Sure enough, this does not compile, and Rust gives us this error:

```
error[E0597]: `string2` does not live long enough
--> src/main.rs:6:44
   |
6  |         result = longest(string1.as_str(), string2.as_str());
   |                                                    ^^^^^^^^^^^^^^^^^^^^^^^^^
   |                                                    borrowed value does not live long enough
7  |     }
   |     - `string2` dropped here while still borrowed
8  |     println!("The longest string is {}", result);
   |                                                    ----- borrow later used here
```

Borrow Checker Example 3

What if we knew (as the programmer) that `string1` is always longer than `string2`?

Let's switch the strings around:

```
let string1 = String::from("long string is long");
let result;
{
    let string2 = String::from("xyz");
    result = longest(string1.as_str(), string2.as_str());
}
println!("The longest string is {}", result);
```



Borrow Checker Example 3

```
let string1 = String::from("long string is long");
let result;
{
    let string2 = String::from("xyz");
    result = longest(string1.as_str(), string2.as_str());
}
println!("The longest string is {}", result);
```



- Even though we know (as a human) that the reference will be valid, the compiler does not know

Avoiding Lifetime Annotations

Suppose we wanted to always return the first input, `x`.

```
fn first<'a>(x: &'a str, y: &str) -> &'a str {  
    x  
}
```

- We don't need to annotate `y` with `'a`, because the return value doesn't care about `y`'s lifetime

Lifetimes of Return Values

The lifetime of a return value *must* match the lifetime of one of the inputs.

```
fn dangling<'a>(x: &str, y: &str) -> &'a str {  
    let result = String::from("really long string");  
    result.as_str()  
}
```

- If it didn't depend on an input, then it would *always* be a dangling reference!



Lifetime Elision

All references must have a lifetime. But we've seen many references without lifetime annotations...

This is a version of a function we saw back in week 2:

```
fn first_word(s: &str) -> &str {  
    let bytes = s.as_bytes();  
  
    for (i, &item) in bytes.iter().enumerate() {  
        if item == b' ' {  
            return &s[0..i];  
        }  
    }  
  
    &s[..]  
}
```

Story Time

Long ago, in the dark ages of the 2010s, every reference needed an explicit lifetime.

```
fn first_word<'a>(s: &'a str) -> &'a str {
```

- Before Rust 1.0, every single `&` needed an explicit `'something` annotation
- This became incredibly repetitive, and so the Rust team programmed the borrow checker to infer lifetime annotation patterns of certain situations
- These patterns are called the *lifetime elision rules*

Lifetime Elision

- Lifetime elision does not provide full inference, it will only infer when it is absolutely sure it is correct
- Lifetimes on function or method arguments are called *input lifetimes*, and lifetimes on return values are called *output lifetimes*
- There are only 3 lifetime elision rules, the first for input lifetimes, the last two for output lifetimes

Lifetime Elision Rule 1

The first rule is that the compiler will assign a different lifetime parameter for each input lifetime.

```
fn foo(x: &i32);  
fn foo<'a>(x: &'a i32);  
  
fn bar(x: &i32, y: &i32);  
fn bar<'a, 'b>(x: &'a i32, y: &'b i32);
```

Lifetime Elision Rule 2

The second rule is that if there is only 1 input lifetime parameter, then it is assigned to all output lifetimes.

```
fn foo(x: &i32) -> &i32;  
fn foo<'a>(x: &'a i32) -> &'a i32;  
  
fn bar(arr: &[i32]) -> (&i32, &i32);  
fn bar<'a>(arr: &'a [i32]) -> (&'a i32, &'a i32);
```


Lifetime Elision Rule 3

If there are multiple input lifetime parameters, but the first parameter is `&self` or `&mut self`, the lifetime of `&self` is assigned to all output lifetimes.

- This only applies to methods
- Makes writing methods much nicer!
- *Examples to come later...*

Lifetime Elision Example 1

Let's pretend we are the compiler, and let's attempt to apply the lifetime elision rules to `first_word`.

```
fn first_word(s: &str) -> &str;
```

Lifetime Elision Example 1

We apply the first rule, which specifies that each parameter gets its own lifetime.

```
fn first_word<'a>(s: &'a str) -> &str;
```

Lifetime Elision Example 1

The second rule specifies that the lifetime of the single input parameter gets assigned to all output lifetimes, so the signature becomes this:

```
fn first_word<'a>(s: &'a str) -> &'a str;
```

- Since all references have lifetime annotations, we're done!

