

MS&E 233

Game Theory, Data Science and AI

Lecture 16

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Management Science and Engineering

(by courtesy) Computer Science and Electrical Engineering

Institute for Computational and Mathematical Engineering

Computational Game Theory for Complex Games

- 1
 - Basics of game theory and zero-sum games (T)
 - Basics of online learning theory (T)
 - Solving zero-sum games via online learning (T)
 - *HW1: implement simple algorithms to solve zero-sum games*
 - Applications to ML and AI (T+A)
 - *HW2: implement boosting as solving a zero-sum game*

- 2
 - Basics of extensive-form games
 - Solving extensive-form games via online learning (T)
 - *HW3: implement agents to solve very simple variants of poker*

- 3
 - General games, equilibria and online learning (T)
 - Online learning in general games
 - *HW4: implement no-regret algorithms that converge to correlated equilibria in general games*

Data Science for Auctions and Mechanisms

- 4
 - Basics and applications of auction theory (T+A)
 - Basic Auctions and Learning to bid in auctions (T)
 - *HW5: implement bandit algorithms to bid in ad auctions*

- 5
 - Optimal auctions and mechanisms (T)
 - Simple vs optimal mechanisms (T)
 - *HW6: implement simple and optimal auctions, analyze revenue empirically*

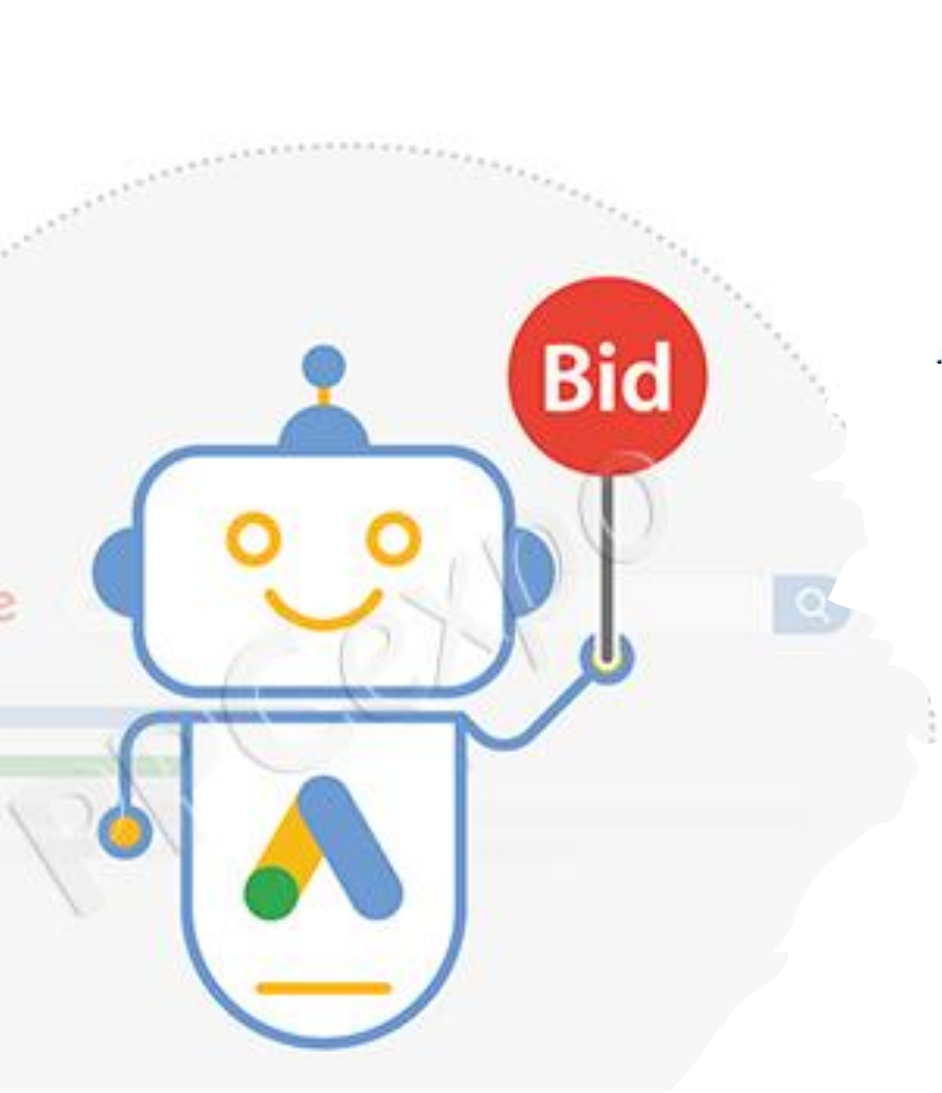
- 6
 - Basics of Statistical Learning Theory (T)
 - Optimizing Mechanisms from Samples (T)
 - *HW7: implement procedures to learn approximately optimal auctions from historical samples*

Further Topics

- 7
 - Econometrics in games and auctions (T+A)
 - **A/B testing in markets (T+A)**
 - *HW8: implement procedure to estimate values from bids in an auction*

Guest Lectures

- Mechanism Design for LLMs, Renato Paes Leme, Google Research
- Auto-bidding in Sponsored Search Auctions, Kshipra Bhawalkar, Google Research



Econometrics in Games and Auctions

Econometrics in Games and Auctions

- We are given data from actions of players in a game (and potentially auxiliary contextual information about the game)
- Multiple instances were players played the same type of game
- We don't know the exact utilities of the players in the game
- We want to use the data to learn the parameters of the utilities of the players in the game or the distribution of these parameters

Why useful?

Scientific: economically meaningful quantities

Perform counter-factual analysis: what would happen if we change the game?

Performance measures: welfare, revenue

Testing game-theoretic models: if theory on estimated quantities predicts different behavior, then in trouble

If I know the equilibrium bid distribution G , then whenever I see a bid b_i , I can *reverse engineer* and *uniquely determine the value* that led to such a bid

unobserved value

$$v_i = b_i + \frac{1}{(n-1) \frac{g(b_i)}{G(b_i)}}$$

observed equilibrium bid

Reverse hazard ratio of distribution of bids
 “Probability that opponent bid is immediately below b_i given that it is below b_i ”

More competition \Rightarrow less “value reduction”

Side Note (Asymmetric Bidders): If I know the equilibrium bid distributions G_i , then whenever I see a bid b_i , I can *reverse engineer* and *uniquely determine the value v_i* that led to such a bid

$$\begin{array}{c} \text{unobserved} \\ \text{value} \end{array} \quad v_i = \begin{array}{c} \text{observed} \\ \text{equilibrium} \\ \text{bid} \end{array} b_i + \frac{1}{\sum_{k \neq i} \frac{g_k(b_i)}{G_k(b_i)}}$$

Reverse hazard ratio
 of distribution of bids
 of k -th opponent
 “Probability that opponent
 bid is immediately below
 b_i given that it is below b_i ”

More competition \Rightarrow less “value reduction”

Estimating CDFs from Truthful Samples

Given truthful bids v_1, \dots, v_m of players in instances of Second Price Auction the CDF of the distribution can be approximated by the empirical CDF to an error of $\approx \frac{1}{\sqrt{n}}$

$$F(z) \stackrel{\text{def}}{=} \Pr(v < z) \approx \frac{1}{n \cdot m} \sum_{i,j} 1\{v_{ij} < z\} \stackrel{\text{def}}{=} \hat{F}(z)$$

Estimating CDFs and PDFs of Bids from FPA Bid Samples

Given bids b_1, \dots, b_m of players in instances of First Price Auction the CDF and PDF of the *bid distribution* can be approximated by empirical CDF and a Kernel Density Estimate

$$G(z) \stackrel{\text{def}}{=} \Pr(b < z) \approx \frac{1}{n \cdot m} \sum_{i,j} 1\{b_{ij} < z\} \stackrel{\text{def}}{=} \hat{G}(z)$$

$$g(z) = \partial_z G(z),$$

$$\hat{g}(z) = \frac{1}{n \cdot m} \sum_{i,j} \frac{1}{h_n} K\left(\frac{b_{ij} - z}{h_n}\right)$$

Fraction of samples that \approx lie within h
from z , divided by region length

Estimating CDFs and PDFs of Values from FPA Bid Samples

Given bids b_1, \dots, b_m of players in instances of First Price Auction the CDF and PDF of the *value distribution* can be approximated using the *plug-in approach*, by approximately “*inverting the bid*” and using the “*recovered value as a truthful sample*”

$$\hat{v}_{ij} = b_{ij} + \frac{\hat{G}(b_{ij})}{(n-1)\hat{g}(b_{ij})}$$

$$\hat{F}(z) \stackrel{\text{def}}{=} \frac{1}{n \cdot m} \sum_{i,j} 1\{\hat{v}_{ij} < z\}, \quad \hat{f}(z) = \frac{1}{n \cdot m} \sum_{i,j} \frac{1}{h_n} K\left(\frac{\hat{v}_{ij} - z}{h_n}\right)$$

Formal Guarantees

- Suppose pdf f has R uniformly bounded continuous derivatives
- If we observed values then error rate of $\left(\frac{nm}{\log(nm)}\right)^{-\frac{R}{2R+1}}$ [Stone'82]
- Now that only bids are observed, [GPV'00] show that best achievable is: $\left(\frac{nm}{\log(nm)}\right)^{-\frac{R}{2R+3}}$
- The density f depends on the derivative of g

Why useful?

Scientific: economically meaningful quantities

Perform counter-factual analysis: what would happen if we change the game?

Performance measures: welfare, revenue

Testing game-theoretic models: if theory on estimated quantities predicts different behavior, then in trouble

What if all we want is to compare
between auctions A and B in
terms of revenue?

What I could potentially do is:

For each auction flip a coin;

If heads, then run auction A else run auction B

After many auctions compare **average**

revenue from A auctions, vs., **average revenue**

from B auctions

Is this correct?

We will see that it can be problematic and needs thought of how to analyze such data or structure such A/B tests!

Experimentation (aka A/B Testing)

The Basics of A/B Testing

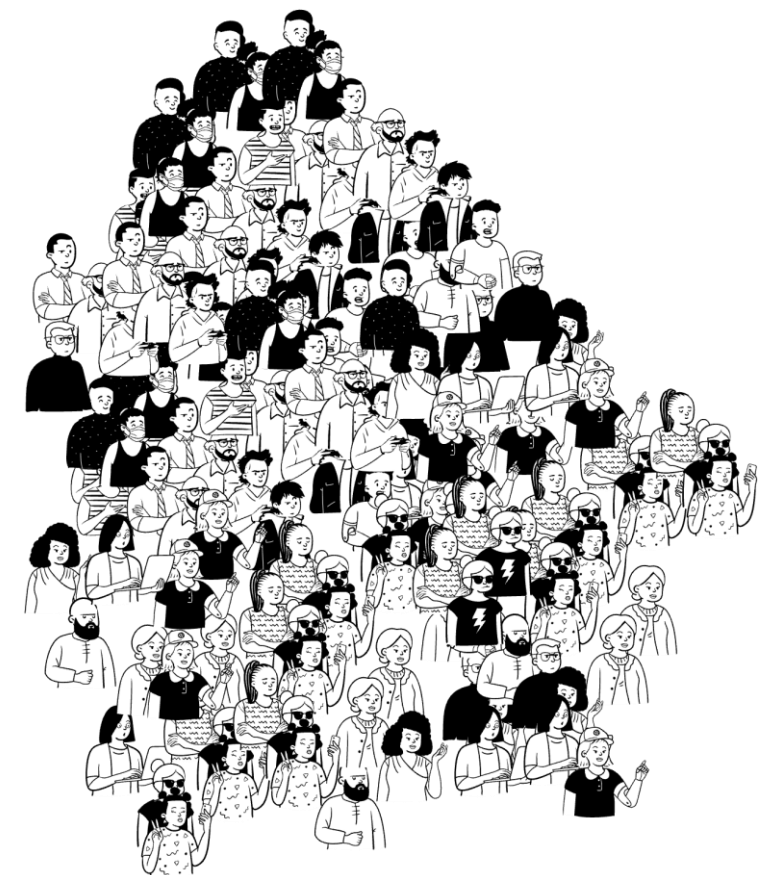
Randomization, Causality, Statistical Inference



The Mechanics

A/B Testing

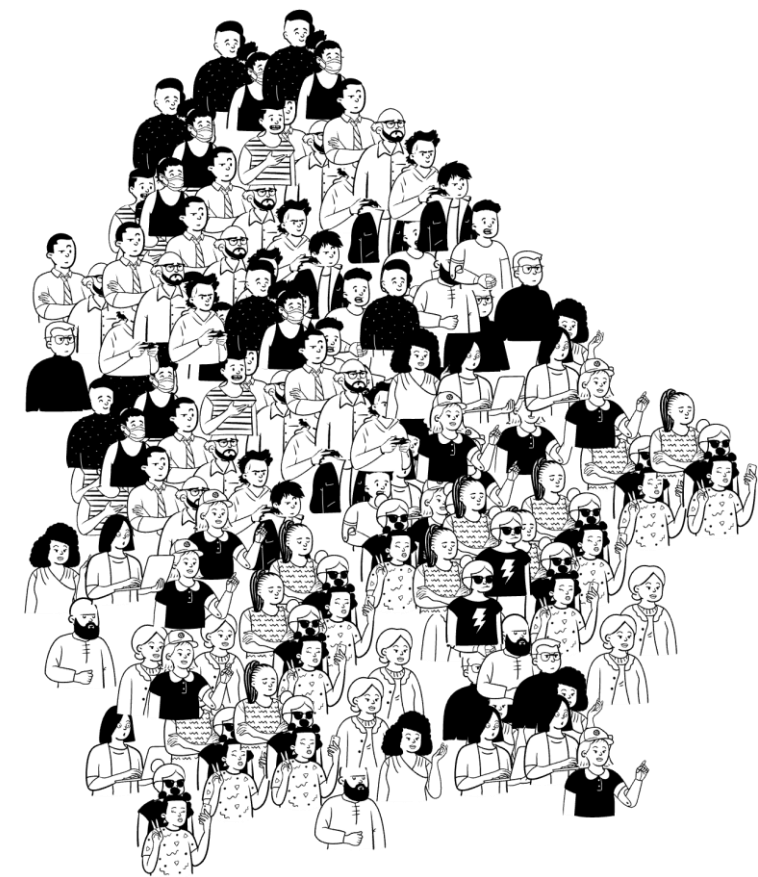
user base



A/B Testing

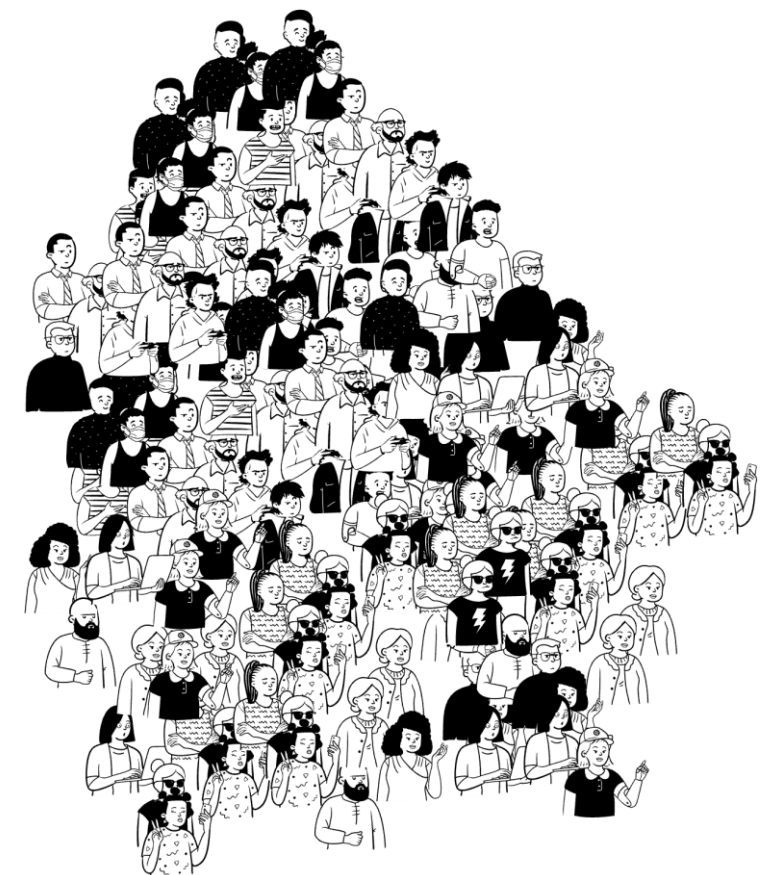
user base

sample



A/B Testing

user base



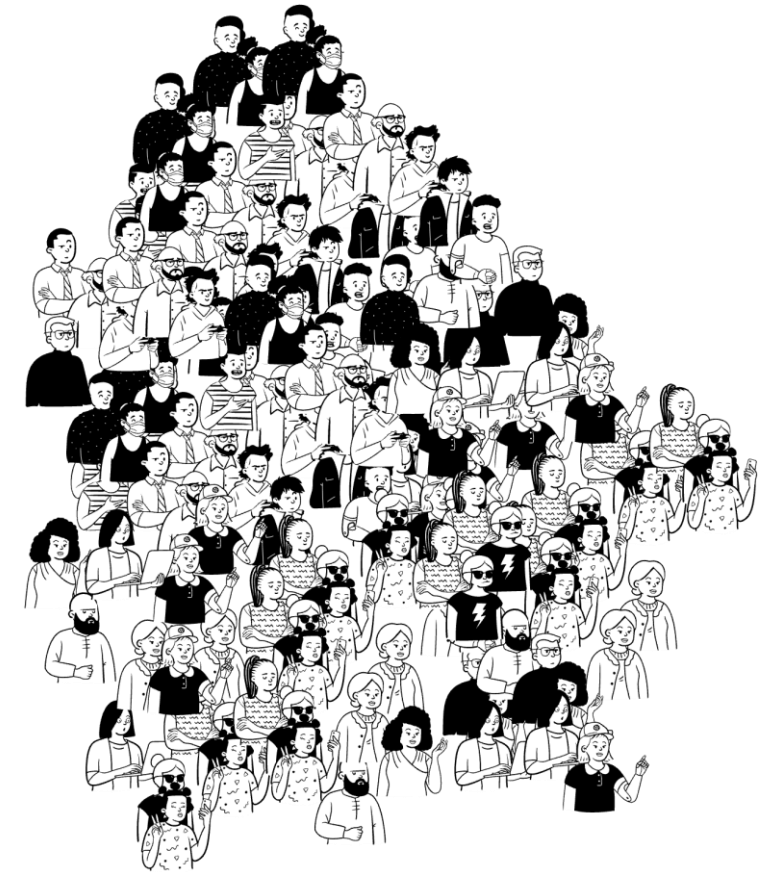
sample



flip a coin for each user

A/B Testing

user base



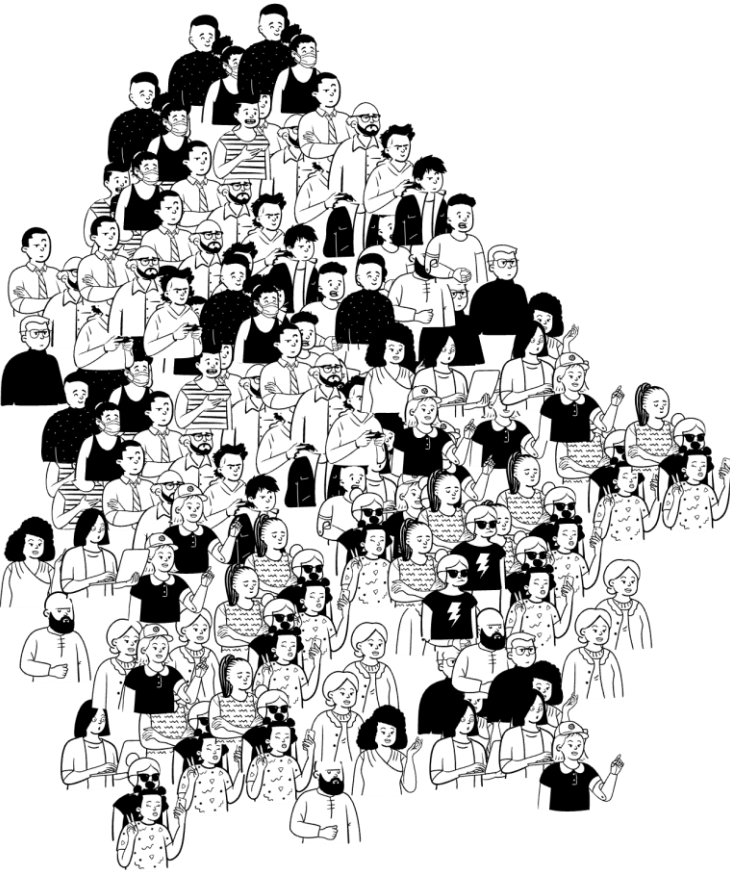
sample



split into groups based on coin

A/B Testing

user base



Group A

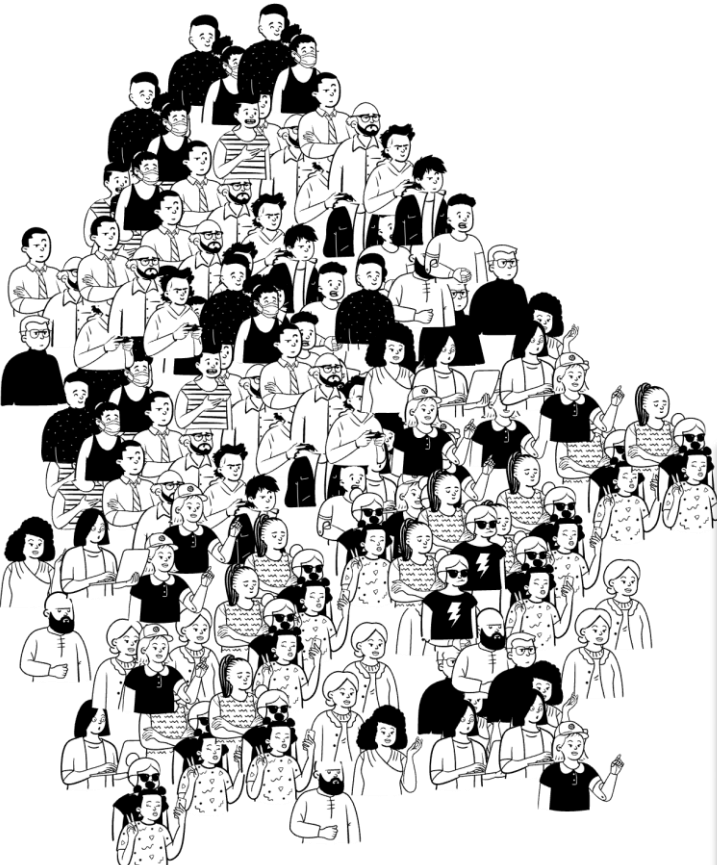


Group B




A/B Testing

user base



A

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Group A



B

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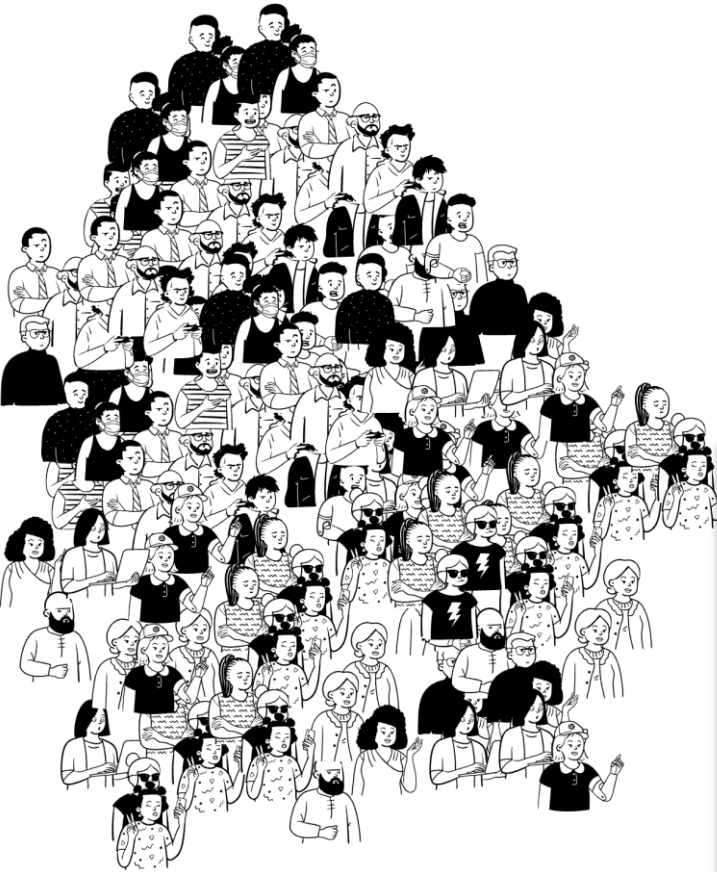
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Group B




