Challenge description and awards Winners attack in short



# $\begin{array}{c} {\rm CHES} \ 2023 \ {\rm challenge} \\ {\rm SMAesH} \ {\rm challenge} \end{array}$

Gaëtan Cassiers Charles Momin François-Xavier Standaert
SIMPLE-Crypto







UCLouvain



G. Cassiers, C. Momin

#### Content

#### Challenge description and awards

Winners attack in short

#### In the CHES2022 Rump session episode



#### CHES Challenge

#### **Call for CHES Challenge Organizers**

Since 2015, a crypto engineering challenge is organized every year in cooperation with CHES. Former editions have focused on practical side-channel attacks, design of countermeasures, deep learning-based attacks, white-box cryptography, and hardware socurity.

#### See for instance

- https://whitex.io/contests
- https://hackatevent.org/hackches21/
- https://ctf.spook.dev/
- https://chesctf.riscure.com/2018/news

#### SIMPLE-Crypto Association

Open Source Secure Implementation of Cryptographic Algorithms

#### Concretely...

#### **Current stage**

- Higher-order masked AES in hardware (soon a CTF?)
- · SCALib: side-channel evaluation library,

SMAesH challenge

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#### SMAesH you said?

- Masked AES HW IP
- ► HPC2 (arbitrary order)
- Provably secure
- ► PRNG included

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SIMPLE-Crypto's Masked AES in Hardware (SMAesH	Packages     No packages published     Pusion your test peckage

amorein Morris Charles

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Figure 9: Global architecture of the MSKass.32bits.stats.stats.statsathmodule. The value bold by the DFF at index i is depicted by the signal sh.reg.out[1] in the HDL.

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#### Masked AES HW IP

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#### SMAesH: technical documentation

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#### SMAesH: preliminary evaluation report

	SIMPLE-Crypto
с	ontents
1	Overview
2	History
3	Evaluation scope
4	Measurement Setup and Traces Pre-processing
5	Evaluation Methodology
6	Results
7	Conclusion
8	Copyright
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This document presents the findings of the preliminary evaluation of the resistance of the SMAesH (ass.anc128.32bits.hpc2) hardware IP to power analysis attacks. The evaluation has been performed the file developers of SMAesH (SIMPLE-Cypto). The terminology for this report is defined in the SMAesH technical dormeenta-

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 $\rightarrow$  See simple-crypto.org/activities/smaesh/

## Challenge description



Goal: SCA attack against SMAesH (first order)

## Challenge description

#### 420.0M SMAesH IP 410.0M CW305 400.0M Artix7 390.0M 370.0M 360.0M 350.0M 340.0M 330.0M 320.0M Mor 1 Apr 3 Apr 21 Aug 3 Aug 25 $< 2^{68}$ $(1 \text{ BTC-H} \cdot s)$ Test (private) Key rank Evaluation Success (Fixed key)

Goal: SCA attack against SMAesH (first order)

## Challenge description



#### Goal: SCA attack against SMAesH (first order)

Challenge description and awards Winners attack in short



#### ▶ 5 teams: $\infty \times$ '20 CTF :D

#### Challenge Stats

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- ▶ 112 submissions

Challenge Stats

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- ► A7\_d2:
  - ightarrow 77 submissions
  - ightarrow 5 successful attacks
- ► S6\_d2:
  - ightarrow 35 submissions
  - ightarrow 6 successful attacks



Challenge Stats

- ▶ 5 teams:  $\infty \times$  '20 CTF :D
- 112 submissions
- ► A7\_d2:
  - ightarrow 77 submissions
  - ightarrow 5 successful attacks
- ► S6\_d2:
  - ightarrow 35 submissions
  - ightarrow 6 successful attacks
- Peak rates:
  - 2 submissions/h/team
  - 12 submissions/day



#### These damn ninjas cutting onions...



#### Winners

#### Prizes<sup>a</sup> for most points and best attack!

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- Prizes<sup>a</sup> for most points and best attack!
- ► A7\_d2:
  - Most points: Valence Cristiani (team mAEStro)
  - Best attack: Valence Cristiani (team mAEStro)
- ► S6\_d2:
  - Most points: Valence Cristiani (team mAEStro)
  - Best attack: Thomas Marquet (team Morningstar)

<sup>a</sup>Teams cannot win more than one prize...

