

# Fantastic Bugs and How to Squash Them; or, the Crimes of Solidity

Evan Sultanik

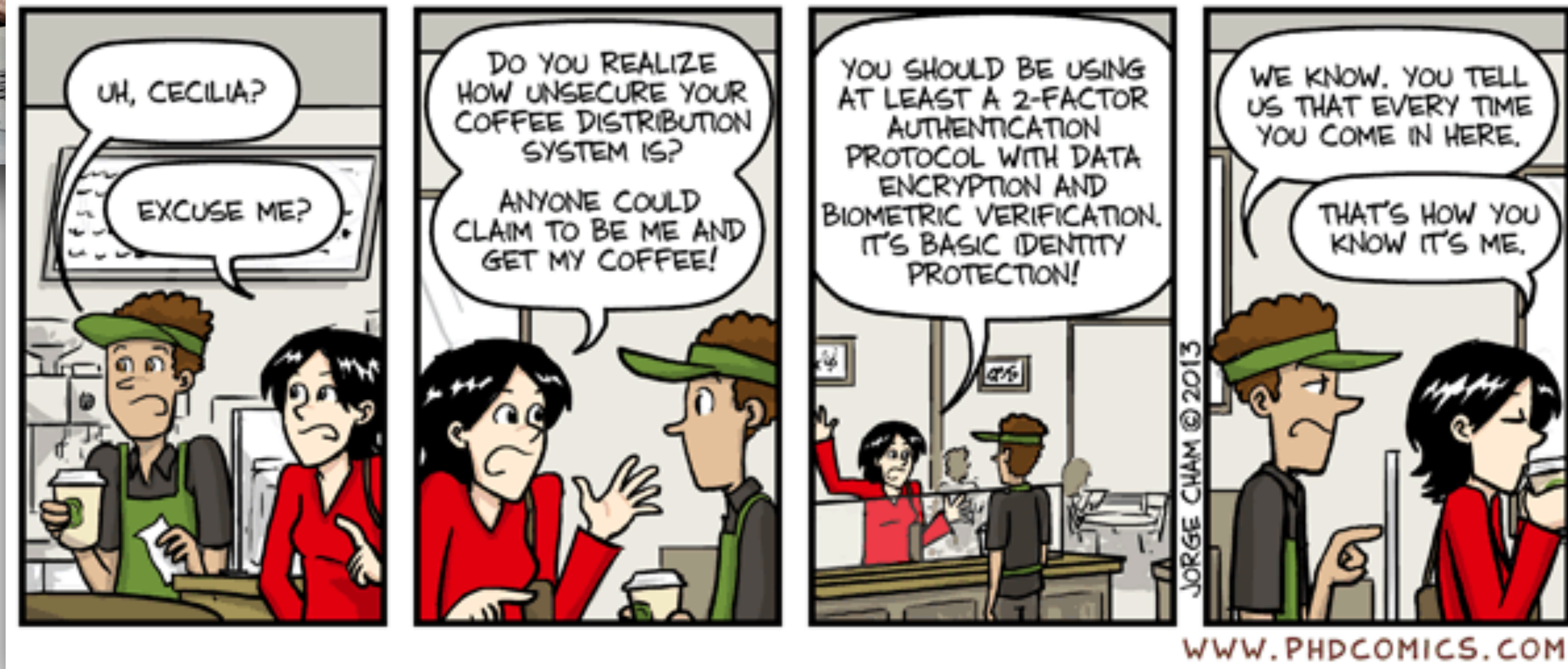
@ESultanik



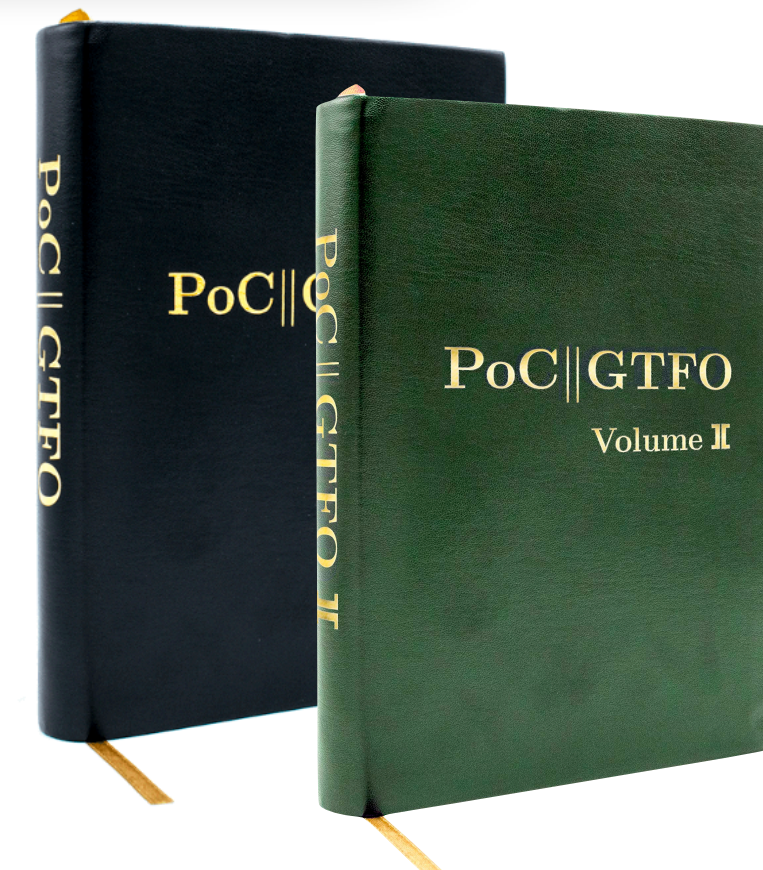




whoami`



TRAIL  
OF BITS





What should you take away  
from this talk?