A/B Testing

user base



A

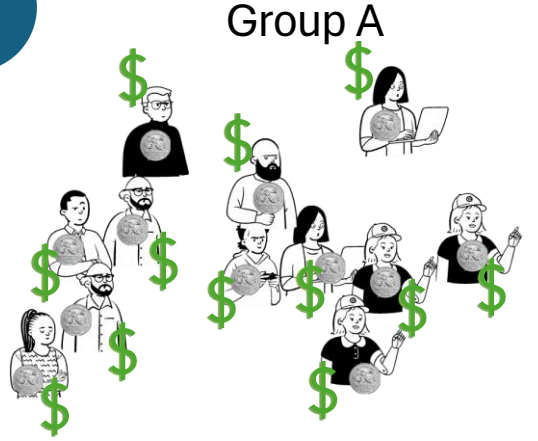
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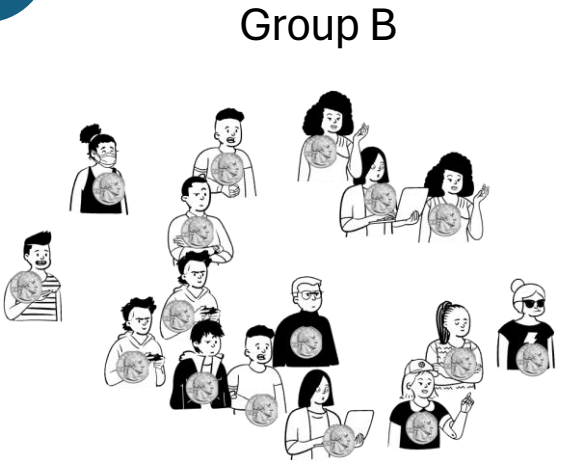
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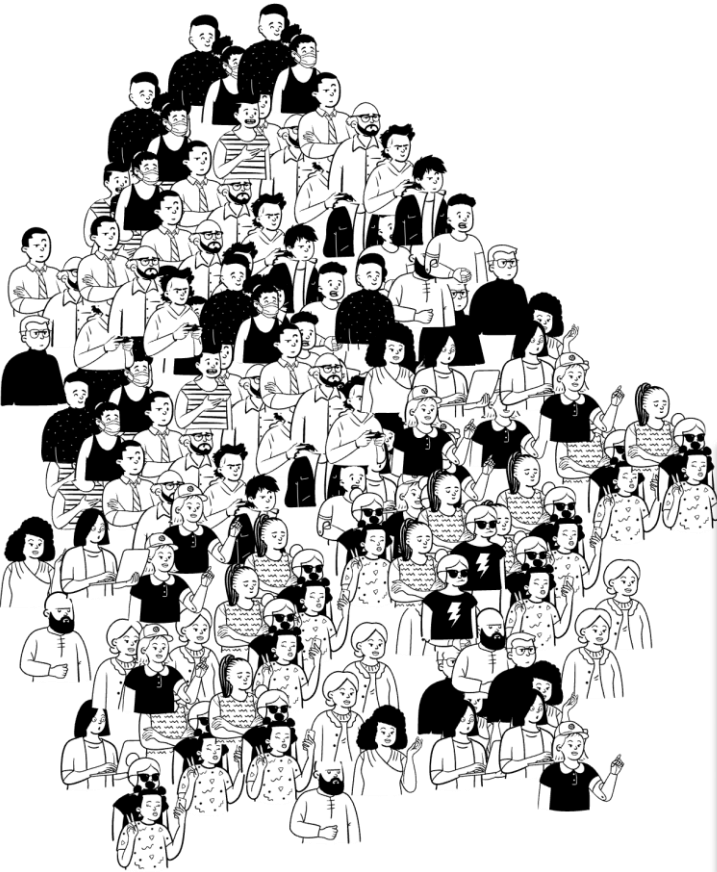
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A/B Testing

user base



A

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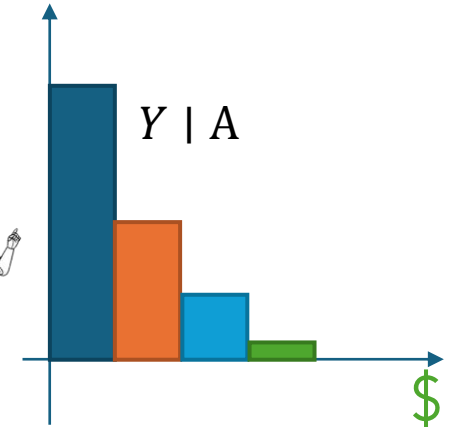
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% of people


Group A



$\mu_A = 10\$$ (average spend)

B

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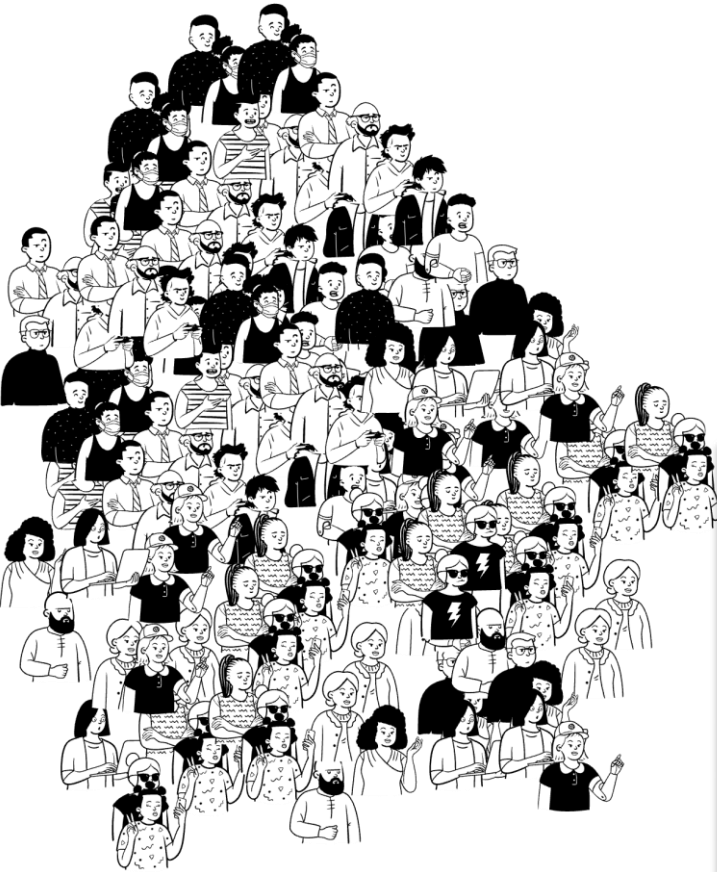
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Group B

$\mu_A = 10\$$ (average spend)


A/B Testing

user base



A

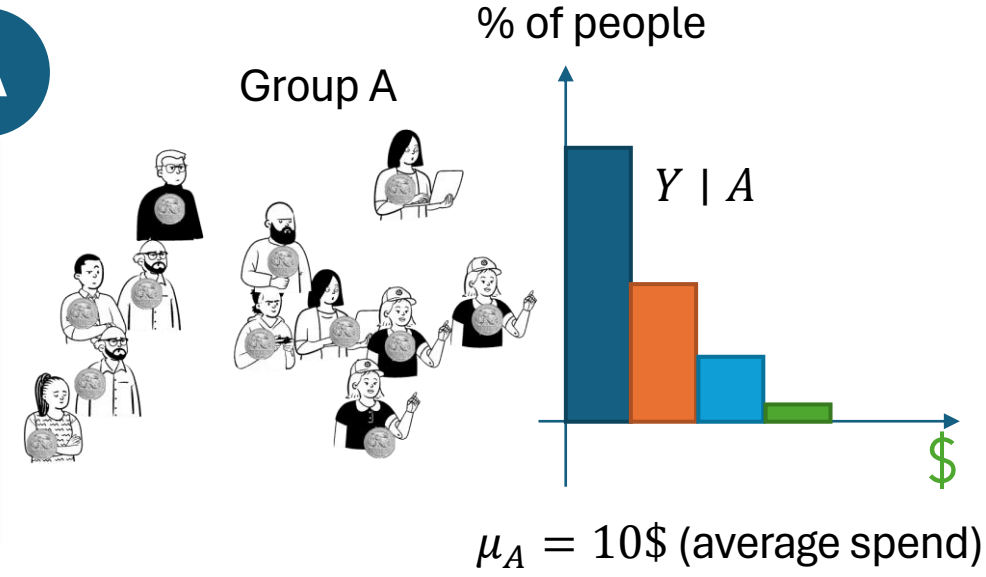
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B

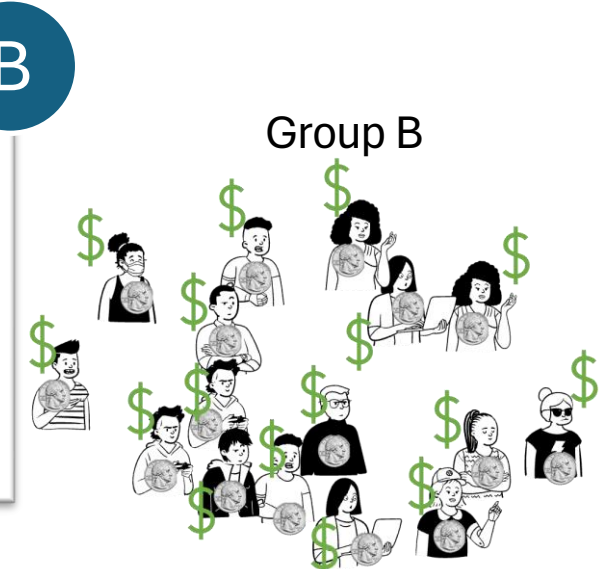
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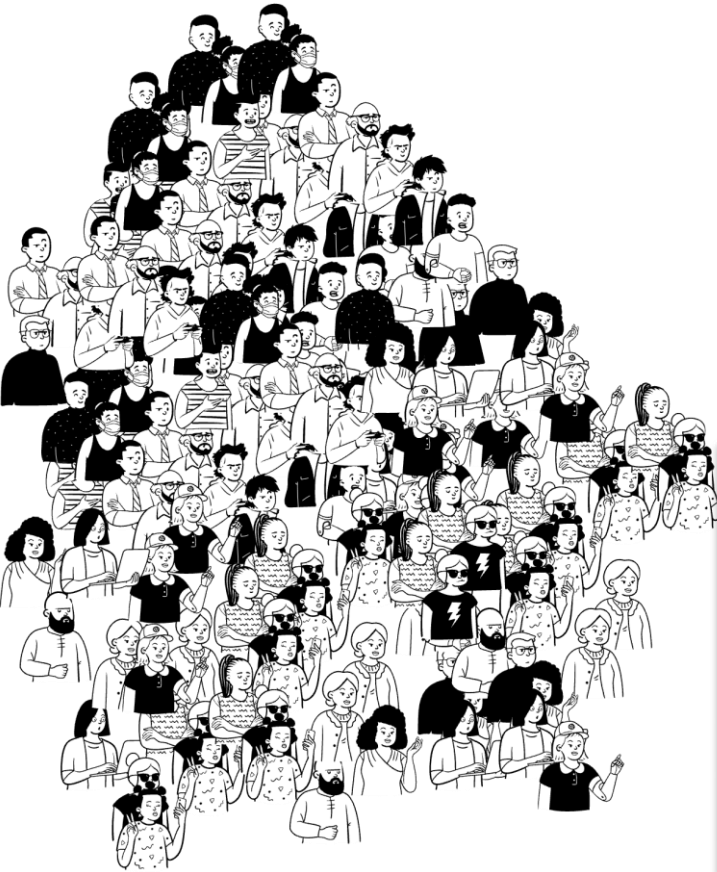
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A/B Testing

user base



A

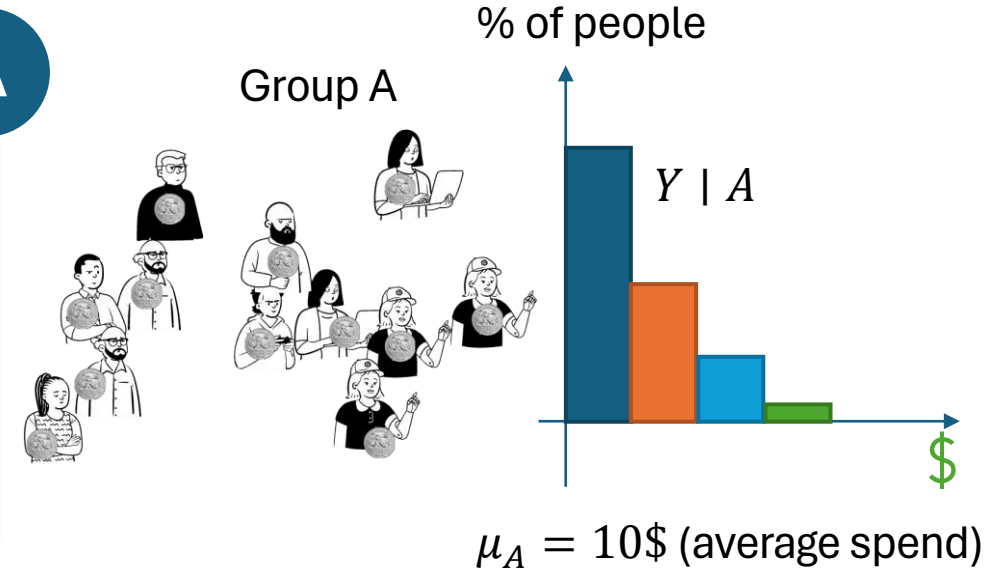
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B

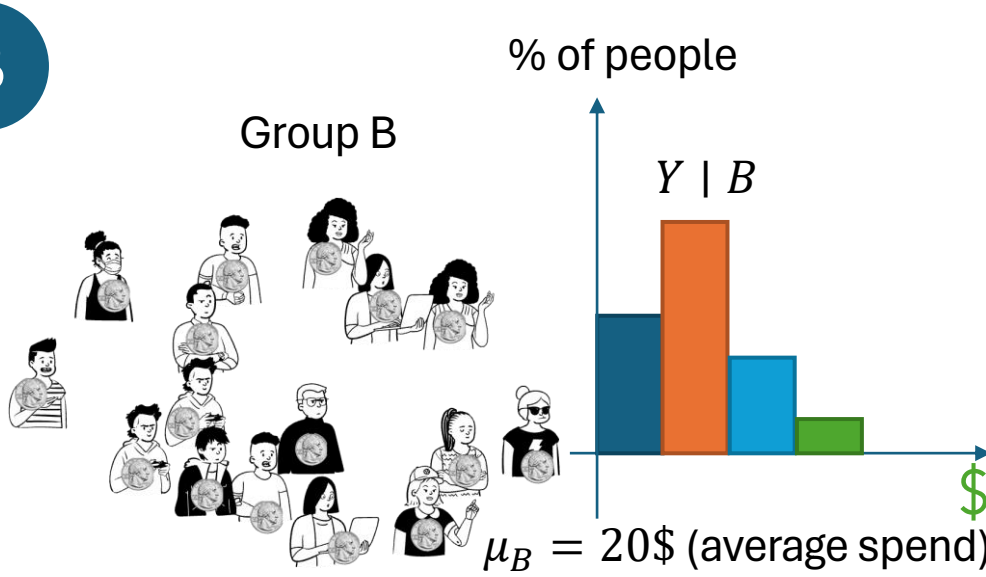
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A/B Testing

Control
Baseline
Status quo

A

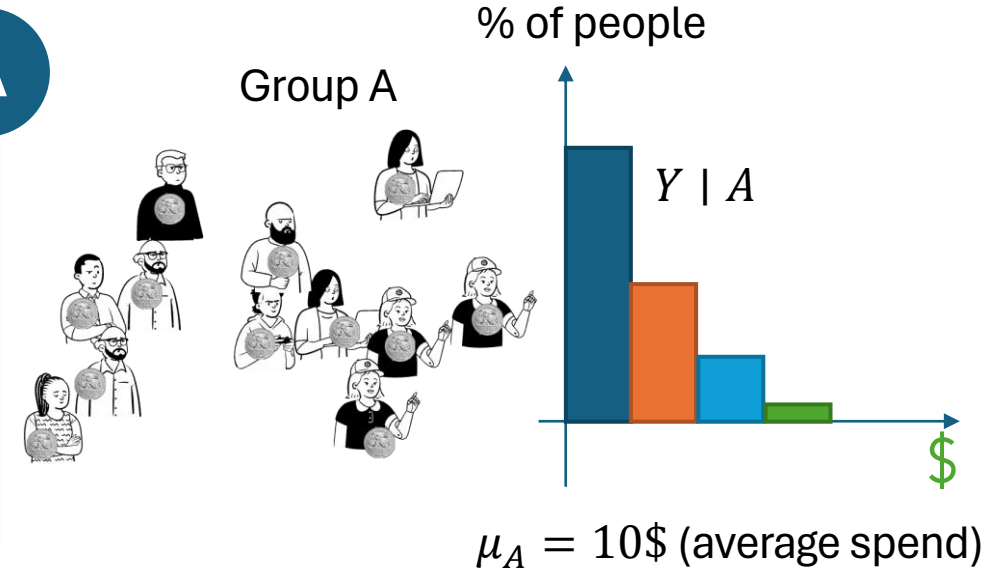
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Treatment
Innovation

B

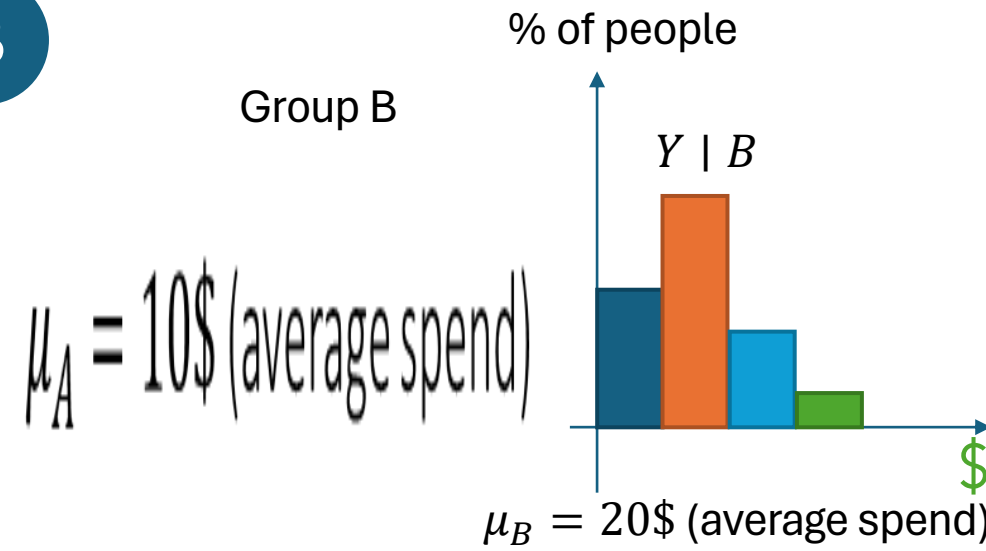
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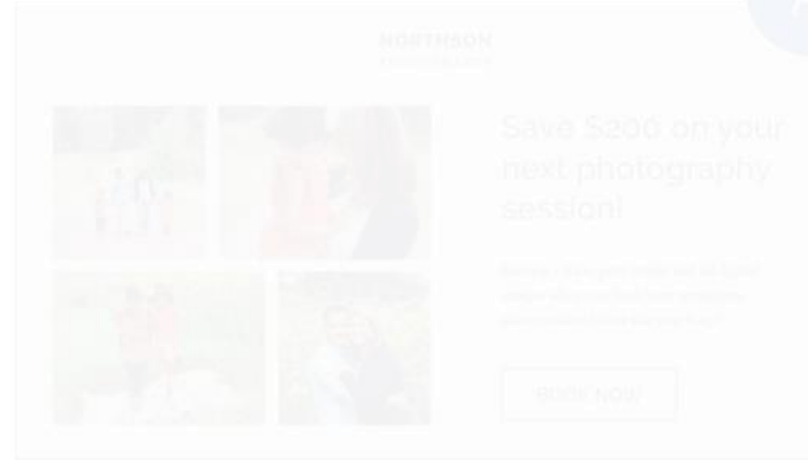
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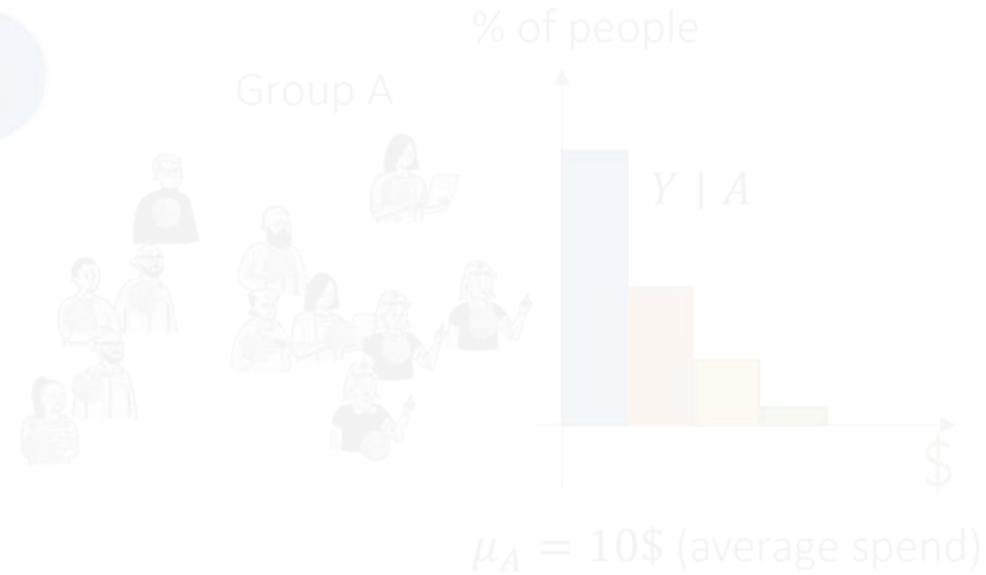