Lifetime Elision Example 2

So why didn't elision work with `longest`? Let's trace it out!

We start with this signature without annotations:

```
fn longest(x: &str, y: &str) -> &str;
```

Lifetime Elision Example 2

Let's apply the first rule and get annotations for all inputs.

```
fn longest<'a, 'b>(x: &'a str, y: &'b str) -> &str;
```

- What now?

Lifetime Elision Example 2

```
fn longest<'a, 'b>(x: &'a str, y: &'b str) -> &str;
```

- The second rule doesn't apply here, because there is more than 1 input lifetime ('a and 'b)
- Since Rust cannot figure out what to do, it gives a compiler error to the programmer so they can write the annotations themselves

Lifetimes in Structs

So far, all of the `struct`s we've looked at have held *owned* type fields.

If we want a `struct` to hold a reference, we need to annotate them.

```
struct ImportantExcerpt<'a> {
    part: &'a str,
    importance: i32,
}

fn main() {
    let novel = String::from("Call me Ishmael. Some years ago...");
    let first_sentence = novel.split('.').next().expect("Could not find a '.'");
    let i = ImportantExcerpt {
        part: first_sentence,
        importance: 42,
    };
}
```


Lifetimes in Structs

```
struct ImportantExcerpt<'a> {  
    part: &'a str,  
    importance: i32,  
}
```

- As with generic data types, we declare the name of the generic lifetime parameter inside angle brackets
- This annotation means an instance of `ImportantExcerpt` can't outlive the reference it holds in its `part` field

Lifetimes in `impl` Blocks

Similarly, we need to annotate `impl` blocks with lifetime parameters.

```
impl<'a> ImportantExcerpt<'a> {  
    fn importance(&self) -> i32 {  
        self.importance  
    }  
}
```

Lifetimes in Methods

Here is an example where the third elision rule is applied:

```
impl<'a> ImportantExcerpt<'a> {  
    fn announce_and_return_part(&self, announcement: &str) -> &str {  
        println!("Attention please: {}", announcement);  
        self.part  
    }  
}
```

- The first rule gives both `&self` and `announcement` their own lifetimes
- The third rule gives the return lifetime the lifetime of `&self`

Putting it all together...

Let's briefly look at the syntax of specifying generic type parameters, trait bounds, and lifetimes all in one function!

```
fn longest_with_an_announcement<'a, T>(x: &'a str, y: &'a str, ann: T) -> &'a str
where
    T: Display,
{
    println!("Announcement! {}", ann);

    if x.len() > y.len() {
        x
    } else {
        y
    }
}
```

Lifetime Bounds

Lifetimes can be bounds, just like traits.

```
#[derive(Debug)]  
struct Ref<'a, T: 'a>(&'a T);
```

- `Ref` contains a reference, with a lifetime of `'a`, to a generic type `T`
- `T` is bounded such that any references *in* `T` must live at least as long as `'a`
- Additionally, the lifetime of `Ref` may not exceed `'a`

Lifetime Bounds

Here is a similar example, but with a function instead of a `struct`.

```
fn print_ref<'a, T>(t: &'a T)
where
    T: Debug + 'a,
{
    println!("print_ref(t) is {:?}", t);
}
```

- `T` must implement `Debug`, and all references *in* `T` must outlive `'a`
- Additionally, `'a` must outlive this function call

Lifetime Bounds

Putting the `Ref` and `print_ref` together:

```
#[derive(Debug)]
struct Ref<'a, T: 'a>(&'a T);

fn print_ref<'a, T>(t: &'a T)
where
    T: Debug + 'a,
{
    println!("print_ref(t) is {:?}", t);
}

fn main() {
    let x = vec![9, 8, 0, 0, 8];
    let ref_x = Ref(&x);
    print_ref(&ref_x);
    // Prints to stdout: print_ref(t) is Ref([9, 8, 0, 0, 8])
}
```