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- Experienced Ethereum developers
  - ✓ Learn from the most common mistakes of your peers
  - ✓ Learn new tooling for improving your SDLC



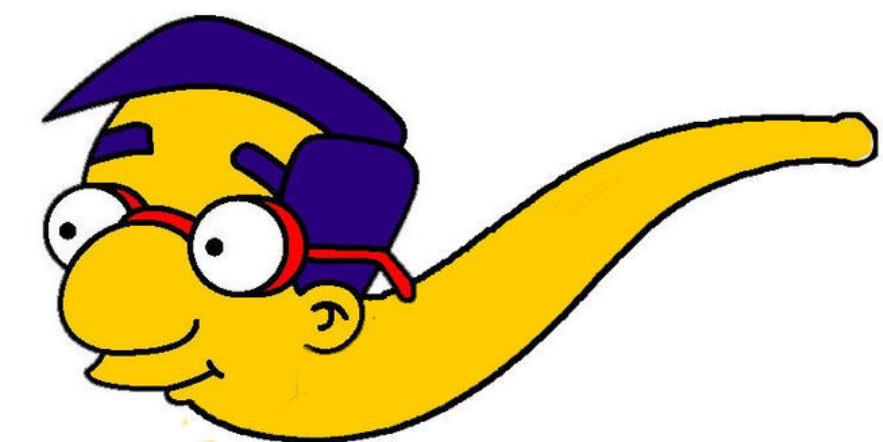
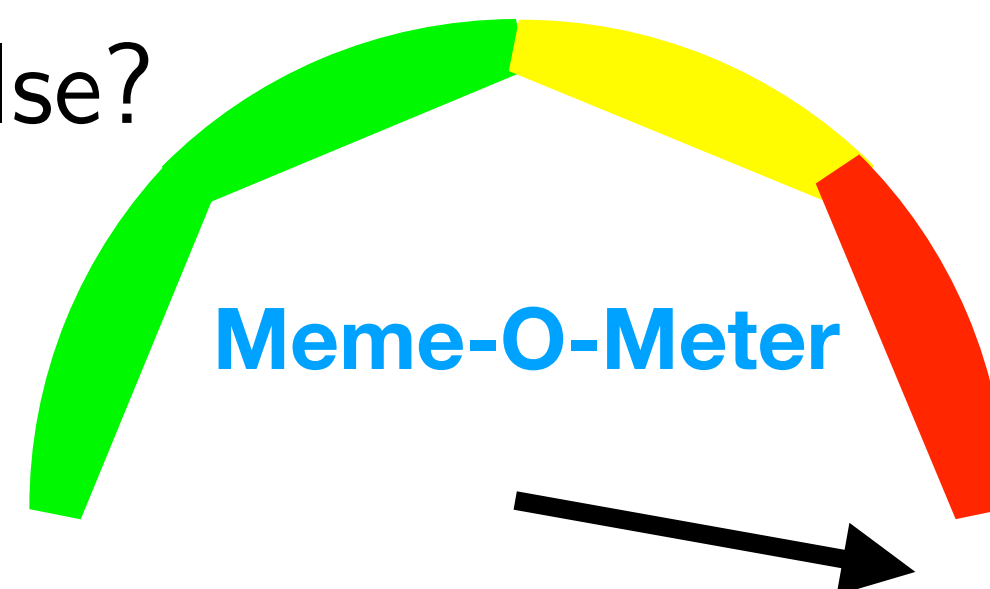
# What should you take away from this talk?