A/B Testing

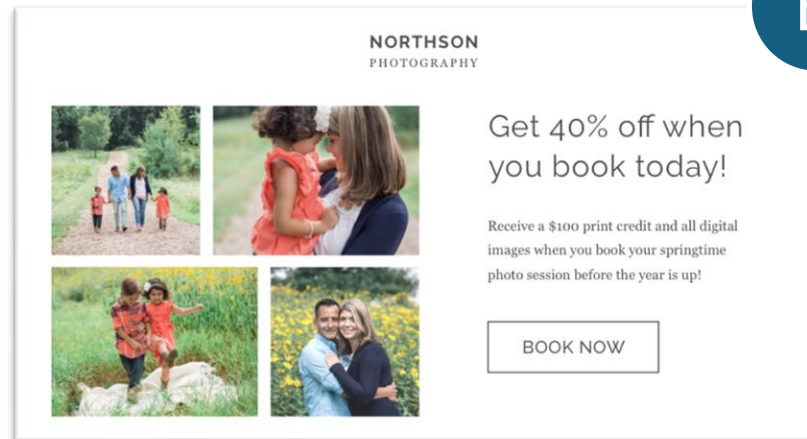
Control
Baseline
Status quo



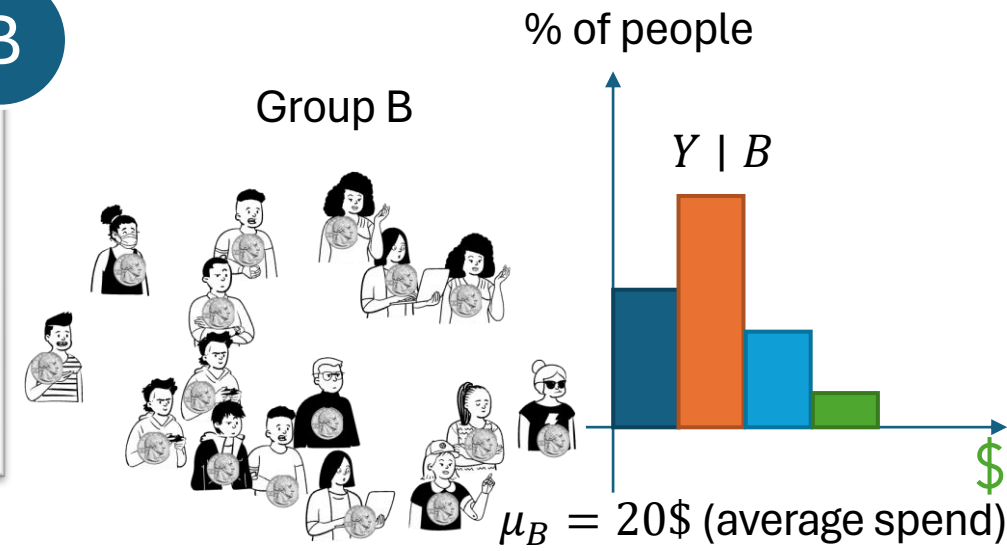
A



Treatment
Innovation

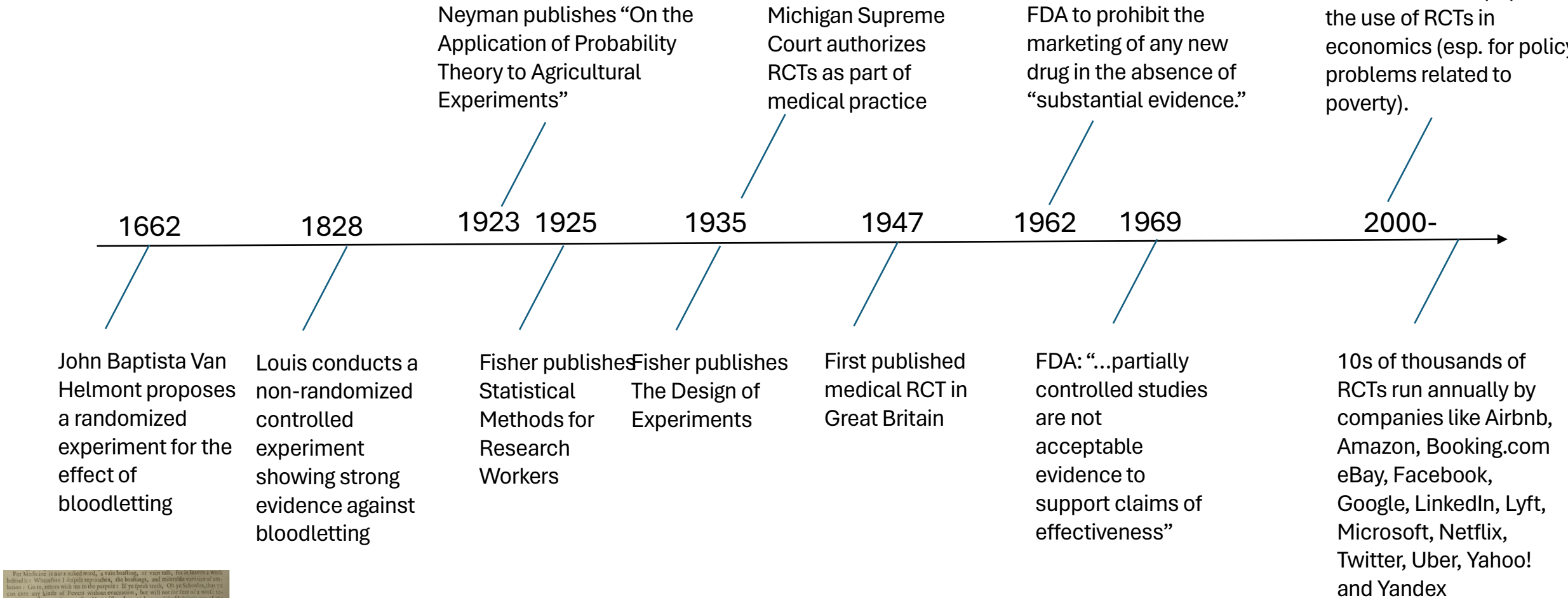


B



A Brief History of Experimentation

Abhijit Banerjee, Esther Duflo, and Michael Kremer, popularize the use of RCTs in economics (esp. for policy problems related to poverty).



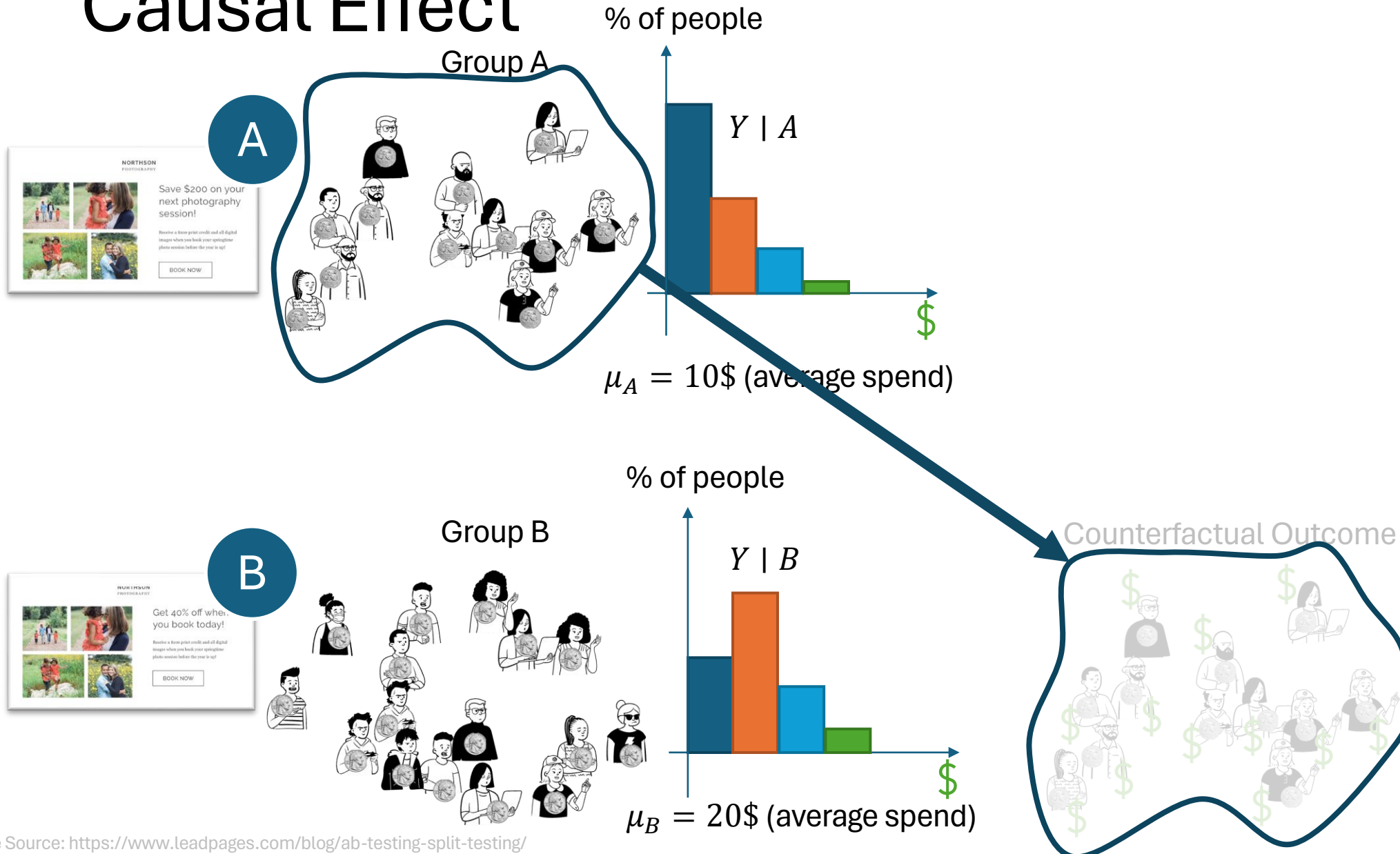
For Medicine is not a naked word, a vain boasting, or vain talk, but it leaves a wide behind it. Wherefore I desire reproaches, the beatings, and miserable verities of temptation. Go to, return with me to the papule. If ye speak truth, Oh ye Scholars, that ye can cure any kind of Fever without evacuation, but will not let us of a word, as ye have, come down to the control ye Humors. Let us take out of the Hospital of the Camp, or from elsewhere, such, or few poor People, that have Fevers, Pleurisy, or any other kind of illness, let us call lots, that one half of them may fall to my share, and the other to yours. I will cure them without blood-letting and terrible evacuation; but do you do, as ye know (for neither do I try you up to the boiling, or to Phlebotomy, or the abundance from primitive Medicine) see shall see how many I shall cure both of us shall have. But let the reward of the contention or wage, be your Pleasure, deposited on both sides. Here your business is decided. Oh ye Magistrates, come

RCTs are the gold standard for measuring the “causal effect” of a “treatment” on an “outcome”

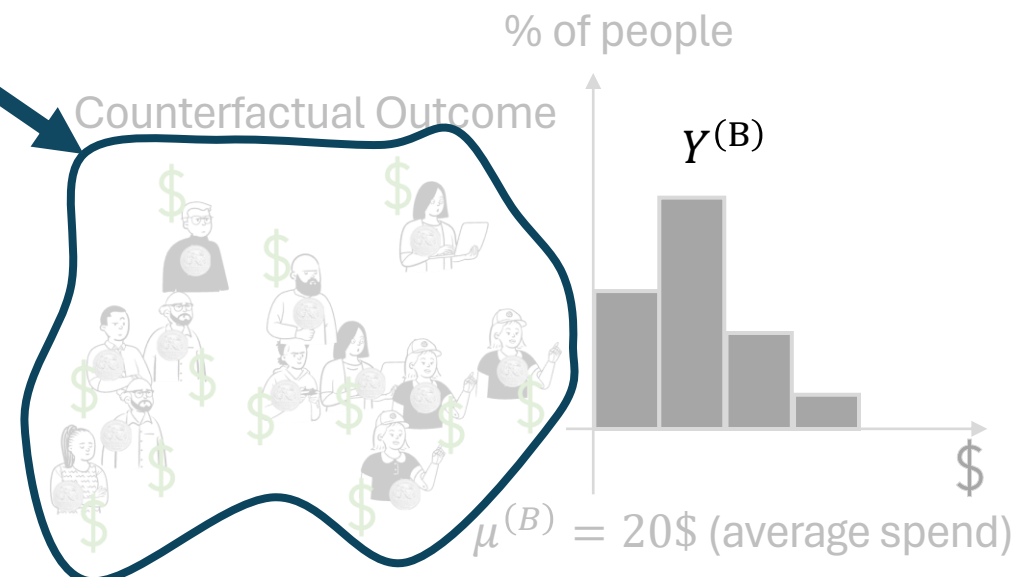
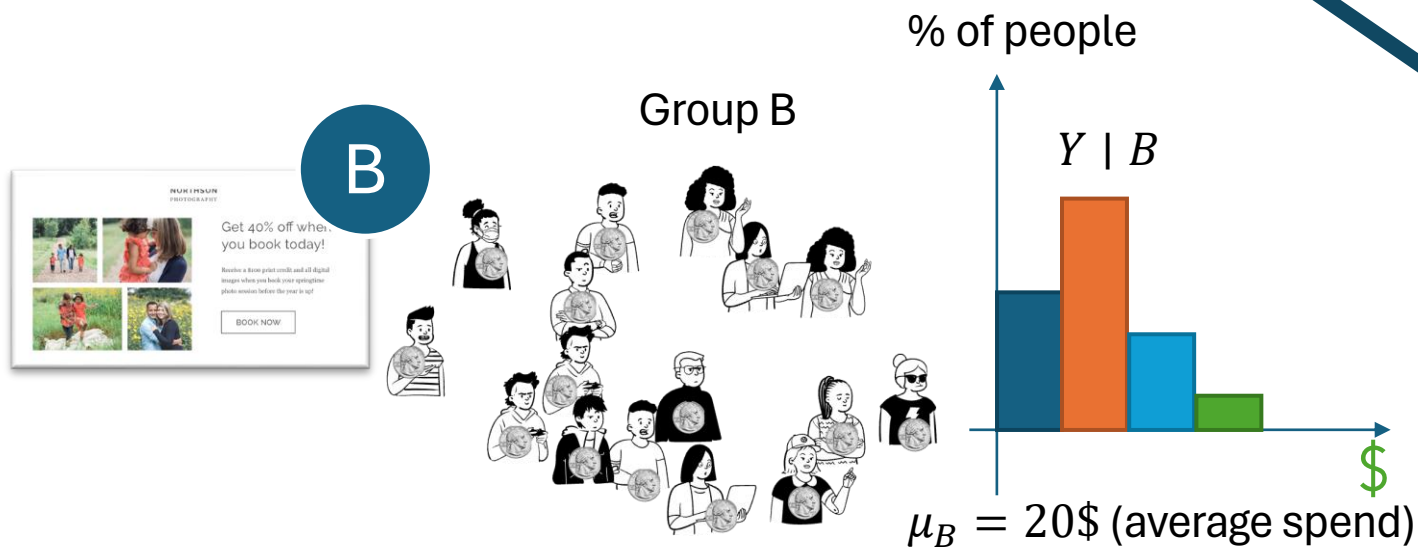
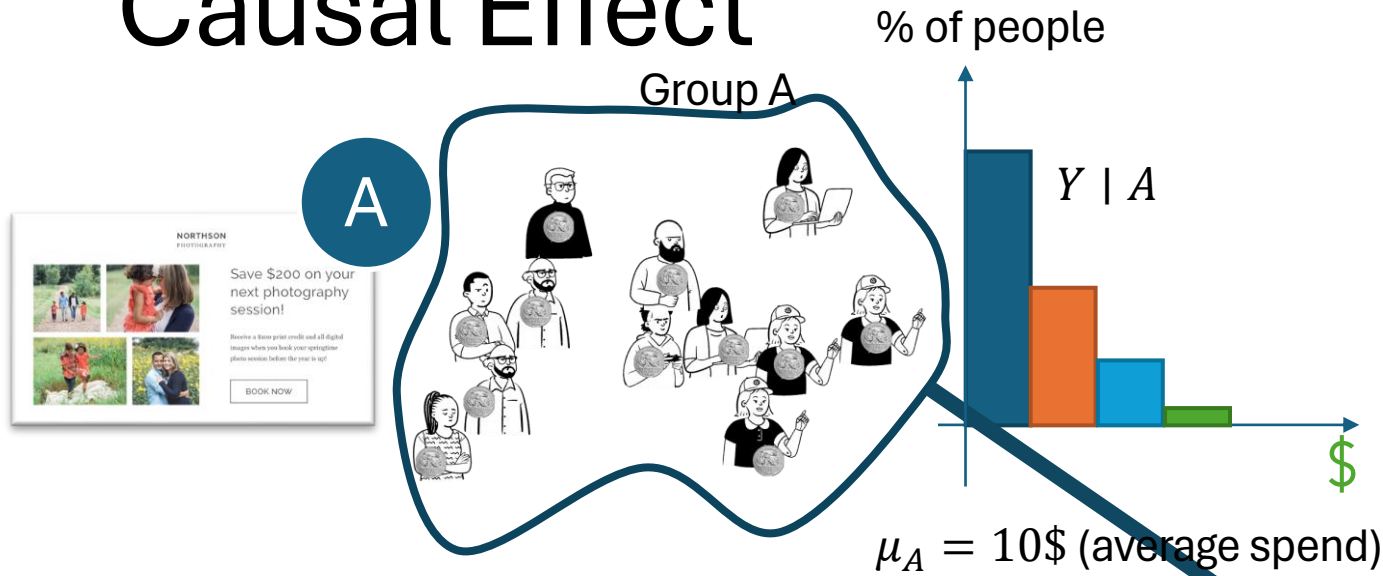
Causality

The background features a complex network of thin, light grey lines that create a sense of interconnectedness. Scattered throughout this network are numerous circles of varying sizes. The color palette for these circles is a gradient, starting with deep reds and maroons in the upper right quadrant and transitioning through purples to dark blues in the lower left quadrant. The overall aesthetic is modern and abstract, suggesting a complex system or a web of relationships.