Lifetime-bounded Lifetimes

We can have lifetimes that are bounded by other lifetimes.

```
// Takes in a `&'a i32` and return a `&'b i32` as a result of coercion
fn choose_first<'a: 'b, 'b>(first: &'a i32, _: &'b i32) -> &'b i32 {
    first
}

fn main() {
    let first = 2; // Longer lifetime
    {
        let second = 3; // Shorter lifetime
        println!("{}", choose_first(&first, &second));
    }
}
```

- `'a: 'b` reads as "lifetime `'a` outlives `'b`"

The 'static' Lifetime

There is a special lifetime called 'static'.

```
let s: &'static str = "I have a static lifetime";
```

- 'static implies that the reference will live until the end of the program (it is valid until the program stops running)
- Here, `s` is stored in the program binary, so it will always be valid!

'static' Error Messages

You may see suggestions to use the 'static' lifetime in error messages.

```
fn foo() -> &i32 {  
    let x = 5;  
    &x  
}
```

```
help: consider using the `static` lifetime, but this is uncommon unless you're  
      returning a borrowed value from a `const` or a `static`
```

```
2 | fn foo() -> &'static i32 {  
  |           ++++++
```

- Before making a change, think about if your reference will *really* live until the end of the program
- You may actually be trying to create a dangling reference!

'static' vs static

There are two common ways to make a variable with a 'static' lifetime.

1. Make a string literal with has type `&'static str`
2. Make a constant with the `static` declaration

'static vs static Example

```
static NUM: i32 = 42;
static NUM_REF: &'static i32 = &NUM;

fn main() {
    let msg: &'static str = "Hello World";
    println!("{msg} {NUM_REF}!");
}
```

Hello World 42!

'static Memory Leaks

There is a third way: we can create 'static values by *leaking memory*.

```
fn random_vec() -> &'static [usize; 100] {
    let mut rng = rand::thread_rng();
    let mut boxed = Box::new([0; 100]);
    boxed.try_fill(&mut rng).unwrap();
    Box::leak(boxed)
}

fn main() {
    let first: &'static [usize; 100] = random_vec();
    let second: &'static [usize; 100] = random_vec();
    assert_ne!(first, second)
}
```

- This allows us to *dynamically* create a 'static reference!

The 'static Bound

'static can also be used as a type bound. However...

- There is a subtle difference between the 'static lifetime and the 'static bound
- The 'static bound means that the type does not contain any non-static references
- This means that all owned data implicitly has a 'static bound, since owned data holds no references

'static Bound Example

Here's an example of using a 'static bound.

```
fn print_it(input: impl Debug + 'static) {  
    println!("'static value passed in is: {:?}", input);  
}  
  
fn main() {  
    // i is owned and contains no references,  
    // thus it has a 'static bound  
    let i = 5;  
    print_it(i);  
  
    // oops, &i only has the lifetime defined by  
    // the scope of main, so it's not 'static  
    print_it(&i);  
}
```



'static' Bound Example

We get a compiler error:

```
error[E0597]: `i` does not live long enough
--> src/lib.rs:15:15
15 |         print_it(&i);
   |                   ^^
   |                   |
   |                   borrowed value does not live long enough
   |                   argument requires that `i` is borrowed for `'static`
16 |     }
   |     - `i` dropped here while still borrowed
```


Review

- Rust has lifetimes to prevent dangling references
- The borrow checker will ensure that lifetimes are always valid
- Rust will allow you elide lifetime annotations in some situations

Further Reading

- You can find some more examples here: [Rust By Example](#)
- If you want to go *really* in depth, read the Rustonomicon chapter on [lifetimes](#)

Another Perspective

What is 'a lifetime'?

- This is a great video made by `leddoo` that explains another way to think about lifetimes!
- Instead of lifetimes as regions of code or scopes, what if we thought about lifetimes as regions of memory?
- Let's watch it together!

Watch Party

What is 'a lifetime?

What is 'a lifetime?

Some quick points:

- Thinking about lifetimes as regions of code can be confusing
- Instead, think about lifetimes as regions of valid memory
- Both interpretations are valid!

Feedback

Please fill out the [feedback form](#) (on Piazza).

- It will help us make this semester better for you
- It will also help make future offerings of this course better for others!
- Feedback is anonymous, so please be honest
- You will receive a homework's worth of extra credit!

Next Lecture: Smart Pointers and Trait Objects

Thanks for coming!

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