- Experienced Ethereum developers
  - ✓ Learn from the most common mistakes of your peers
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- Programmers who are new to smart contracts
  - ✓ Learn what *not* to do
  - ✓ Learn what *to do*



# What should you take away from this talk?

- Experienced Ethereum developers
  - ✓ Learn from the most common mistakes of your peers
  - ✓ Learn new tooling for improving your SDLC
- Programmers who are new to smart contracts
  - ✓ Learn what *not* to do
  - ✓ Learn what *to do*
- People interested in the technology
  - ✓ Learn about the state of the ecosystem
- Everyone else?



*Ceci n'est pas une meme*

# Outline

- Solidity the Language
- Solidity Implementation and Tooling
- On the Horizon
- Bugs!
- What You Can Do About It

*Hi, I'm Trippy, your programming assistant. I help you not get tripped up on Solidity.*



# Outline

- Solidity the Language
- Solidity Implementation and Tooling
- On the Horizon
- Bugs!
- What You Can Do About It

**It looks like you are  
trying to write a bug-  
free Solidity contract...**





# Solidity, the Language

by [Colin McMillen](#), [Jason Reed](#), and [Elly Jones](#).

<input type="checkbox"/>	functional	<input type="checkbox"/>	imperative	<input checked="" type="checkbox"/>	object-oriented	<input checked="" type="checkbox"/>	procedural	<input checked="" type="checkbox"/>	stack-based
<input type="checkbox"/>	"multi-paradigm"	<input type="checkbox"/>	lazy	<input type="checkbox"/>	eager	<input checked="" type="checkbox"/>	statically-typed	<input type="checkbox"/>	dynamically-typed

[X] non-programmer-friendly [1] completely incomprehensible

programming language Your language will not work. Here is why it will not work.

[ ] Syntax is what makes programming difficult

[X] Nobody really needs:

[ ] to interact with code not written in your language

[ ] Scaling up to large software projects will be easy

[ ] Convincing programmers to adopt a language-specific IDE will be easy

- [ ] Specifying behaviors as "undefined" means that programmers won't rely on them

[X] "Spooky action at a distance" makes programming more fun

You appear to believe that:

- ☐ Syntax is what makes programming difficult
- ☐ Garbage collection is free
- ☐ Computers have infinite memory
- ☒ Nobody really needs:
  - ☐ concurrency
  - ☐ a REPL
  - ☒ debugger support
  - ☐ IDE support
  - ☐ I/O
  - ☐ to interact with code not written in your language
- ☐ The entire world speaks 7-bit ASCII
- ☐ Scaling up to large software projects will be easy
- ☐ Convincing programmers to adopt a new language will be easy
- ☐ Convincing programmers to adopt a language-specific IDE will be easy
- ☐ Programmers love writing lots of boilerplate
- ☐ Specifying behaviors as "undefined" means that programmers won't rely on them
- ☒ "Spooky action at a distance" makes programming more fun

Unfortunately, your language (has/lacks):

- ☐ comprehensible syntax
- ☐ semicolons
- ☐ significant whitespace
- ☐ macros
- ☐ implicit type conversion
- ☐ explicit casting
- ☒ type inference
- ☐ goto
- ☐ exceptions
- ☒ closures
- ☐ tail recursion
- ☐ coroutines
- ☐ reflection
- ☒ subtyping
- ☐ multiple inheritance
- ☒ operator overloading
- ☐ algebraic datatypes
- ☒ recursive types
- ☐ polymorphic types
- ☐ covariant array typing
- ☒ monads
- ☐ dependent types
- ☐ infix operators
- ☐ nested comments
- ☐ multi-line strings
- ☒ regexes
- ☐ call-by-value
- ☐ call-by-name
- ☐ call-by-reference
- ☐ call-cc

The following philosophical objections apply:



The following philosophical objections apply:

- ☐ Programmers should not need to understand category theory to write "Hello, World!"
- ☐ Programmers should not develop RSI from writing "Hello, World!"
- ☐ The most significant program written in your language is its own compiler
- ☐ The most significant program written in your language isn't even its own compiler
- ☒ No language spec
- ☒ "The implementation is the spec"
  - ☐ The implementation is closed-source
  - ☐ covered by patents
  - ☐ not owned by you
- ☒ Your type system is unsound
- ☒ Your language cannot be unambiguously parsed
  - ☒ a proof of same is attached
    - ☐ invoking this proof crashes the compiler
- ☐ The name of your language makes it impossible to find on Google
- ☐ Interpreted languages will never be as fast as C
- ☐ Compiled languages will never be "extensible"
- ☐ Writing a compiler that understands English is AI-complete
- ☐ Your language relies on an optimization which has never been shown possible
- ☐ There are less than 100 programmers on Earth smart enough to use your language
- ☐ \_\_\_\_\_ takes exponential time
- ☐ \_\_\_\_\_ is known to be undecidable