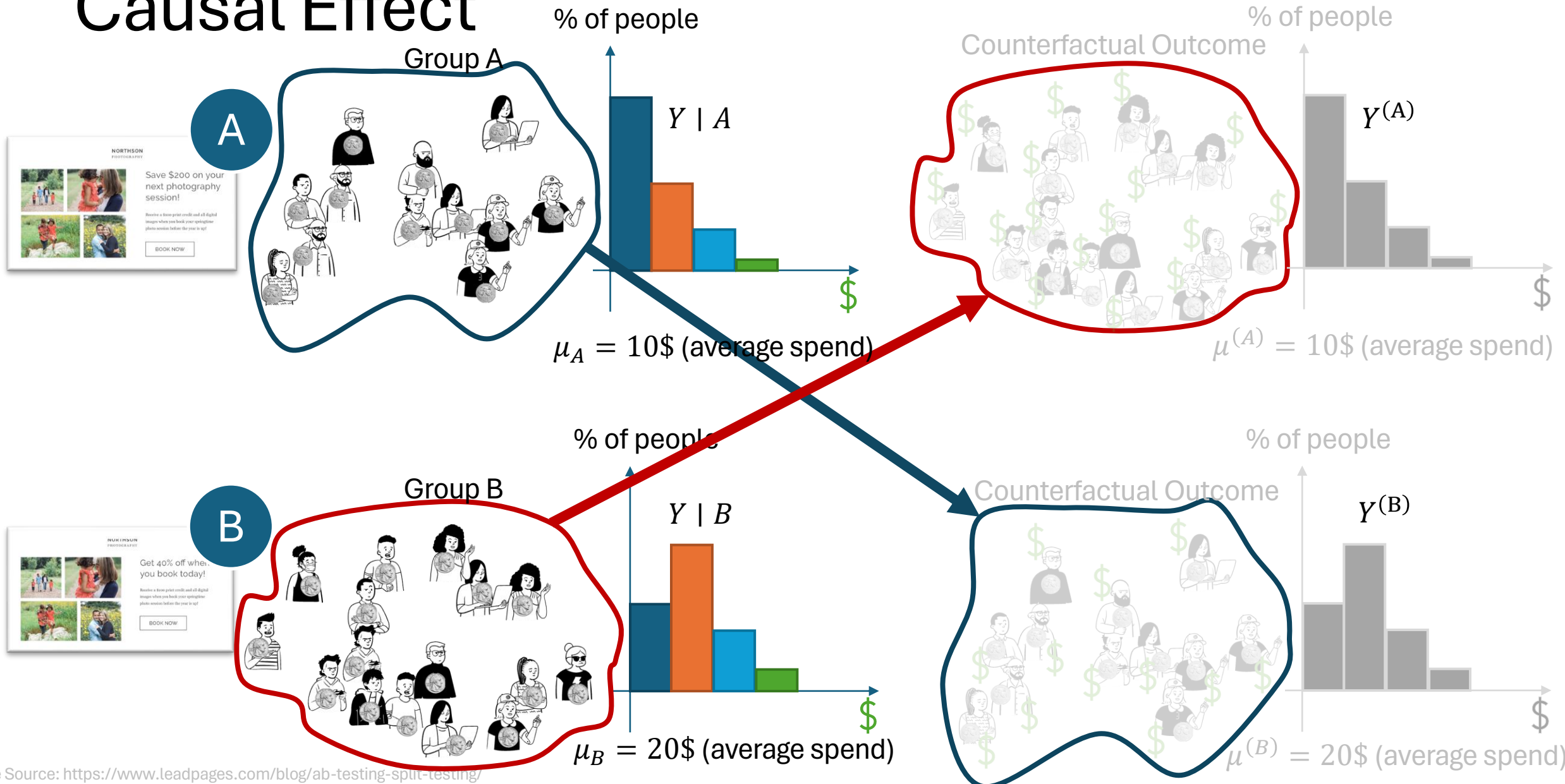
Causal Effect



Causal Effect



Causal Effect



Randomization implies

$$Y|A \sim Y^{(A)}$$

$$Y|B \sim Y^{(B)}$$

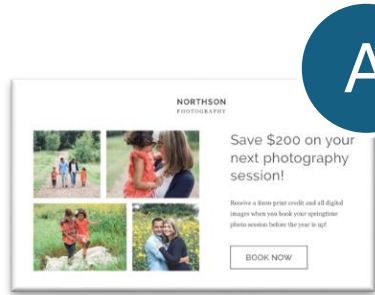
Aggregate differences between groups

$$E[Y|A] - E[Y|B]$$

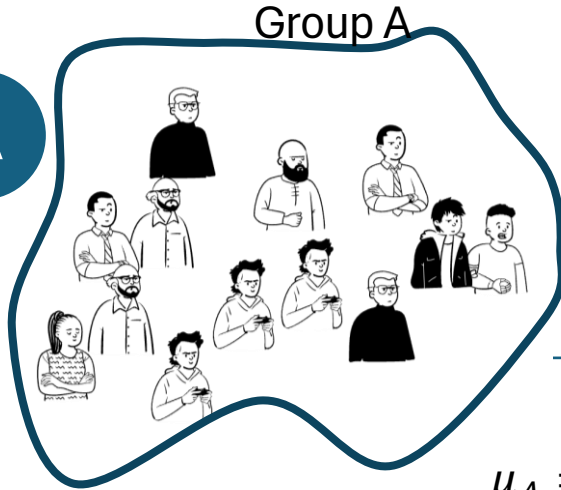
Equal aggregate causal effects

$$E[Y^{(A)} - Y^{(B)}]$$

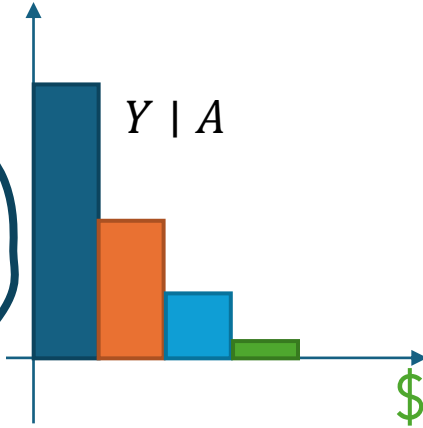
Historical Data



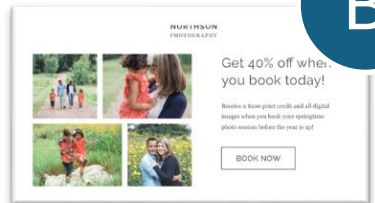
A



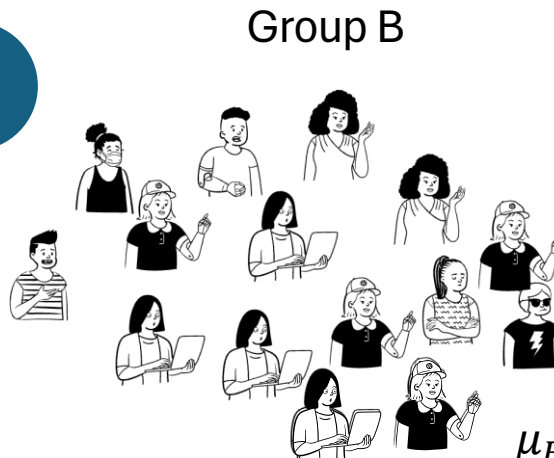
% of people



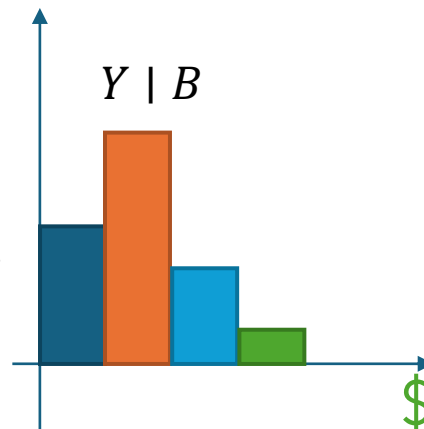
$$\mu_A = 10\$ \text{ (average spend)}$$



B



% of people

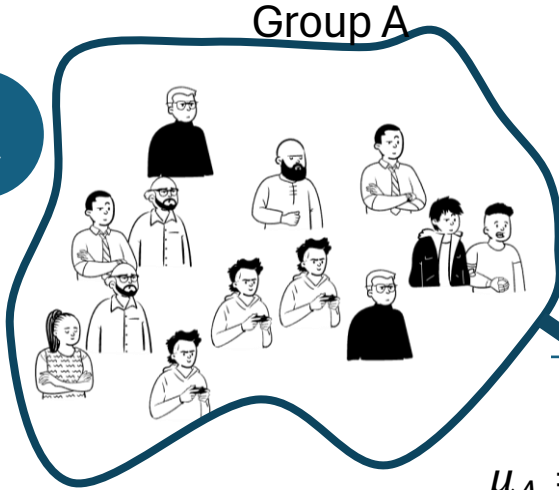
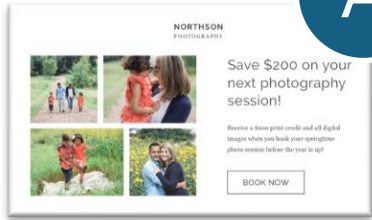


$$\mu_B = 20\$ \text{ (average spend)}$$

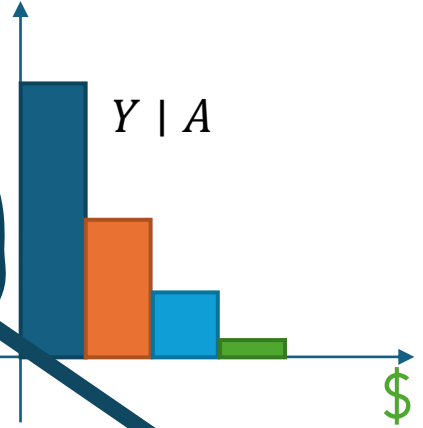
Historical Data

% of people

A

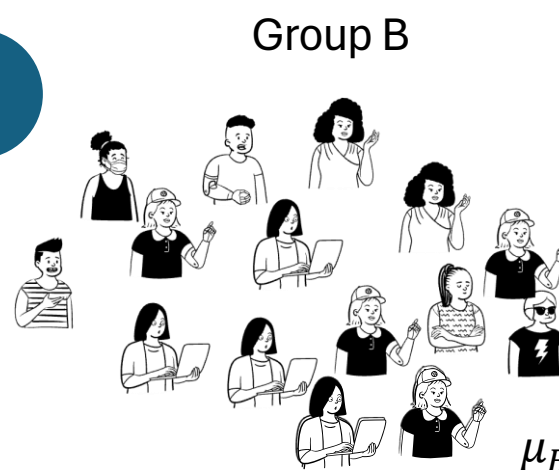
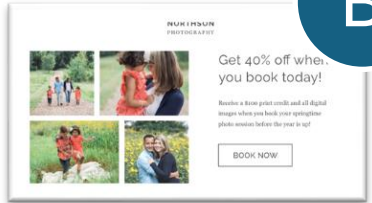


Group A

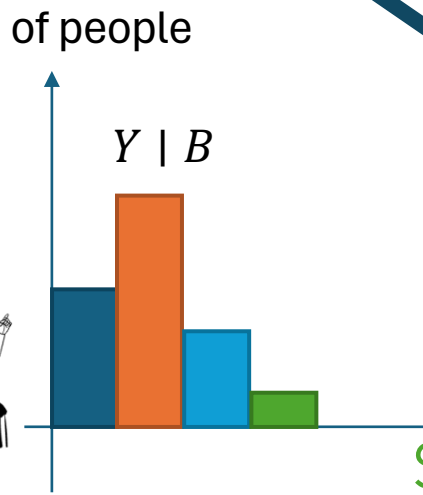


$\mu_A = 10\$$ (average spend)

B

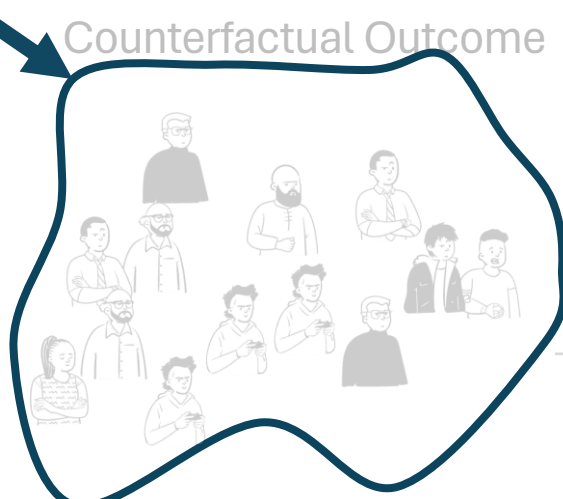


Group B

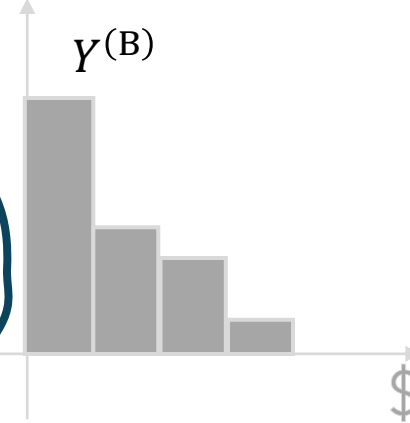


$\mu_B = 20\$$ (average spend)

% of people



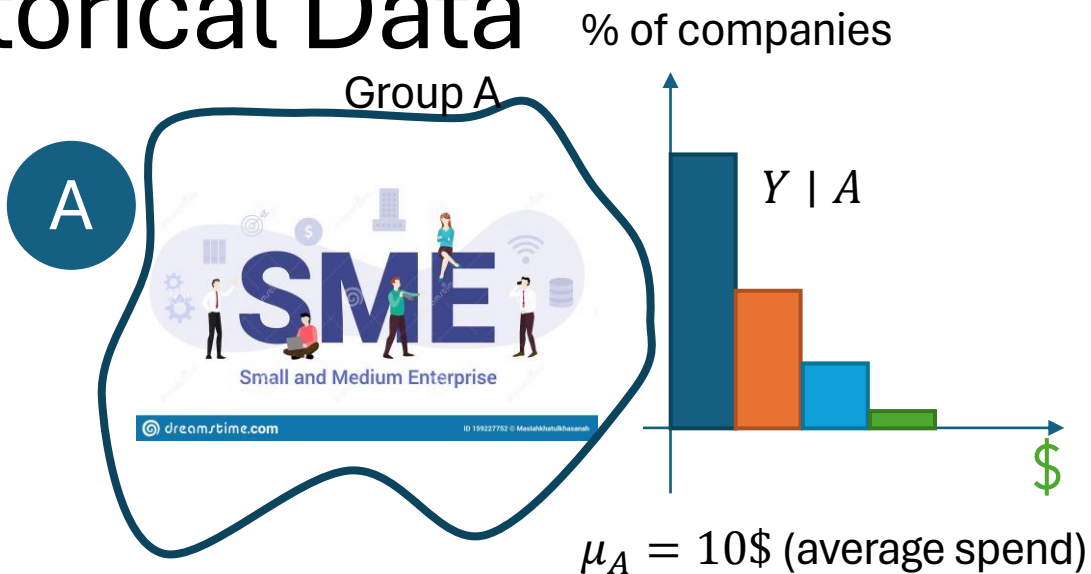
Counterfactual Outcome



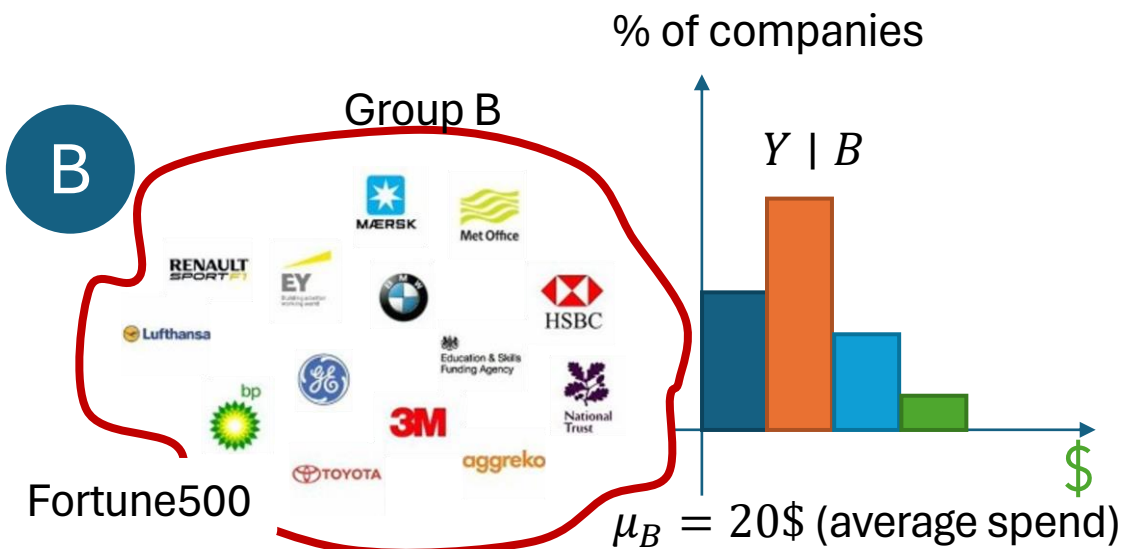
\$

Historical Data

No cloud specialist



Assign Cloud Specialist



Historical Data

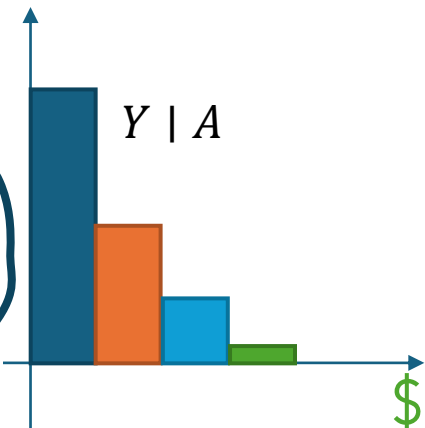
No cloud specialist



A



% of companies

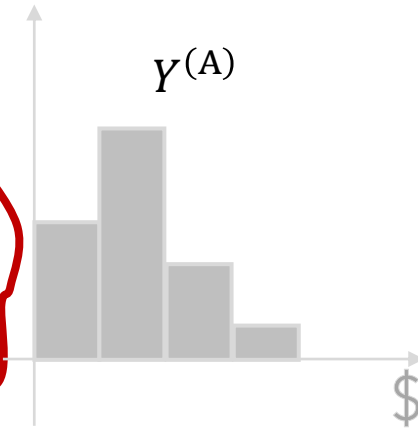


$\mu_A = 10\$$ (average spend)

Counterfactual outcome



% of companies



$\mu^{(A)} = 20\$$ (average spend)

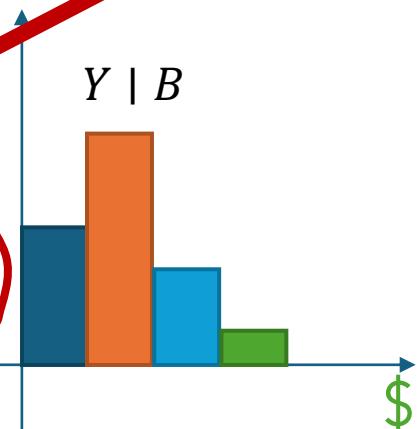
Assign Cloud Specialist



B



% of companies

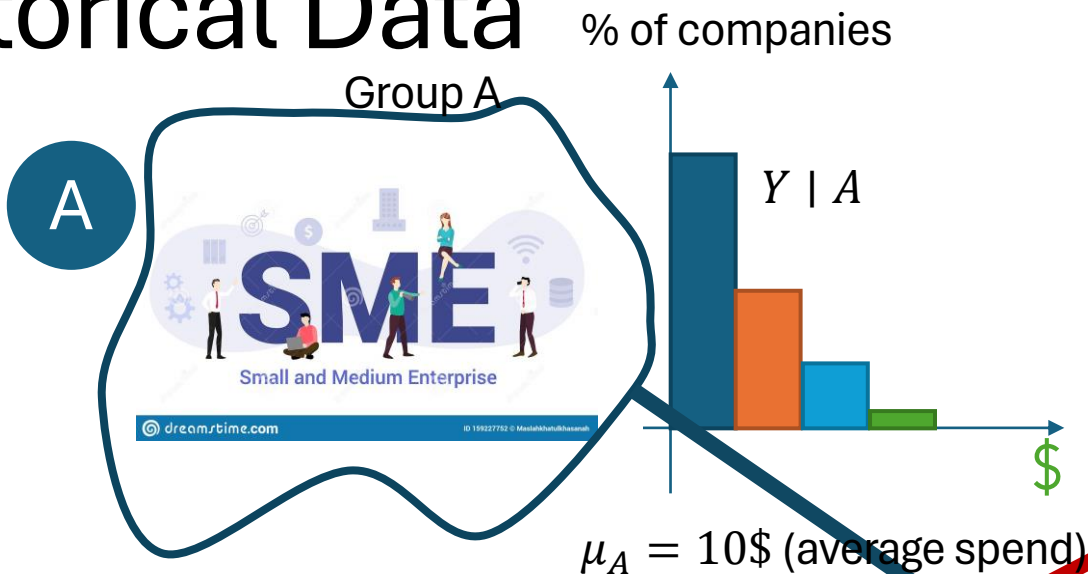


Fortune500

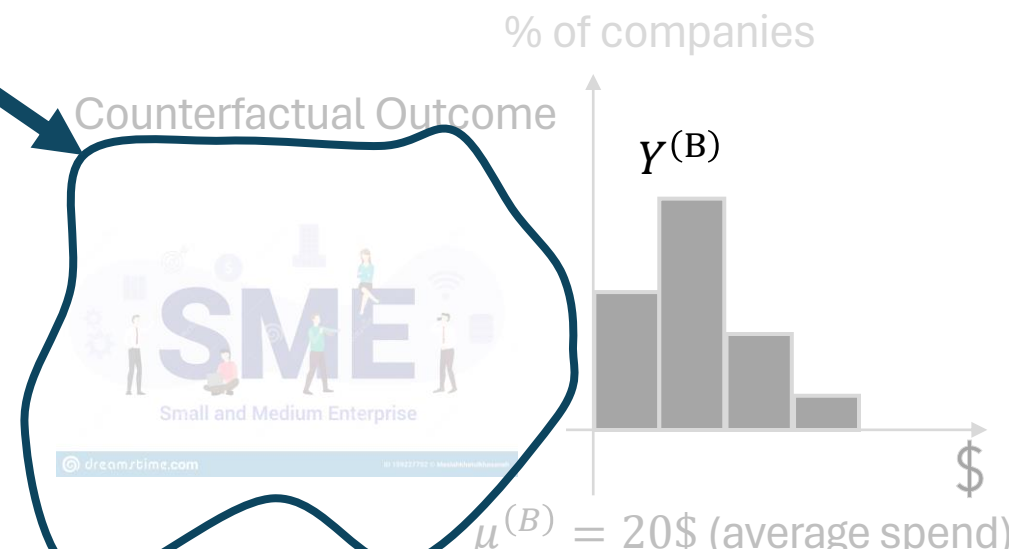
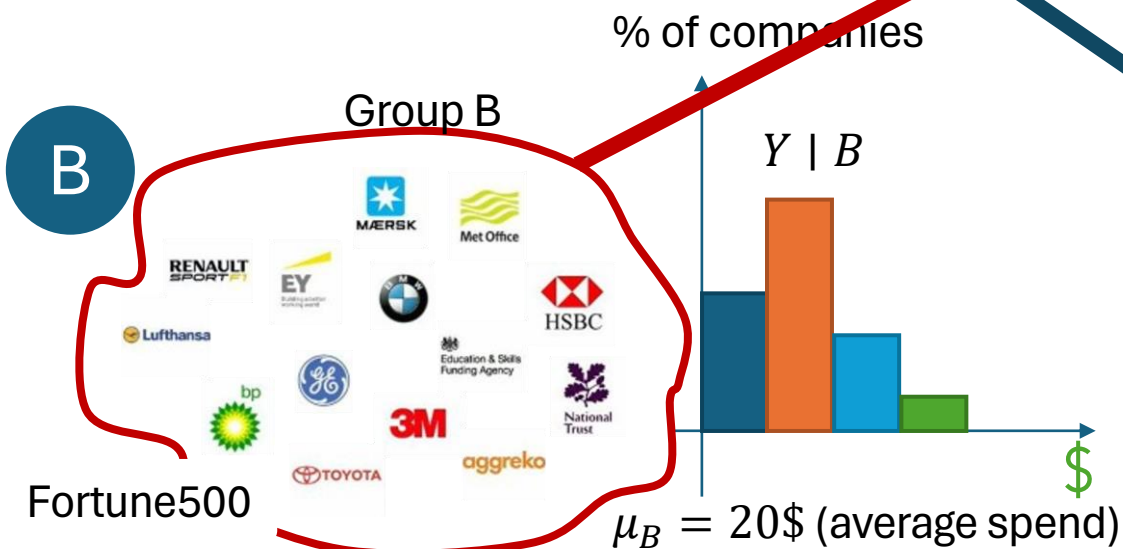
$\mu_B = 20\$$ (average spend)