- Valence Cristiani (NinjaLab) is awarded 1000 USD
- Thomas Marquet (AAU) is awarded 500 USD

#### Congratulations!

#### Content

#### Challenge description and awards

Winners attack in short



Valence Cristiani | Ches 2023

NinjaLab



Sbox

T1 = U0 + U3	T8 = U7 + T6	T15 = T5 + T11	T22 = T7 + T21
T2 = U0 + U5	T9 = U7 + T7	T16 = T5 + T12	T23 = T2 + T22
T3 = U0 + U6	T10 = T6 + T7	T17 = T9 + T16	T24 = T2 + T10
T4 = U3 + U5	T11 = U1 + U5	T18 = U3 + U7	T25 = T20 + T17
T5 = U4 + U6	T12 = U2 + U5	T19 = T7 + T18	T26 = T3 + T16
T6 = T1 + T5	T13 = T3 + T4	T20 = T1 + T19	T27 = T1 + T12
T7 = U1 + U2	T14 = T6 + T11	T21 = U6 + U7	



T23 = U0 + U3	T19 = T22 + R5	T17 = U2 # T19	T6 = T22 + R17
T22 = U1 # U3	T9 = U7 # T1	T20 = T24 + R13	T16 = R13 + R19
T2 = U0 # U1	T10 = T2 + T24	T4 = U4 + T8	T27 = T1 + R18
T1 = U3 + U4	T13 = T2 + R5	R17 = U2 # U5	T15 = T10 + T27
T24 = U4 # U7	T3 = T1 + R5	R18 = U5 # U6	T14 = T10 + R18
R5 = U6 + U7	T25 = U2 # T1	R19 = U2 # U4	T26 = T3 + T16
T8 = U1 # T23	R13 = U1 + U6	Y5 = U0 + R17	

Figure 6: Top linear transform in reverse direction.

 $Z = K \bigoplus P$ 

M1 = T13 x T6	M17 = M5 + T24	M33 = M27 + M25	$M49 = M43 \times T16$
M2 = T23 x T8	M18 = M8 + M7	$M34 = M21 \times M22$	M50 = M38 x T9
M3 = T14 + M1	M19 = M10 + M15	$M35 = M24 \times M34$	M51 = M37 x T17
M4 = T19 x D	M20 = M16 + M13	M36 = M24 + M25	$M52 = M42 \times T15$
M5 = M4 + M1	M21 = M17 + M15	M37 = M21 + M29	M53 = M45 x T27
M6 = T3 x T16	M22 = M18 + M13	M38 = M32 + M33	$M54 = M41 \times T10$
M7 = T22 x T9	M23 = M19 + T25	M39 = M23 + M30	$M55 = M44 \times T13$
M8 = T26 + M6	M24 = M22 + M23	M40 = M35 + M36	$M56 = M40 \times T23$
$M9 = T20 \times T17$	$M25 = M22 \times M20$	M41 = M38 + M40	$M57 = M39 \times T19$
M10 = M9 + M6	M26 = M21 + M25	M42 = M37 + M39	$M58 = M43 \times T3$
$M11 = T1 \times T15$	M27 = M20 + M21	M43 = M37 + M38	$M59 = M38 \times T22$
$M12 = T4 \times T27$	M28 = M23 + M25	M44 = M39 + M40	$M60 = M37 \times T20$
M13 = M12 + M11	$M29 = M28 \times M27$	M45 = M42 + M41	$M61 = M42 \times T1$
$M14 = T2 \times T10$	$M30 = M26 \times M24$	$M46 = M44 \times T6$	$M62 = M45 \times T4$
M15 = M14 + M11	$M31 = M20 \times M23$	$M47 = M40 \times T8$	$M63 = M41 \times T2$
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Sbox tower fileds implementation



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T7 = U1 + U2	T14 = T6 + T11	T21 = U6 + U7	

Figure 5: Top linear transform in forward direction.

T23 = U0 + U3	T19 = T22 + R5	T17 = U2 # T19	T6 = T22 + R17
T22 = U1 # U3	T9 = U7 # T1	T20 = T24 + R13	T16 = R13 + R19
T2 = U0 # U1	T10 = T2 + T24	T4 = U4 + T8	T27 = T1 + R18
T1 = U3 + U4	T13 = T2 + R5	R17 = U2 # U5	$T15 = T10 + T2^{-1}$
T24 = U4 # U7	T3 = T1 + R5	R18 = U5 # U6	T14 = T10 + R18
R5 = U6 + U7	T25 = U2 # T1	R19 = U2 # U4	T26 = T3 + T16
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Build the huge and horrible graph from the equations



Make more than 4000 Gaussian templates (2 for each node since it's masked)





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Build the huge and horrible graph from the equations



Make more than 4000 Gaussian templates (2 for each node since it's masked)

290k traces

Apply belief propagation algorithm (SASCA) and get the key

# But it...

 Requires to understand a lot of theory (graphs, BP algorithm, dealing with the loops etc...)