Historical Data

No cloud specialist



Assign Cloud Specialist



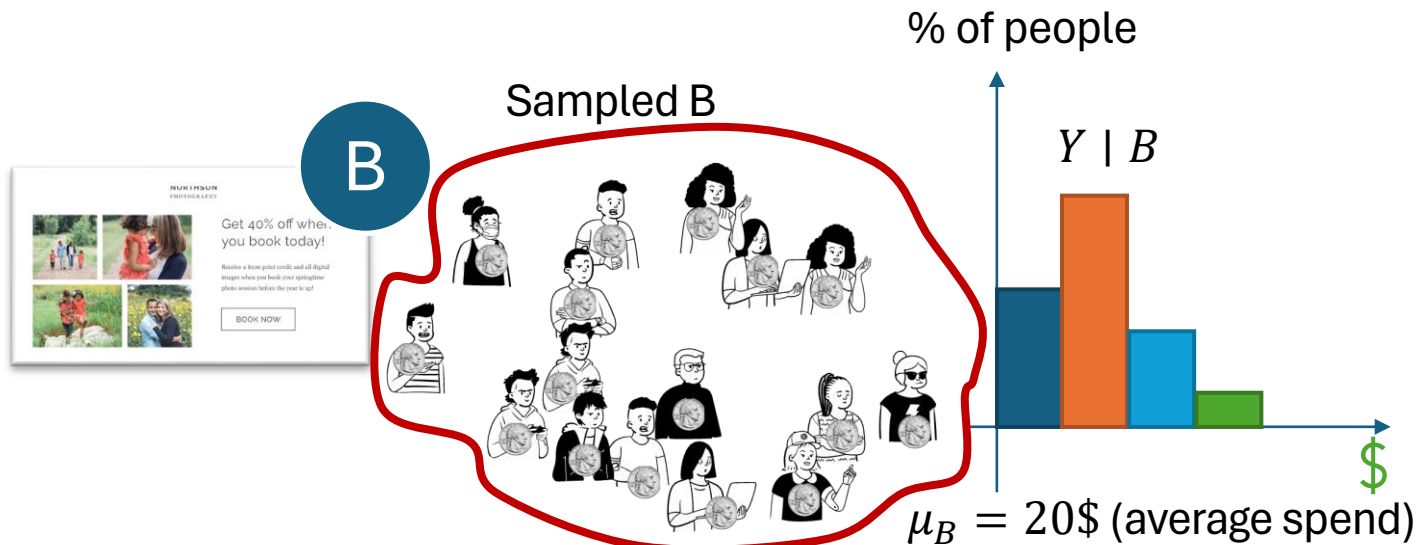
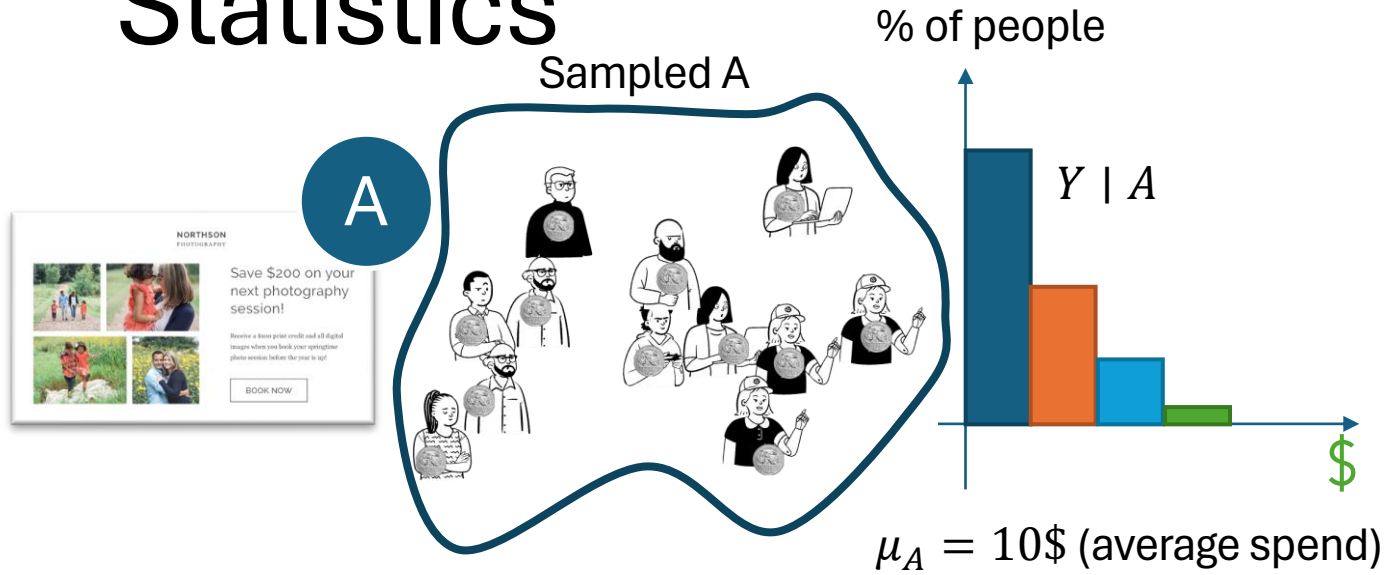


Statistics

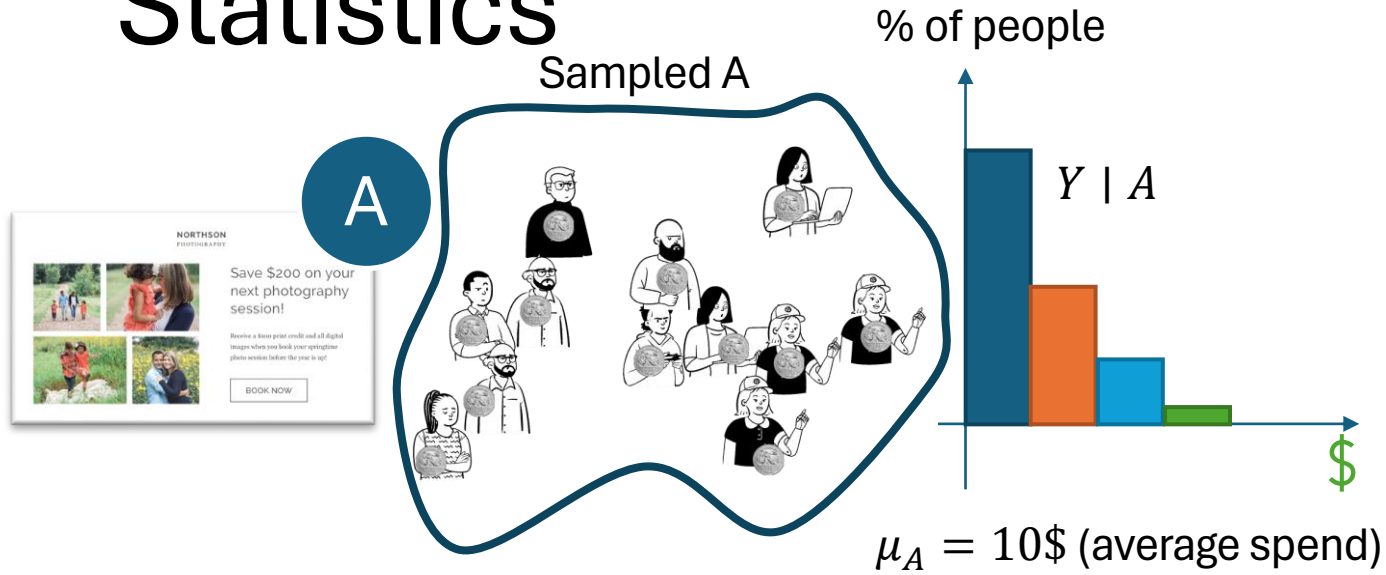
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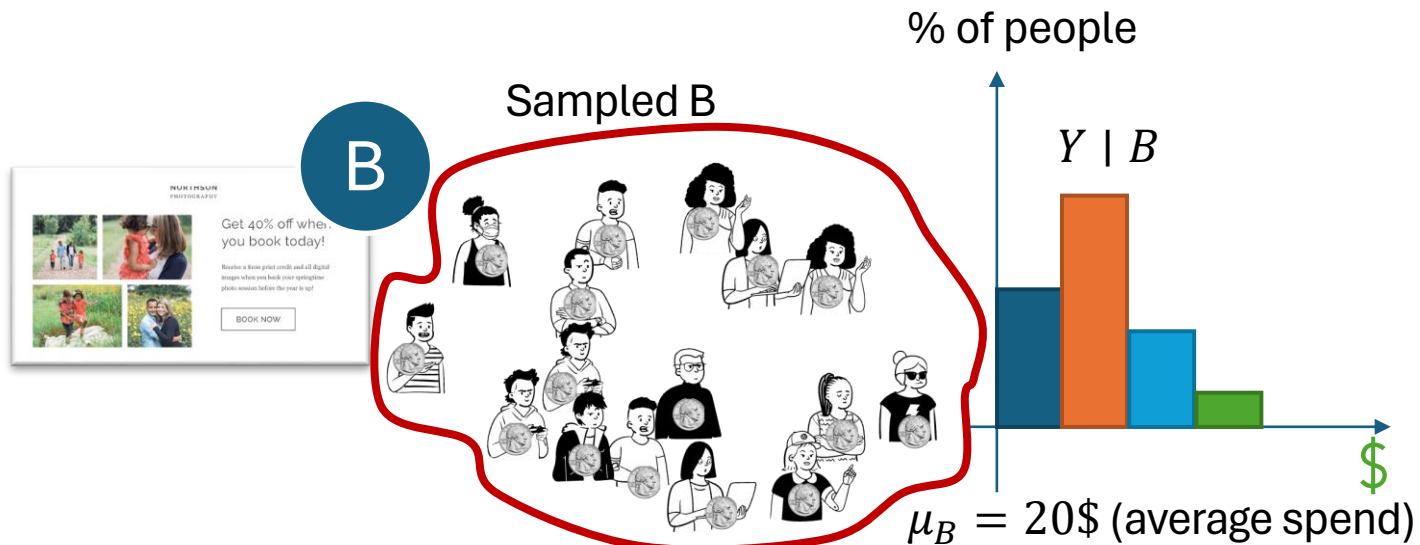
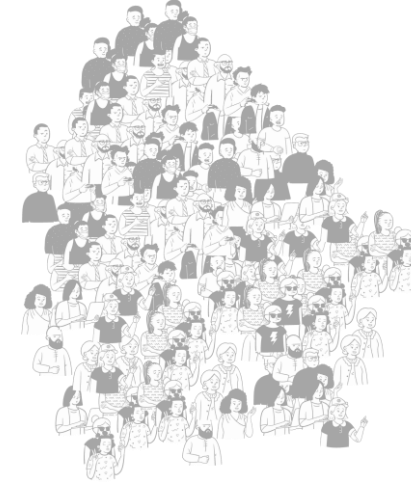
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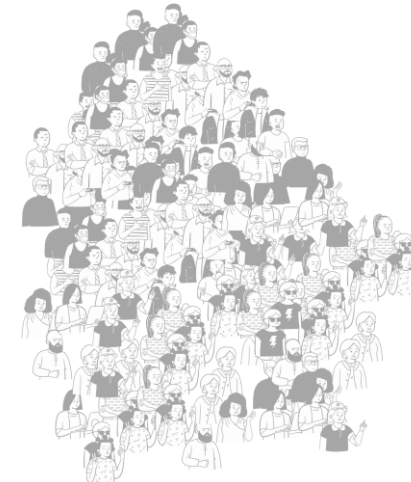
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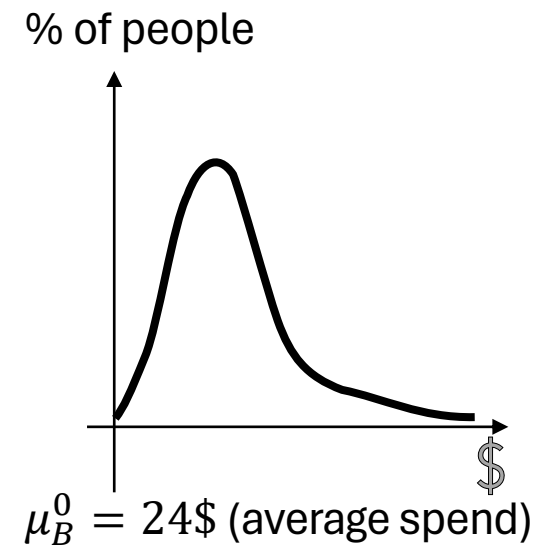
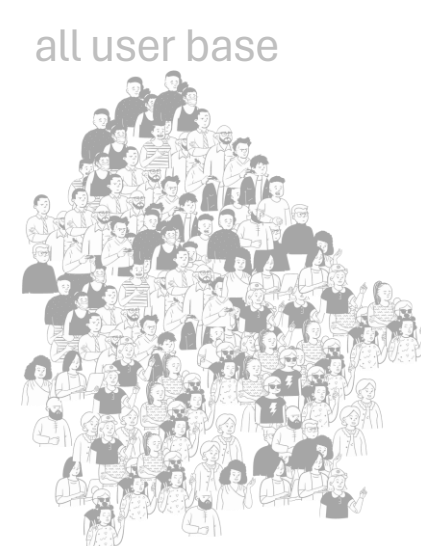
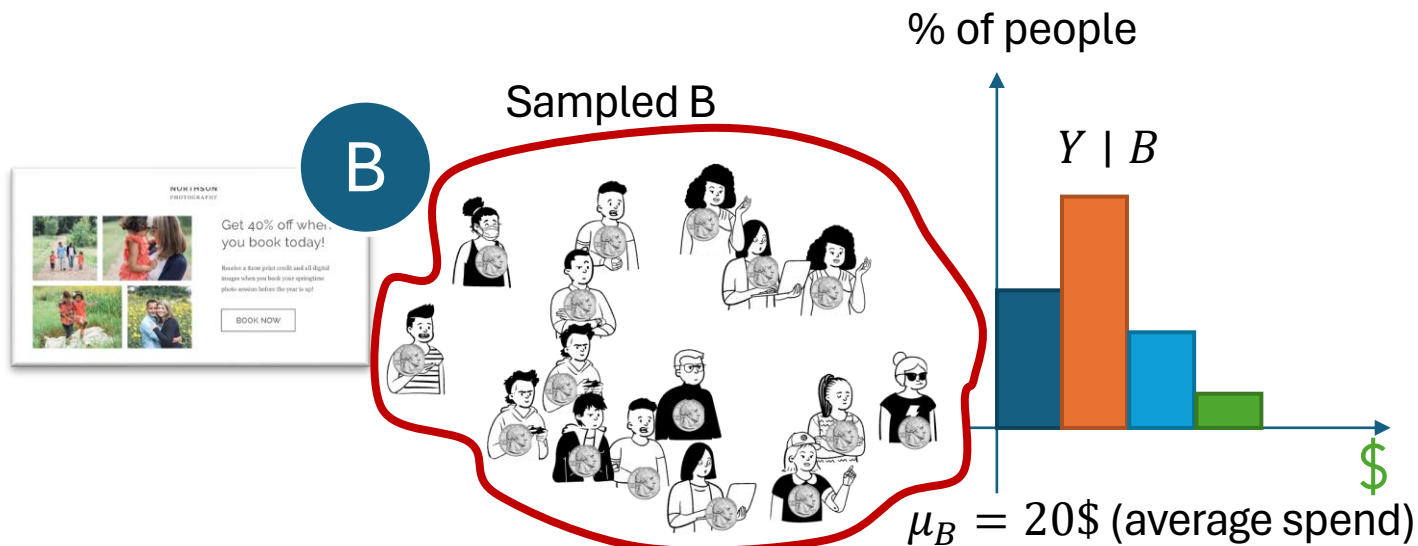
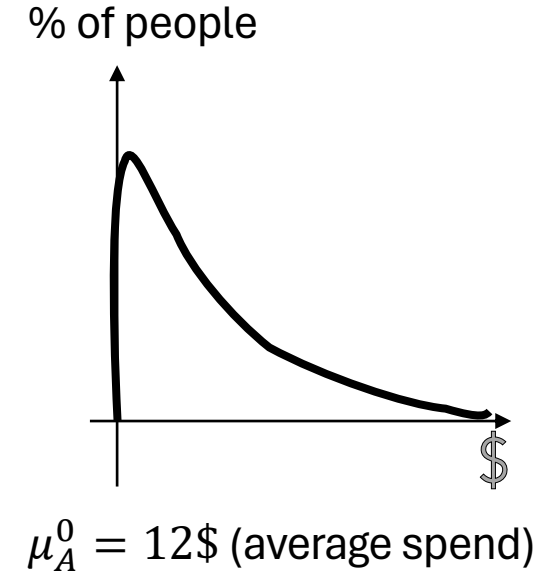
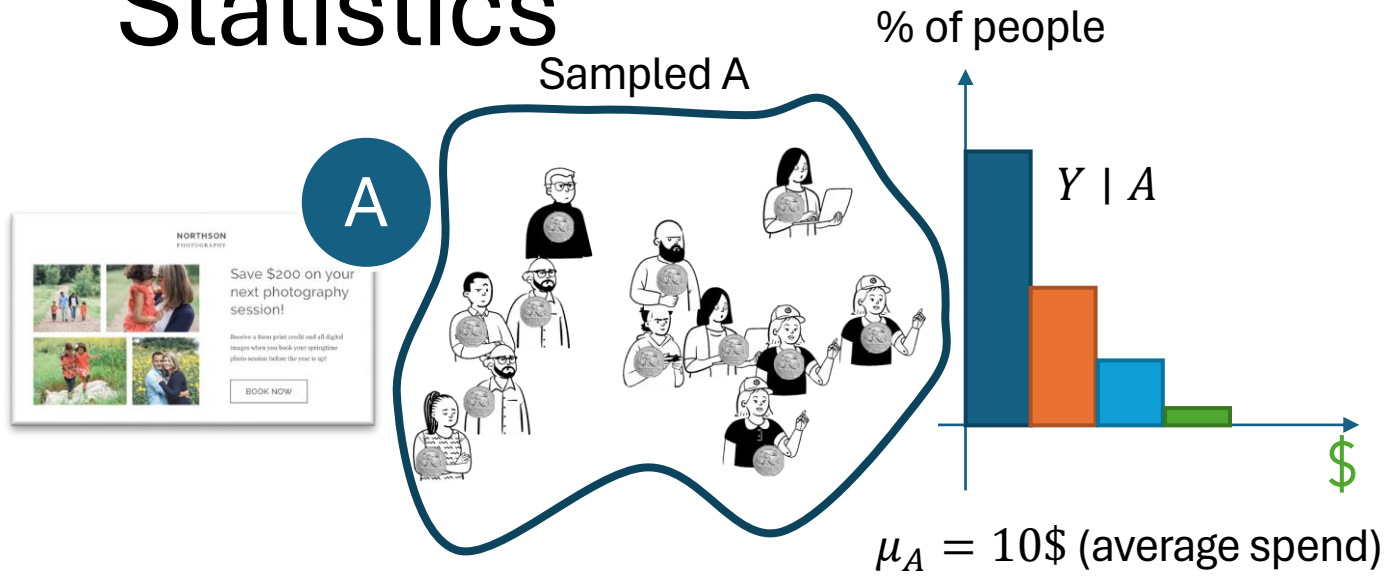
all user base



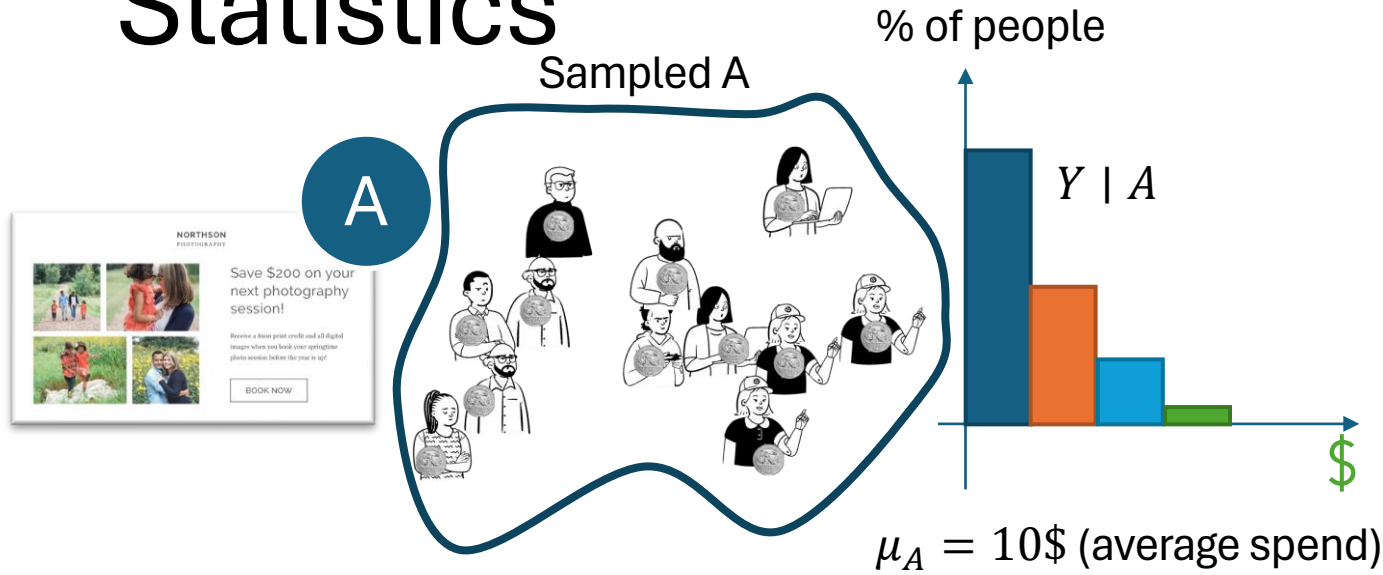
all user base



Statistics



Statistics



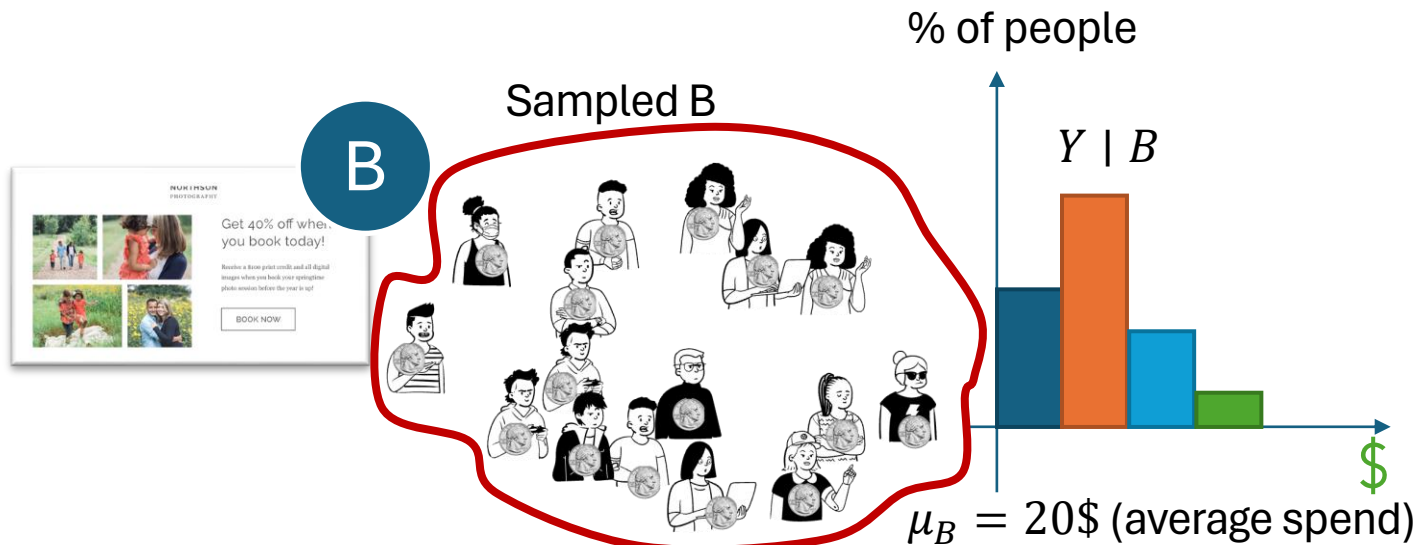
with probability 95%

population mean

$$\mu_A^0 \approx \mu_A \pm 2 \sqrt{\frac{\text{Var}(Y|A)}{N_A}}$$

sample mean

sampling error



population mean

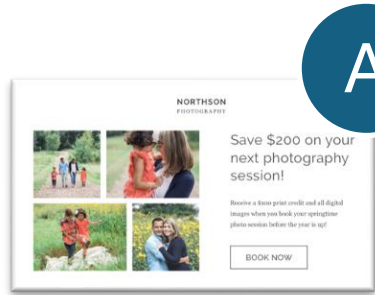
$$\mu_B^0 \approx \mu_B \pm 2 \sqrt{\frac{\text{Var}(Y|B)}{N_B}}$$

sample mean

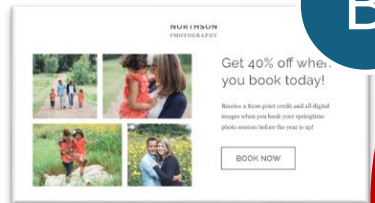
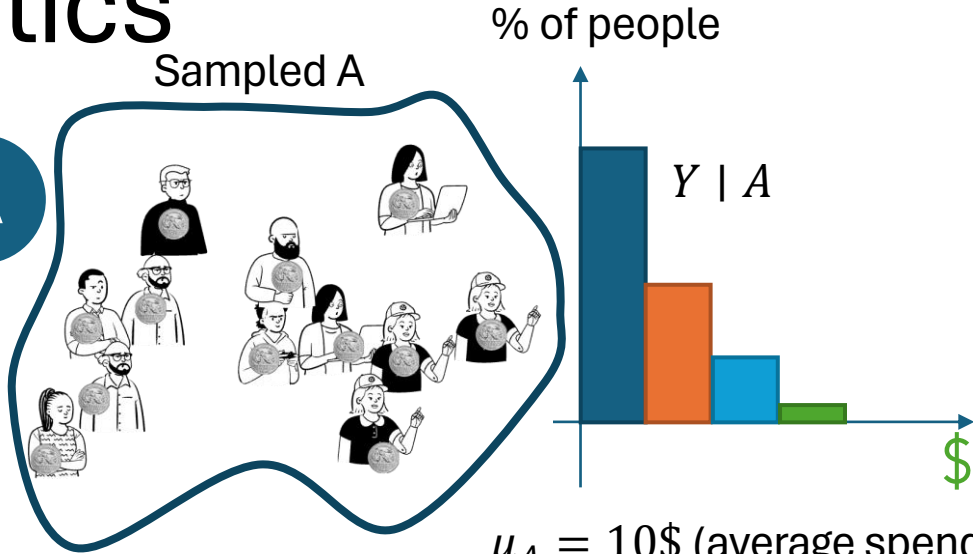
sampling error

Statistics

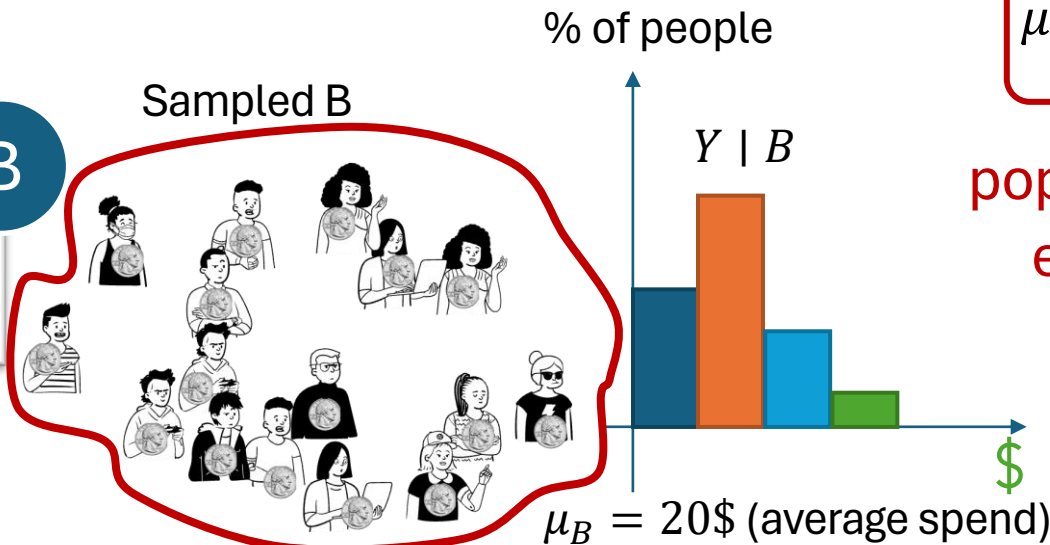
with probability 95%



A



B



$$\mu_B^0 - \mu_A^0$$

population effect

$$\approx \mu_B - \mu_A \pm$$

sample effect

$$2 \sqrt{\frac{\text{Var}(Y|A)}{N_A} + \frac{\text{Var}(Y|B)}{N_B}}$$

sampling error

Experimentation Platform

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Experimentation Platform (ExP) is a team of 60+ Data Scientists, Software Engineers and Program Managers. Our mission is to accelerate innovation through trustworthy experimentation. Most major products such as Bing, Cortana, Edge, Exchange, Identity, MSN, Office client, Office online, Photos, Skype, Speech, Store, Teams, Visual Studio Code, Windows, Xbox use our platform ExP to run trustworthy Online Controlled Experiments – aka A/B tests.



Interference!
The Big Challenge of A/B Testing in
Markets and Platforms



Interference

- Social Network interference
- Equilibrium effects
- Stateful systems and time effects

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Advertise on Reddit

Reach over 100K communities – Connect with passionate communities that deliver results for brands across all industries. Create impact & own top communities in your target category for 24 hours. Try Reddit **ads**.

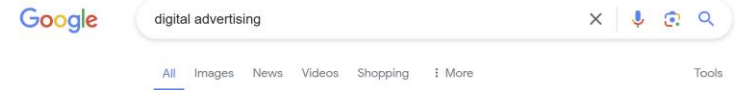
$$v_1 \sim F_1$$

 **Microsoft**
<https://about.ads.microsoft.com> > advertising > start-now

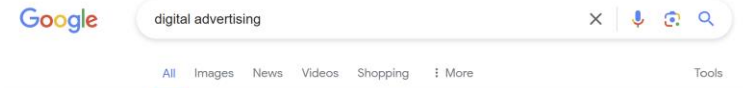
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We'll Help You Find Your Customers and Reach Searchers Across The Microsoft Network. Plus, Receive a \$500 Microsoft **Advertising** Credit When You Spend Just \$250! Free Sign Up.

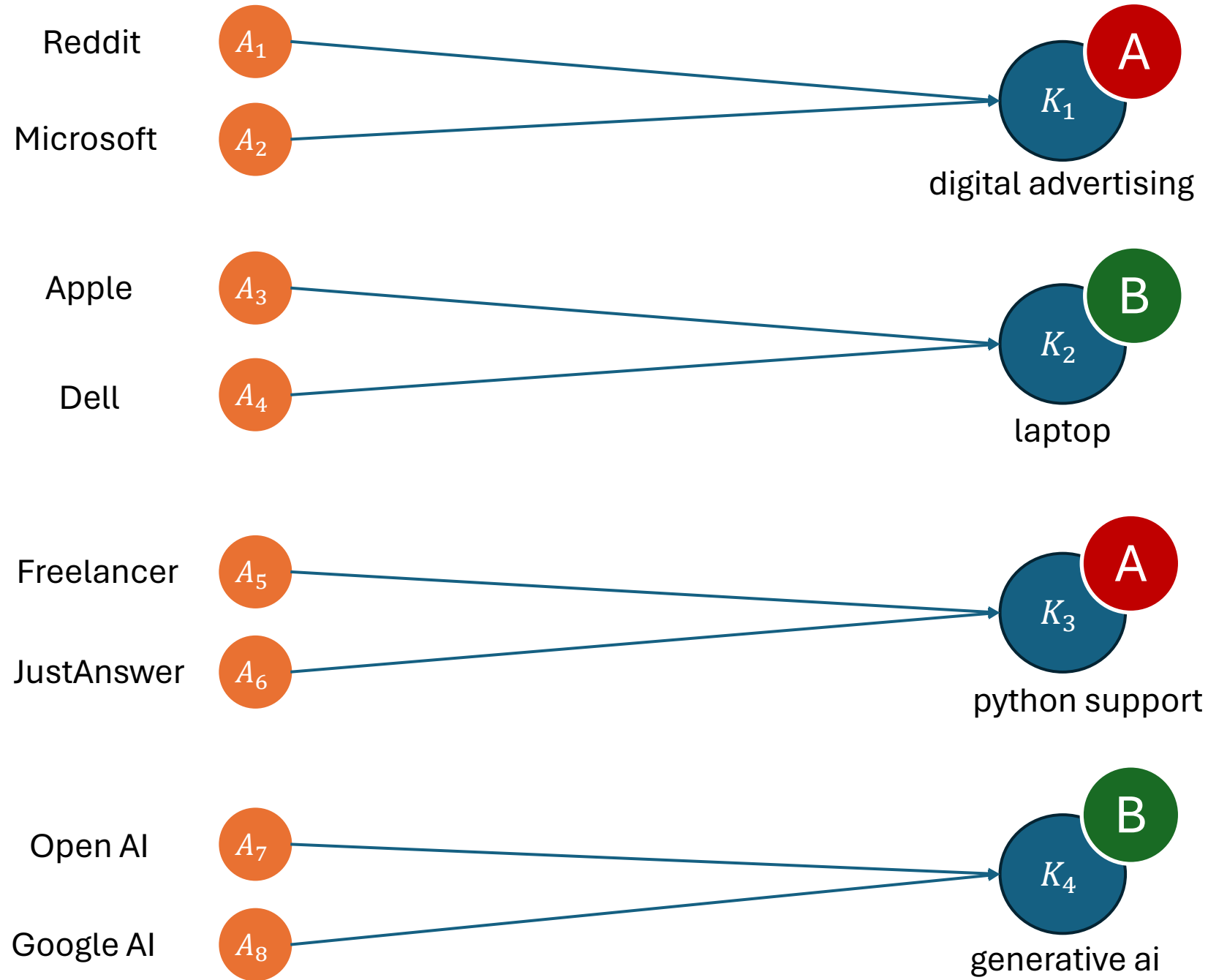
$$v_2 \sim F_2$$

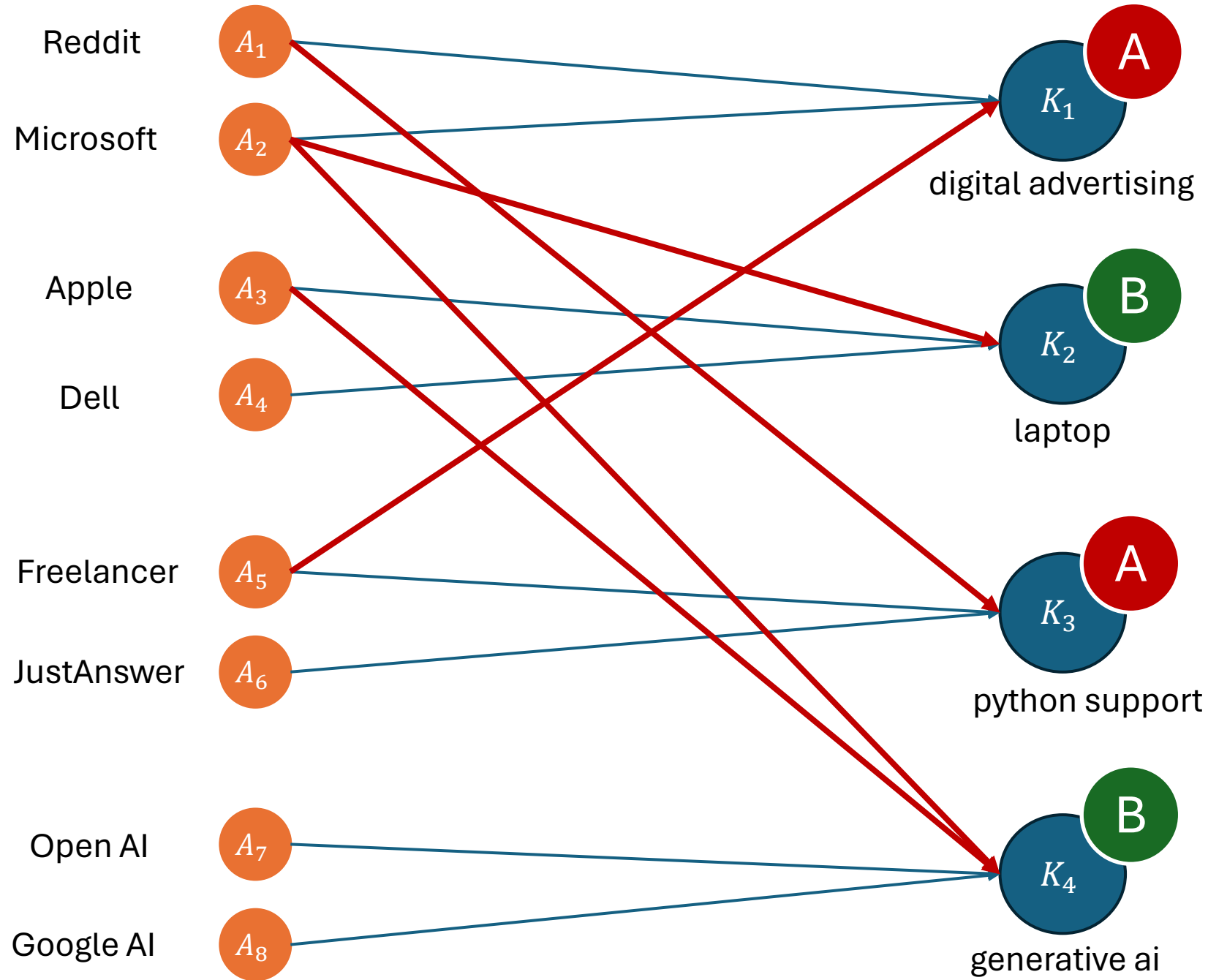


A



B





About 15,700,000 results (0.27 seconds)

Wool Socks at Walmart
www.walmart.com/Underwear_&_Socks
 4.3 ★★★★★ rating for walmart.com
 Save On Wool Socks at Walmart. Free Shipping Site to Store.

Wool Socks Superstore - SocksAddict.com
www.socksaddict.com/Wool-Socks
 4.5 ★★★★★ rating for socksaddict.com
 Get Free S&H + 99% Ship Same Day. SmartWool, Darn Tough & More Wool!
 Guaranteed lowest prices - 99% ship same day - 180 day return policy
 Brands: Outdoor Series, Running Series, Kids Series, Darn Tough, ...
 Injinji Wool Socks - Wigwam Wool Socks - Women's Wool Socks - Shop All Socks

Official Site: SmartWool® Clothing & Socks
www.smartwool.com/ SmartWool
 Extraordinarily comfortable SmartWool® socks & apparel: High performance Merino
 hiking, skiing, outdoor sport, running, walking, cycling & daily clothing.
 Hike - Run - Baselayers - Socks







Amazon.com: Merino Wool Blend Socks: Clothing
www.amazon.com/Gilbin...
 3 Pairs Merino Wool Blend
 Merino Wool Blend Socks, perfect for wearing with hiking, hunting, skiing, ...

SmartWool Socks at Sierra Trading Post
www.sierratradingpost.com/smartwool-socks-bs-19440-231/
 Improving upon nature's finest insulator, combine the best Merino wool and the latest
 advanced technology and SmartWool is born. It is the ultimate wool for ...

REI Lightweight Merino Wool Hiking Crew Socks at REI.com
www.rei.com/.../rei-lightweight-merino-wool-hiking-crew-socks REI
 ★★★★★ Rating: 5 - Review by stringbreaker - Mar 21, 2014 - \$12.50

Shopping Ads

Shop for wool socks on Google Sponsored

 Merino Wool Ragg Sock, 12" Two-Pack... \$22.95 L.L. Bean	 SmartWool Men's New Classic Rib So... \$18.95 Socks Addict	 carhartt men's all terrain wool blend s... \$10.50 Sears
 Hue Cable Knit Merino Wool Blend... \$10.00 Bloomingdale's	 Dickies Men's 1pk Merino Wool Socks... \$11.00 Target	 Men's Over-the-calf Wool Rib Dress Soc... \$9.99 Lands' End

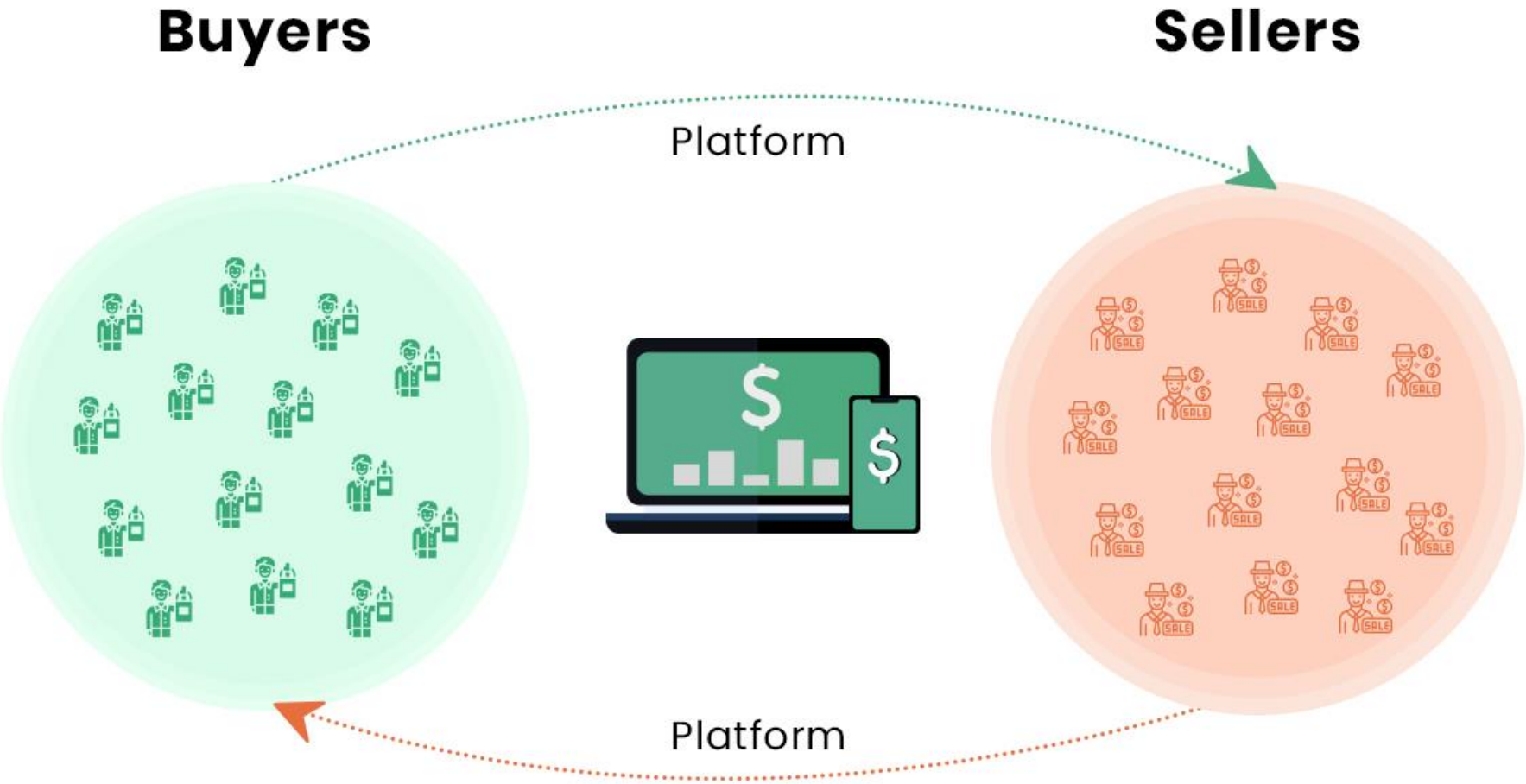
Search Network Ads

Alpaca Therapeutic Socks
www.alpacasofmontana.com/
 Warm, Soft, Alpaca Diabetic Socks
 Non-restrictive promotes blood flow

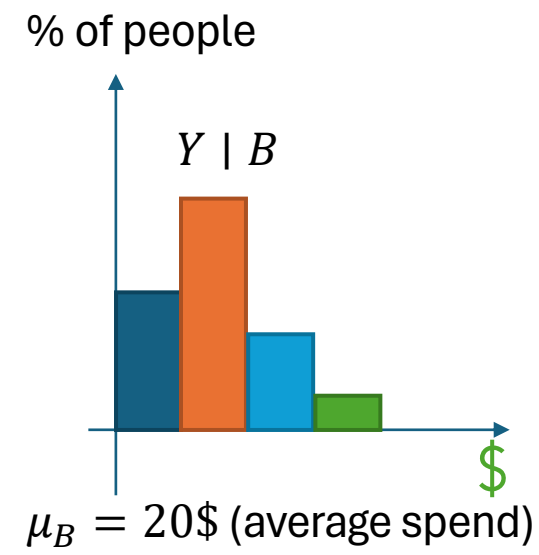
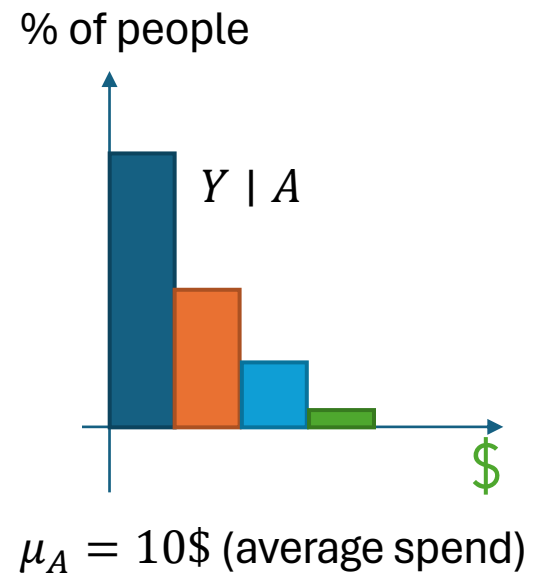
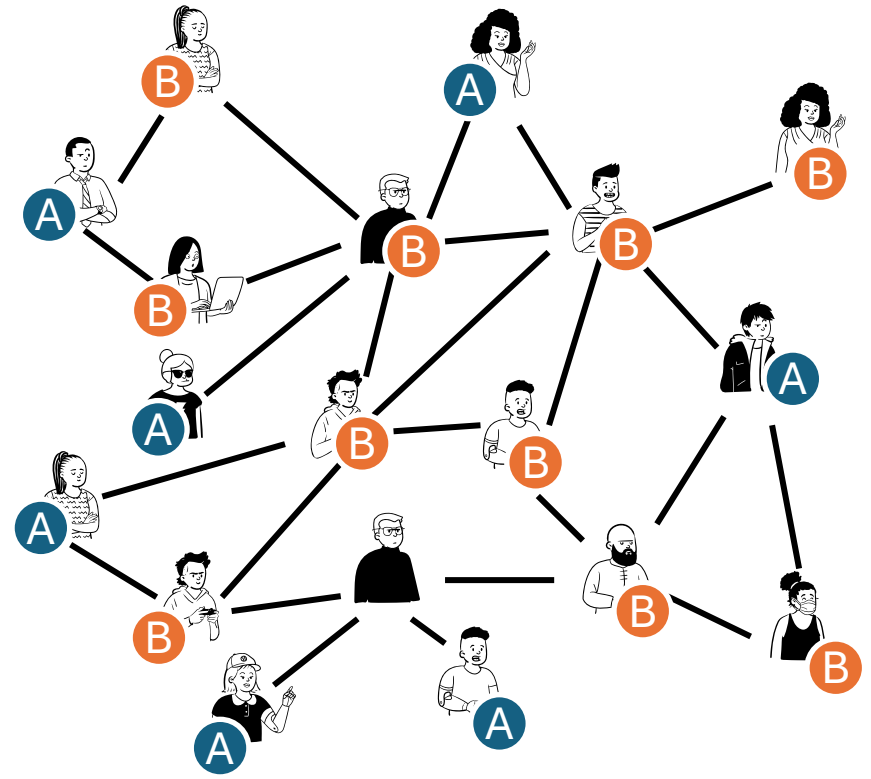
Women's SmartWool Socks
www.sierratradingpost.com/SmartWool
 4.6 ★★★★★ advertiser rating
 Great Selection of SmartWool Socks.
 Shop SmartWool Socks For Women Now.