 $\succ$  Is very long

Does not even guarantee to win

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 Requires to understand a lot of theory (graphs, BP algorithm, dealing with the loops etc...)

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Does not even guarantee to win



## **BEING A BAD GUY**

Let's use another side-channel ? Power leakage is so old shcool...



## Upper bound of log<sub>2</sub>(KeyRank)

Aggregating many well-crafted submissions may allow to extract enough information on the key



# **BEING AN BAD GUY**

# How many submissions?

- Uniform probability for all bytes except one
- Return a different score for each of the 256 values with a uniform spacing (ex: 1, 2 ..., 256)

Obfuscate this behind a neural network...

Upload the submission and store the log<sub>2</sub>(KeyRank)

Avergage of **4.9** bits of information per submissions

4.9 x 13 = 63.7

Require 13 submissions !

## **BEING A BAD GUY**

Read it backwards...

## I created a new account named Sec-artorez

Hawai	A7_d2	200000	×	128.0
Everest	A7_d2	210000	×	126.7
Dubai	A7_d2	220000	×	123.8
Inazawa	A7_d2	225000	×	127.7
Bahamas	A7_d2	215000	×	127.8
Zanzibar	A7_d2	200000	×	127.0
Antarctica	A7_d2	180000	×	127.3
Capri	A7_d2	205000	×	128.0
Faliraki	A7_d2	220000	×	125.2
Gaios	A7_d2	180000	×	127.9
Jakarta	A7_d2	189000	×	125.0
Kuala Lumpur	A7_d2	230000	×	123.3

- First letter is a reminder for the concerned byte
- Space the submission by ~ 2 days...

Local analysis reveals that the we gained 66.1 bits. Means that we should have :

 $\log_2(KeyRank) = 61.9$ 

Aggregate the results and mount the final attack.

# **BEING A BAD GUY**



## The SMAesH challenge has been SMASHED





#### SMAesH Challenge : Or how to enjoy your summer

Thomas Marquet

September 11, 2023





#### Spartan-6 dataset

- Hardware masked AES with two shares (r and  $x \oplus r$ )
- No access to r
- Perfectly synchronized traces
- Low SnR
- Problem : How to pick up enough signal ?
- Solution : Praying to the deep learning god



#### Intermediates under attack

$\texttt{bytes\_to\_SB}[8d-1:0]$	_	0	${\rm RK}_{13}^{i-1}$	$S_0^i$
$\texttt{bytes\_to\_SB}[16d-1:8d]$	_	0	${\tt RK}_{14}^{i-1}$	$S_5^i$
$\texttt{bytes\_to\_SB}[24d-1:16d]$	_	0	${ m RK}_{15}^{i-1}$	$S_{10}^i$
$\verb+bytes_to_SB[32d-1:24d]$	_	0	${\tt RK}_{12}^{i-1}$	$S_{15}^i$



Strategy :

- Recover  $S_{12}, S_1, S_6, S_{11}$
- Recover  $S_i$  from  $S_i \oplus S_{i+4 \pmod{16}}$



S

si si

S<sup>i</sup><sub>7</sub> S<sup>i</sup><sub>1</sub>

Deep learning when randomness isn't known Single-task

With  $m_{\theta_{x_i}}$  the set of layers expected to fit the intermediate  $x_i$ 



(a) A model that do not work (most of the time)

(b) A model that do work (sometimes)

Figure: Hard encoding of the masking scheme inside the network

Dont Learn What You Already Know : Masure and the gangesta



# Simply better model I swear it's not that ugly

$$l_{c_0} = \text{clk 3 to 11}$$
,  $l_{c_1} = \text{clk 4 to 12}$ ,  $l_{c_2} = \text{clk 5 to 13}$   
 $x_i = S_i \oplus S_{i+4 \pmod{16}}$ 



Figure: Multi-task model to recover the transitions  $x_i$ 

Eprint : A Comparison of Multi-task and Single-task Approaches



#### Conclusion

- It leaks less than ASCAD
- Cross entropy go from 5.5452 to 5.5452
- Sun light is overrated

#### Acknowledgments

 Supported by the European Research Council (ERC) under the European Union's Horizon 2020 research and innovation program (grant agreement No 725042)



## What's next?

- Secret fixed key datasets will be published.
- Leaderboard will be updated with SOTA attack.
- SMAesH public evaluation continues...
- ► ... And more are coming!
  → What are you waiting for?



## SIMPLE-crypto

Interested? Want to participate? Question or suggestion?

► SIMPLE-crypto website

```
https://www.simple-crypto.org/
```

SMAesH challenge website

https://smaesh-challenge.simple-crypto.org/

Contact

```
info@simple-crypto.org
```

(or with a beer now ;) )

## THANKS