Image source: https://googleadstrategy.com/google-adwords-search-network-vs-display-network/

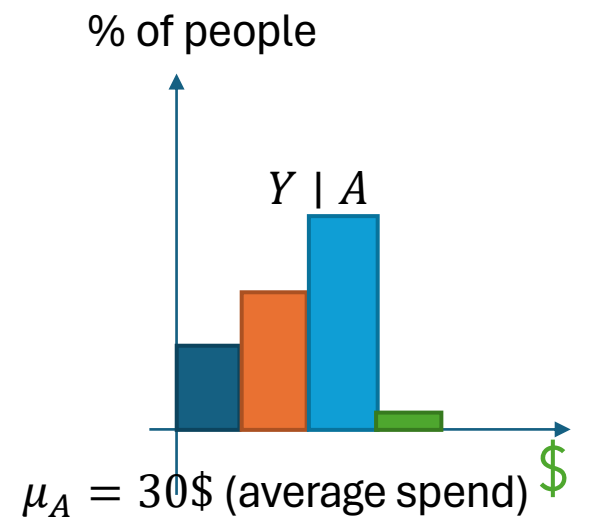
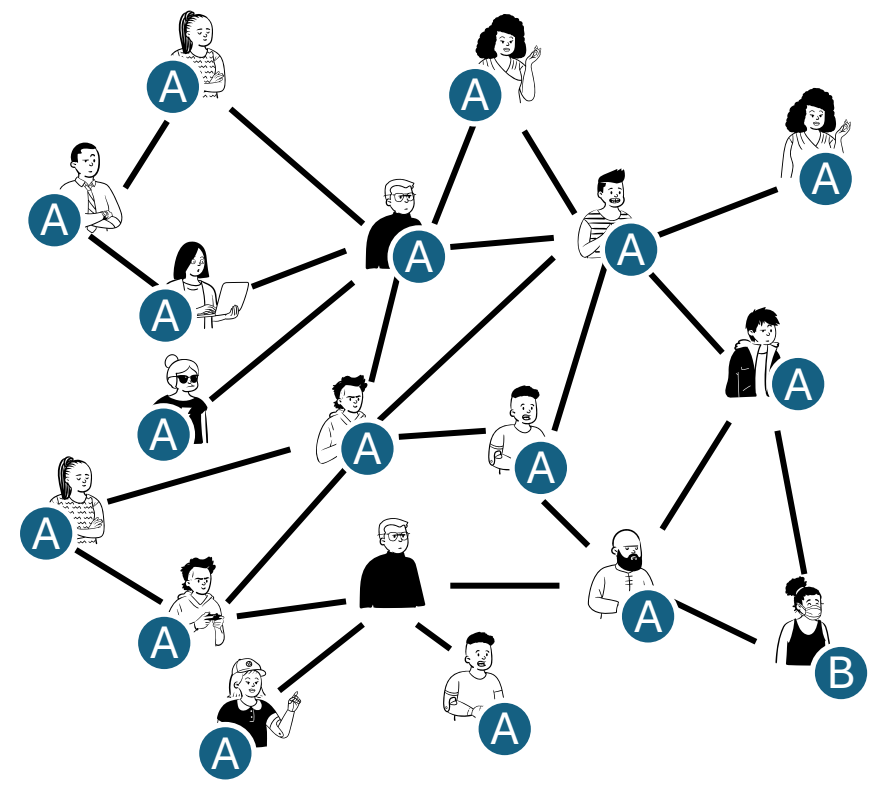
Two-Sided Matching Markets



Social Network Interference

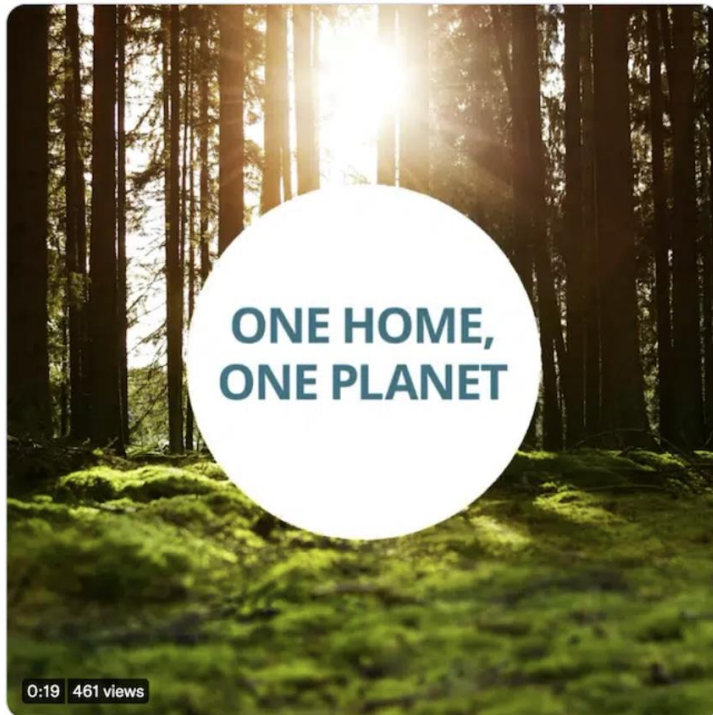


Counterfactual world





Big challenges call for big solutions. Tune in to [#OHOP21](#) on 9 November to hear thinkers, doers and leaders discuss the global response to climate change. Watch the event, get inspired and discover how we can take action [ingka.com/one-home-one-p...](#)
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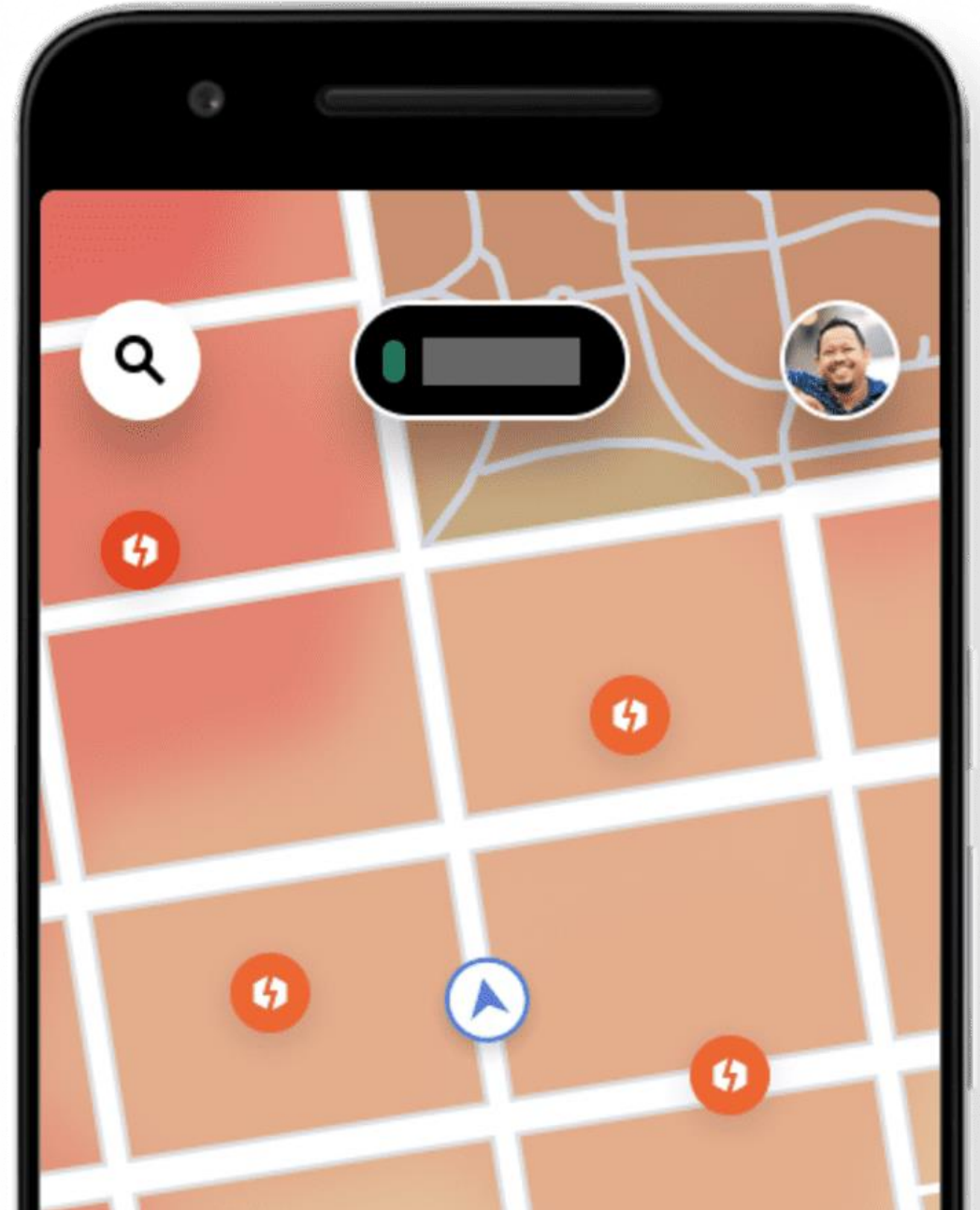
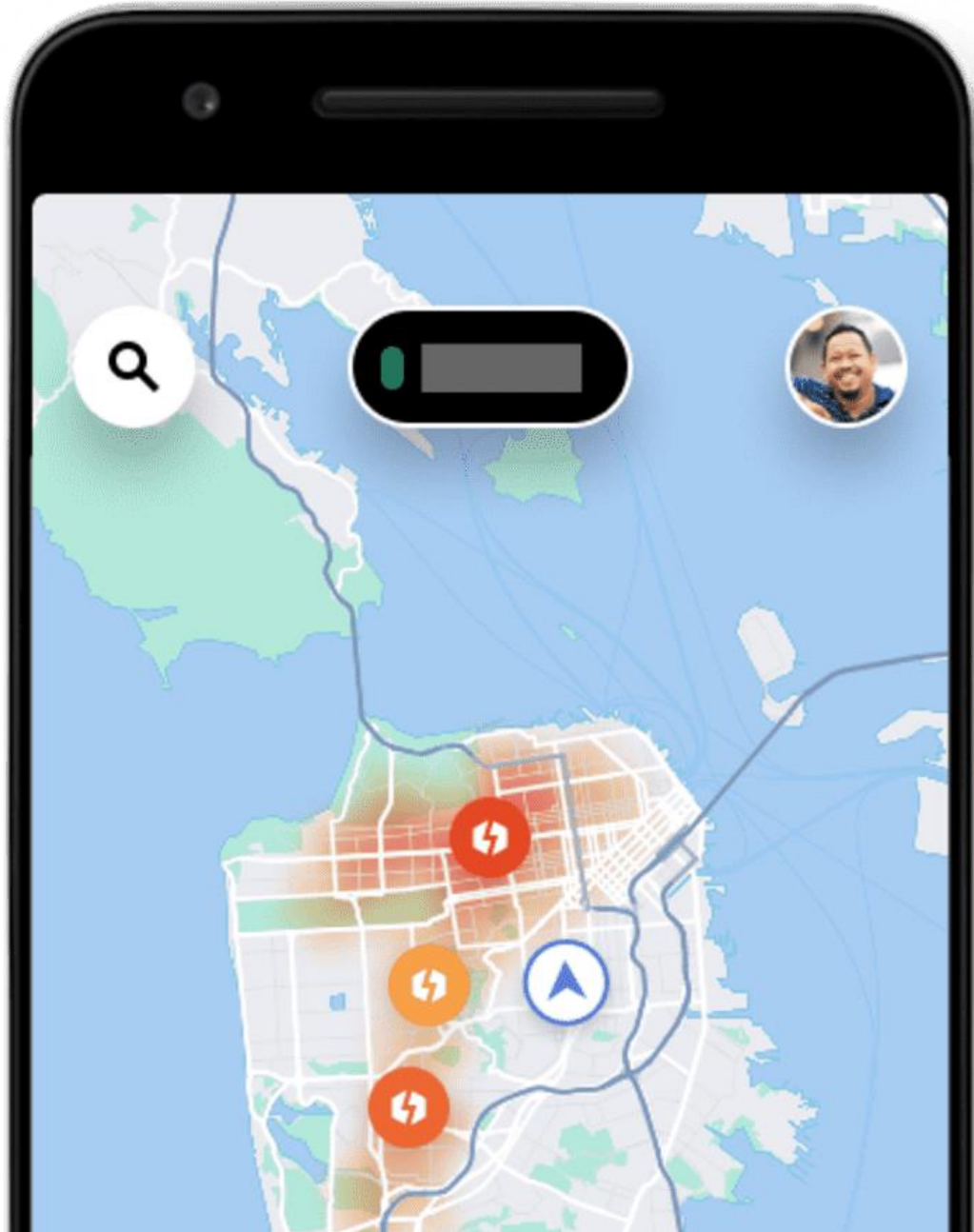
Source: [@IKEA](#)



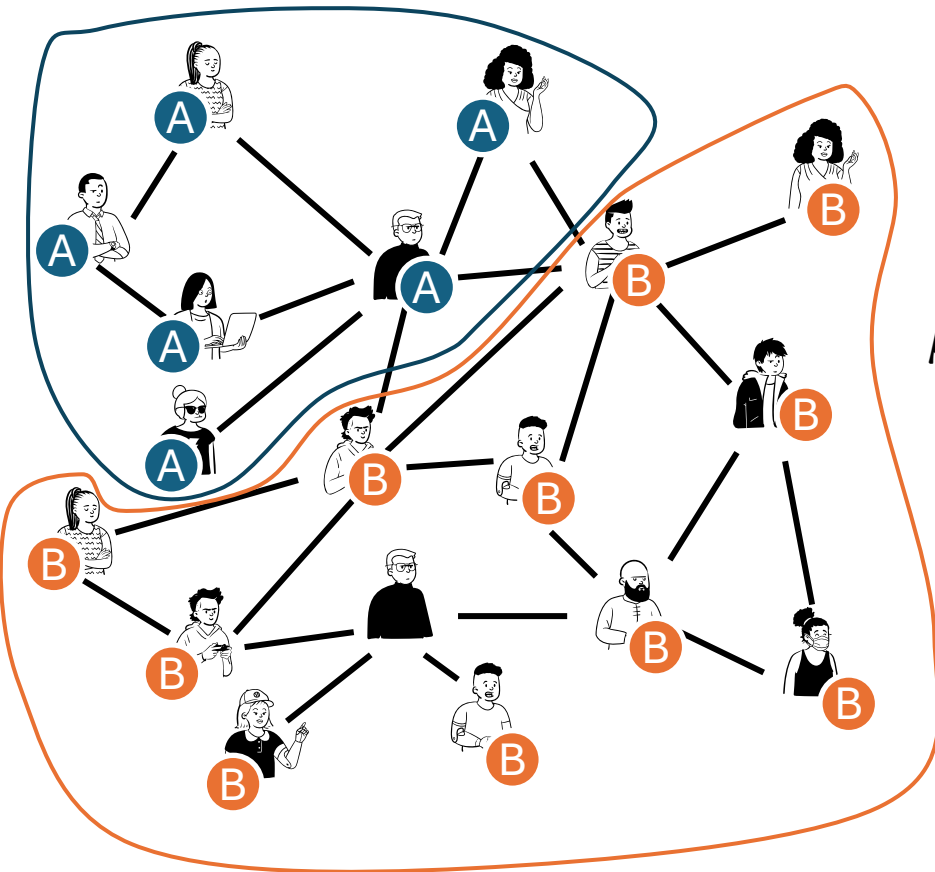
The moment is now. Climate action can't wait any longer. Join global thinkers, doers & leaders at [#OHOP21](#) on 9 Nov – where they'll discuss the need for urgent change & action to help create a better future. Learn more: [ingka.com/one-home-one-p...](#)
[#COP26](#) [#AssembleABetterFuture](#)



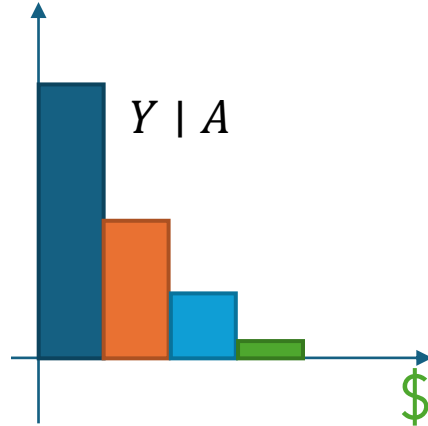
Source: [@IKEA](#)



Approach: Clustering

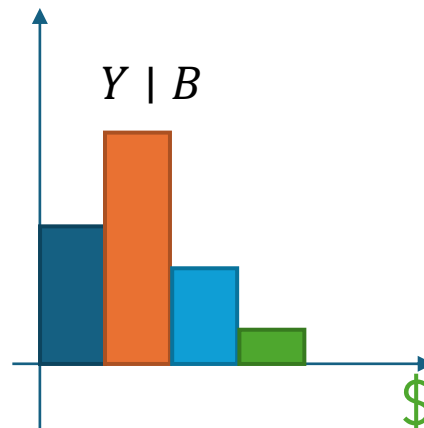


% of people



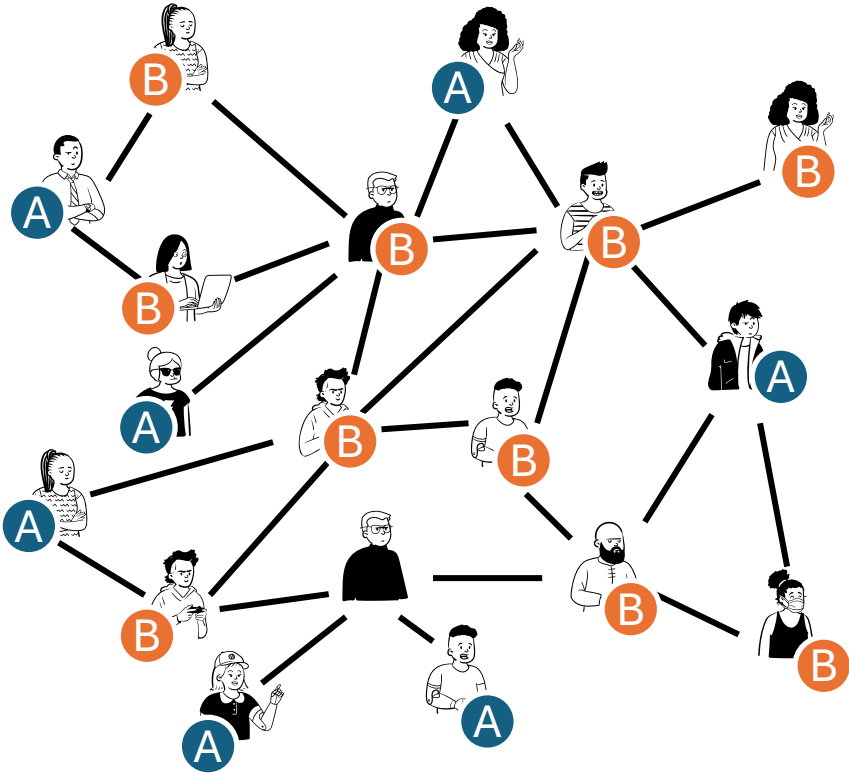
$\mu_A = 10\$$ (average spend)

% of people

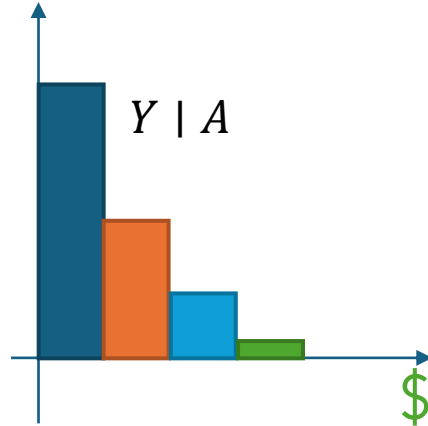


$\mu_B = 20\$$ (average spend)

Approach: Structural Bias Correction



% of people

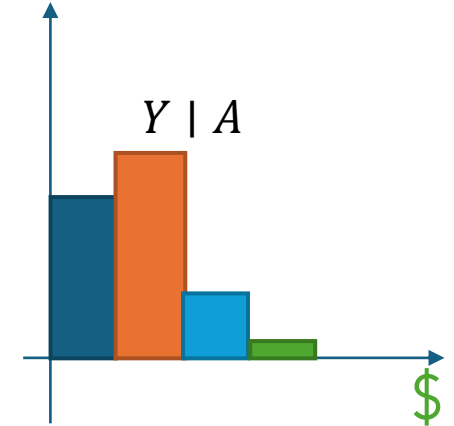


$\mu_A = 10\$$ (average spend)

Correct
Spill-Over
Bias

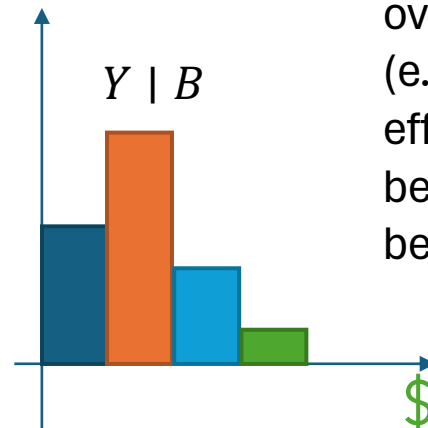


% of people



$\mu_A = 10\$$ (average spend)

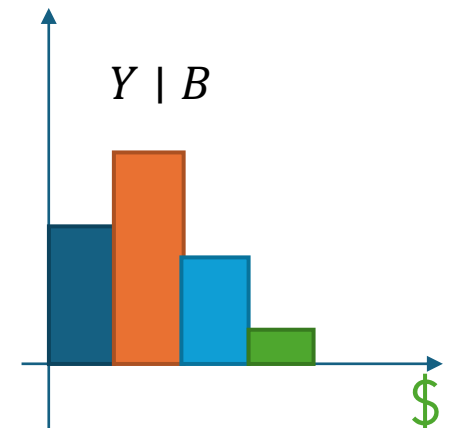
% of people



$\mu_B = 20\$$ (average spend)

Use Network Information +
Assumptions on how spill-
overs change outcome
(e.g. additive homophily
effects, market equilibrium
behavior, Nash equilibrium
behavior)

% of people



$\mu_B = 20\$$ (average spend)

A/B Testing in Auctions

A/B Testing over Position Auction Formats

Context A/B Testing for Position Auctions

- We want to optimize over the space of position auctions
- We are allowed to play with the click probabilities of slots
 - By reducing or increasing the space allocated to each slot
 - Changing the probability that the slot appears on the impression
 - Randomizing which slot the k -th highest bidder gets

High-Level Idea

- We will see that we can run a single randomized auction
- Using data that contain m samples of bids from that single randomized auction, we can estimate the revenue for every other auction in the design space at an estimation rate of $\frac{1}{\sqrt{m}}$
- Hence, we can choose the best auction in the space, with only a few rounds of experimentation!
- To do that we will need to use optimal auction theory!

Formal Setting

- We have N bidders and N slots (wlog) with CTRs $a_1 \geq \dots \geq a_N \geq a_{N+1} = 0$
- Bidders are charged their bid-per-click (GFP)
- k -th highest bidder assigned with some distribution to one of the slots
- Slot distributions are solely determined by bid rank
- k -th highest bidder gets an implicit expected CTR of x_k

$$x_k = p_{k1}a_1 + \dots + p_{kN}a_N$$

- These expected CTRs are monotone decreasing, $x_1 \geq x_2 \geq \dots \geq x_N$
- No bidder is over-assigned $\sum_j p_{kj} \leq 1$
- No slot is over-assigned $\sum_k p_{kj} \leq 1$

Feasibility Characterization

- They must be feasible: for each prefix, x_1, \dots, x_k I cannot allocate a total probability more than the cumulative top k highest slots

$$\sum_{i=1}^k x_i \leq \sum_{i=1}^k a_i$$



Feasibility Characterization

- They must be feasible: for each prefix, x_1, \dots, x_k I cannot allocate a total probability more than the cumulative top k highest slots

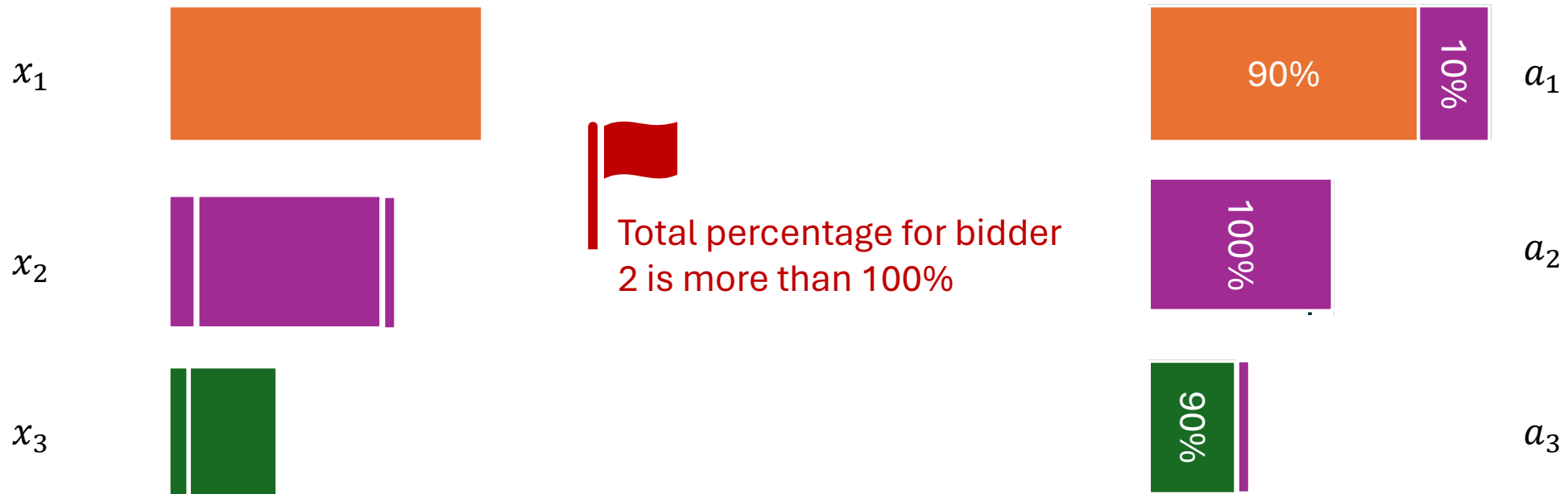
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Feasibility Characterization

- They must be feasible: for each prefix, x_1, \dots, x_k I cannot allocate a total probability more than the cumulative top k highest slots

$$\sum_{i=1}^k x_i \leq \sum_{i=1}^k a_i$$



Equivalently: Position Auction with Flexible CTRs

- We have N bidders and N slots
- Bidders are charged their bid-per-click (GFP)
- Slots are assigned in decreasing order of bidders
- k -th slot has CTR x_k . CTR of k -th slot is part of the design space
- Can choose the CTRs in any manner that satisfies $\forall k \leq N$:

$$\sum_{j=1}^k x_j \leq \sum_{j=1}^k a_j$$

for some set of predefined quantities $a_1 \geq \dots \geq a_N \geq a_{N+1} = 0$

Equivalently: Distribution over k -Unit Auctions

- In a k -unit auction we are selling k -units of the same good
- The top- k bidders win a unit and pay their bid

Theorem. Position auction with $x_1 \geq \dots \geq x_N \geq x_{N+1} = 0$, equivalent to distribution over k -unit auctions. k -th unit auction chosen w.p.

$$w_k = x_k - x_{k+1}, \quad k \geq 1, \quad \text{and,} \quad w_0 = 1 - x_1$$

Proof. If you are the j -th bidder in position auction, you win w.p. x_j
If you are the j -th bidder in random k -unit auction, you win if $k \geq j$

$$\Pr(k \geq j) = \sum_{k \geq j} w_k = \sum_{k \geq j} x_k - x_{k+1} = x_j$$

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If you are the j -th bidder in random k -unit auction, you win if $k \geq j$

$$\Pr(k \geq j) = \sum_{k \geq j} w_k = \sum_{k \geq j} x_k - x_{k+1} = x_j$$

Equivalently: Distribution over k -Unit Auctions

- In a k -unit auction we are selling k -units of the same good
- The top- k bidders win a unit and pay their bid
- We run k -unit auction with probability w_k
- When bidders are symmetric, every such auction has a symmetric monotone equilibrium (in fact it has a unique equilibrium that is symmetric and monotone)

Revenue of Randomized k -unit Auction

- By Myerson, revenue of any auction is expected virtual welfare

$$\text{Rev} = \sum_i E[\phi_i(v_i) \cdot x_i(v_i)] = \sum_i \sum_k w_k E[\phi_i(v_i) \cdot x_{i,k}(v_i)]$$

- Allocation function is solely determined by rank

$$\begin{aligned} x_{i,k}(v) &= \Pr(\leq k - 1 \text{ bidders above you}) \\ &= \sum_{t=1}^{k-1} \binom{n-1}{t} (1 - F(v))^t F(v)^{n-1-t} \end{aligned}$$

- Expected allocation only depends on quantile $q(v) = F(v)$

Revenue of Randomized k -unit Auction

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- Expected allocation only depends on quantile $q(v) = F(v)$
- Convenient to re-express everything in quantiles instead of values

Revenue of Randomized k -unit Auction

- By Myerson, revenue of any auction is expected virtual welfare

$$\text{Rev} = \sum_i E[\phi_i(q_i) \cdot x_i(q_i)] = \sum_i \sum_k w_k E[\phi_i(q_i) \cdot x_{i,k}(q_i)]$$

- Allocation function is solely determined by rank

$$x_{i,k}(q) = \sum_{t=1}^{k-1} \binom{n-1}{t} (1-q)^t q^{n-1-t}$$

- Quantiles q are uniformly distributed in $[0,1]$:

$$v(q) = F^{-1}(q), \quad \Pr(Q \leq q) = \Pr(v \leq v(q)) = F(v(q)) = q$$

- Virtual values simplify, since by derivative of inverse $v'(q) = (F^{-1}(q))' = 1/f(v(q))$

$$\phi_i(q) = v(q) - \frac{1 - F(v(q))}{f(v(q))} = v(q) - (1 - q) \cdot v'(q) = -(v(q) \cdot (1 - q))'$$

Suffices to Analyze Estimation of Revenue of k -th unit Auction

- The revenue is the weighted sum of terms (using also symmetry)

$$R_k = E[\phi(q) \cdot x_k(q)]$$

- The function $x_k(q)$ is known in closed form
- The function $\phi(q)$ is negative derivative of the **revenue function**

$$\phi(q) = -R'(q), \quad R(q) = v(q) \cdot (1 - q)$$

- Integration-by-Parts yields

$$E[\phi(q) \cdot x_k(q)] = - \int_0^1 R'(q) \cdot x_k(q) dq = \int_0^1 R(q) \cdot x'_k(q) dq = E[R(q) \cdot x'_k(q)]$$

- It suffices that we estimate terms

$$R_k := E[v(q) \cdot (1 - q) \cdot x'_k(q)]$$

For any randomized k-unit first-price auction among symmetric bidders, we have that:

$$\text{Rev} = n \sum_{k \leq N} w_k E[v(q) \cdot (1 - q) \cdot x'_k(q)]$$

Estimating $R_k = E[v(q) \cdot (1 - q) \cdot x'_k(q)]$

- The value function $v(q) = F^{-1}(q)$ relates to distribution of values
- Only observed from data distribution of bids with CDF G and pdf g
- Define the bid function $b(q) = G^{-1}(q)$: *what is my bid if I'm at the bottom q -th percentile of the distribution of values, equivalently, if I'm at the q -th percentile of the distribution of bids*
- Want to relate value of quantile q to bid of quantile q
- Similar to bid inversion question in last lecture

Estimating $R_k = E[v(q) \cdot (1 - q) \cdot x'_k(q)]$

- At symmetric equilibrium

$$b(q) = \operatorname{argmax}_z (v(q) - z) \cdot x(b^{-1}(z))$$

- The first order condition (using derivative of inverse):

$$(v(q) - b(q)) \cdot x'(q) \frac{1}{b'(q)} - x(q) = 0$$

- We can write a similar bid inversion formula

$$v(q) = b(q) + \frac{b'(q)x(q)}{x'(q)}$$

- **Reminder:** The functions $x(q)$ and $x'(q)$ are known in closed form

Estimating $R_k = E[v(q) \cdot (1 - q) \cdot x'_k(q)]$

- We can write a similar bid inversion formula

$$v(q) = b(q) + \frac{b'(q)x(q)}{x'(q)}$$

- Need to estimate $b(q)$ and $b'(q)$ from data
- **Reminder:** $b(q) = G^{-1}(q)$, $b'(q) = \frac{1}{g(G^{-1}(q))}$
- Estimating $b(q)$ and $b'(q)$ is the same as estimating G, g
- **Main message.** The quantity R_k for any k depends only on $b(q)$ and not on $b'(q)$ because it is an integral over q ! Leads to much faster rates.

Estimating $R_k = E[v(q) \cdot (1 - q) \cdot x'_k(q)]$

- We can write

$$R_k = E[b(q) \cdot (1 - q) \cdot x'_k(q)] + E\left[\frac{b'(q)x(q)}{x'(q)} \cdot (1 - q) \cdot x'_k(q)\right]$$

- First part only depends on $b(q)$. Analogous to estimating a CDF
- Second part seemingly problematic. But **integration-by-parts**

$$E\left[\frac{b'(q)x(q)}{x'(q)} \cdot (1 - q) \cdot x'_k(q)\right] = -E\left[b(q) \left(\frac{x(q)(1 - q) \cdot x'_k(q)}{x'(q)}\right)'\right]$$

- This only depends on $b(q)$ and known quantities

For any randomized k-unit first-price auction among symmetric bidders, we have that:

$$\text{Rev} = n \sum_{k \leq N} w_k E[b(q) \cdot f(q)]$$

for a function $f(q)$ known in closed form

We can estimate Rev by estimating the CDF of bids using the empirical CDF \hat{G} . Then use $\hat{b} = \hat{G}^{-1}$ and

$$\widehat{\text{Rev}} = n \sum_{k \leq N} w_k \int_0^1 \hat{b}(q) \cdot f(q) dq$$

for a function $f(q)$ known in closed form

Assuming $f(q)$ is bounded (e.g. holds if original auction chooses each k with positive probability), then

$$|\widehat{\text{Rev}} - \text{Rev}| \lesssim 1/\sqrt{m}$$

Conclusion

- Run a single randomized auction as our experimentation strategy
- Using data that contain m samples of bids from that single randomized auction, we can estimate the revenue for every other auction in the design space at an estimation rate of $\frac{1}{\sqrt{m}}$
- Hence, we can choose the best auction in the space, with only a few rounds of experimentation!
- To do that we used optimal auction theory!

A/B Testing across Many Keywords with Budgets

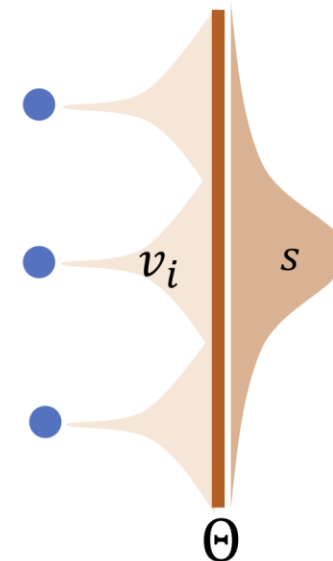
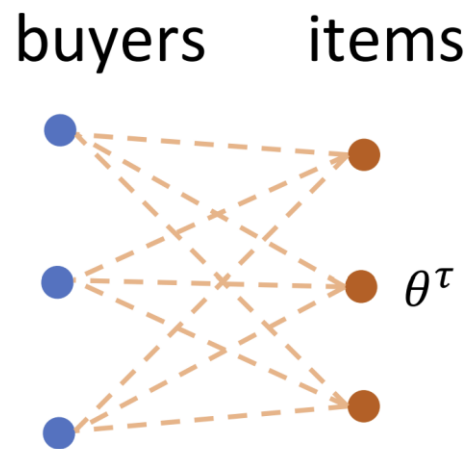
Budgets!

- So far we did not place any budget constraints on bidders
 - In practice, budget constraints are very important
 - Bidders participate in many auctions and have a budget limit
 - Can only spend at most B_i in total across all the auctions
-
- This couples the bidding strategy across auctions
 - Makes learning (e.g. no-regret learning hard)
 - In its full generality a stochastic dynamic program

Simplified Budgets: Pacing Equilibria

[Interference Among First-Price Pacing Equilibria: A Bias and Variance Analysis \(arxiv.org\)](#)

- In practice, people use the following simplification
- We have n bidders and a continuum of items
- Items have type θ which follows some distribution with measure s
- $v_i(\theta)$ is bidder i 's value for an item of type θ



Simplified Budgets: Pacing Equilibria

The multipliers $\beta = (\beta_1, \dots, \beta_n)$ and price function $p(\theta)$ are a *pacing equilibrium* if there exists an allocation function $x(\theta)$ such that

- First-price payment: $p(\theta) = \max_i \beta_i v_i(\theta)$
- Highest-bidder wins: $x_i(\theta) \geq 0 \Rightarrow \beta_i v_i(\theta) = \max_k \beta_k v_k(\theta)$

- Budgets are respected

$$\int_{\theta} x_i(\theta) p(\theta) s(\theta) d\theta \leq B_i$$

- No-overselling: $\sum_i x_i(\theta) \leq 1$
- Full-allocation of competitive items: $p(\theta) > 0 \Rightarrow \sum_i x_i(\theta) = 1$
- No un-necessary pacing: $\int_{\theta} x_i(\theta) p(\theta) s(\theta) d\theta < B_i \Rightarrow \beta_i = 1$

Characterization of Pacing Equilibria

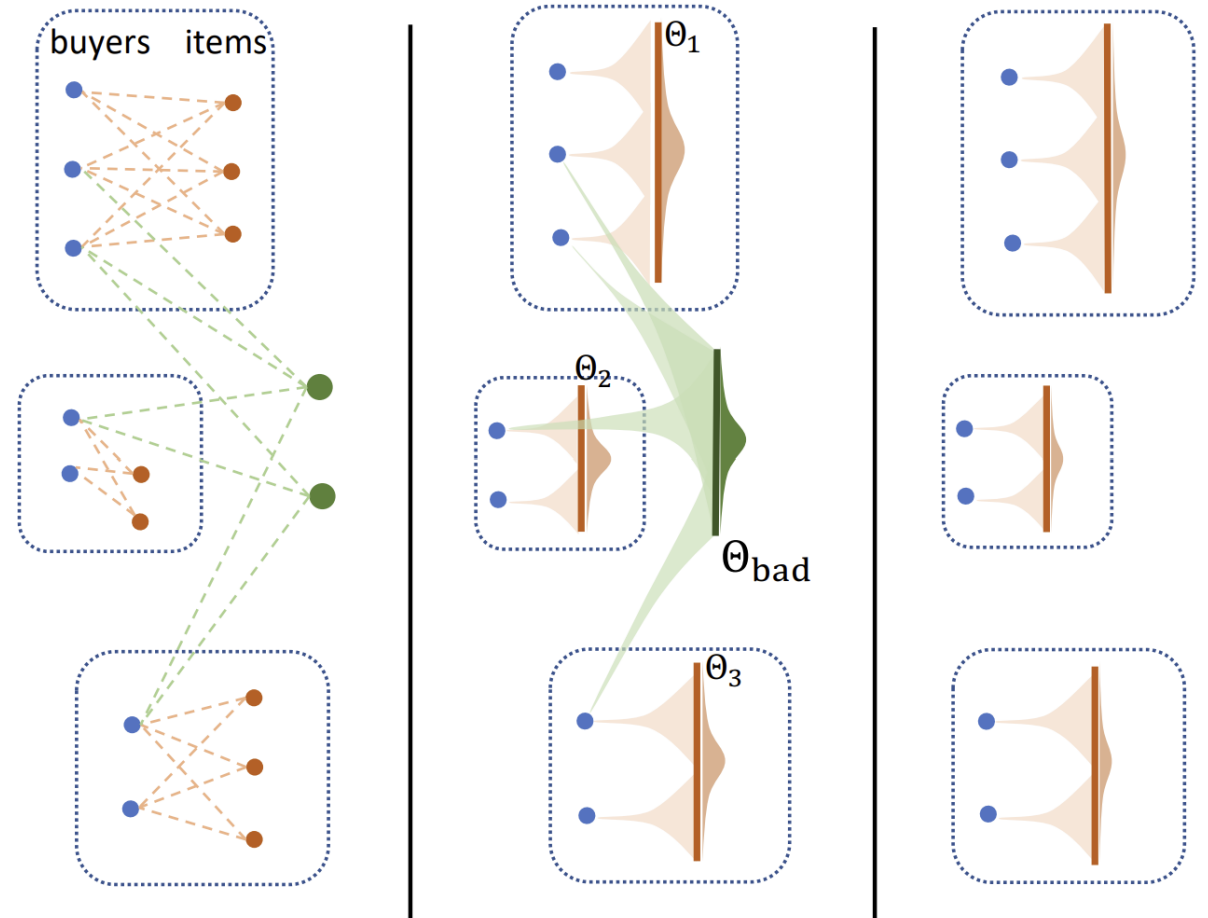
Multipliers in pacing equilibrium are characterized as solutions to a convex optimization problem (related to market equilibrium)

$$\beta_* = \operatorname{argmin}_{\beta \in (0,1]^n} E \left[\max_i \beta_i v_i(\theta) \right] - \sum_i B_i \log(\beta_i)$$

Clustered Experiment Designs and Debiasing

[Interference Among First-Price Pacing Equilibria: A Bias and Variance Analysis \(arxiv.org\)](#)

1. For each sub-market want pacing multipliers as if the bad items don't exist
2. With such multipliers, can estimate idealized revenue for each sub-market, as if isolated
3. Characterization of multipliers as minimizers of market equilibrium program \Rightarrow closed form first-order bias that bad items introduce
4. Subtract bias and measure revenue of A and B clusters using debiased multipliers



A/B Testing in Two-Sided Matching Markets

Two-Sided Randomized Designs

[Experimental Design in Two-Sided Platforms: An Analysis of Bias | Management Science \(informs.org\)](#)