Your implementation has the following flaws:

- ☐ CPUs do not work that way

Your implementation has the following flaws:

- [ ] CPUs do not work that way
- [ ] RAM does not work that way
- [ ] VMs do not work that way
- [X] Compilers do not work that way
- [ ] Compilers cannot work that way
- [ ] Shift-reduce conflicts in parsing seem to be resolved using `rand()`
- [ ] You require the compiler to be present at runtime
- [ ] You require the language runtime to be present at compile-time
- [X] Your compiler errors are completely inscrutable
- [X] Dangerous behavior is only a warning
- [ ] The compiler crashes if you look at it funny
- [ ] The VM crashes if you look at it funny
- [X] You don't seem to understand basic optimization techniques
- [X] You don't seem to understand basic systems programming
- [ ] You don't seem to understand pointers
- [ ] You don't seem to understand functions

Additionally, your marketing has the following problems:

- [ ] Unsupported claims of increased productivity
- [ ] Unsupported claims of greater "ease of use"
- [ ] Obviously rigged benchmarks
  - [ ] Graphics, simulation, or crypto benchmarks where your code just calls handwritten assembly through your FFI

Additionally, your marketing has the following problems:

- ☐ Unsupported claims of increased productivity
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- ☐ Obviously rigged benchmarks
  - ☐ Graphics, simulation, or crypto benchmarks where your code just calls handwritten assembly through your FFI
  - ☐ String-processing benchmarks where you just call PCRE
  - ☐ Matrix-math benchmarks where you just call BLAS
- ☐ Noone really believes that your language is faster than:
  - ☐ assembly
  - ☐ C
  - ☐ FORTRAN
  - ☐ Java
  - ☐ Ruby
  - ☐ Prolog
- ☐ Rejection of orthodox programming-language theory without justification
- ☐ Rejection of orthodox systems programming without justification
- ☐ Rejection of orthodox algorithmic theory without justification
- ☐ Rejection of basic computer science without justification

Taking the wider ecosystem into account, I would like to note that:

- ☐ Your complex sample code would be one line in: \_\_\_\_\_
- ☐ We already have an unsafe imperative language
- ☐ We already have a safe imperative OO language
- ☐ We already have a safe statically-typed eager functional language
- ☐ You have reinvented Lisp but worse
- ☒ You have reinvented Javascript but worse
- ☐ You have reinvented Java but worse
- ☐ You have reinvented C++ but worse
- ☐ You have reinvented PHP but worse



☐ assembly    ☐ C    ☐ FORTRAN    ☐ Java    ☐ Ruby    ☐ Prolog  
☐ Rejection of orthodox programming-language theory without justification  
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**[X] You have reinvented Javascript but worse**  
☐ You have reinvented Java but worse  
☐ You have reinvented C++ but worse  
☐ You have reinvented PHP but worse  
☐ You have reinvented PHP better, but that's still no justification  
☐ You have reinvented Brainfuck but non-ironically

In conclusion, this is what I think of you:

**[X] You have some interesting ideas, but this won't fly.**  
**[X] This is a bad language, and you should feel bad for inventing it.**  
**[X] Programming in this language is an adequate punishment for inventing it.**

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[X] Programming in this language is an adequate punishment for inventing it.

```
if (1 | 0 < 1) {  
    /* case 1 */  
} else {  
    /* case 2 */  
}
```



```
if (1 | 0 < 1) {
```

```
    C, C++, Javascript, Java, ...
```

```
} else {
```

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    /* case 2 */
```

```
}
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```
if (1 | 0 < 1) {
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C, C++, Javascript, Java, ...

```
} else {
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**LEEEROYYY JENKINS!**

Solidity

*Such Language!*

*Much Bugs!*

```
}
```

**WOW!**

```
if (1 | 0 < 1) {
```

C, C++, Javascript, Java, ...

```
} else {
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Solidity

**WOW!**

**LEEEROOYY**

Lesson: Don't assume  
Solidity behaves like  
most other languages!



```
for (var i = 0; i < foo.length; ++i) {  
    foo[i] = i;  
}
```

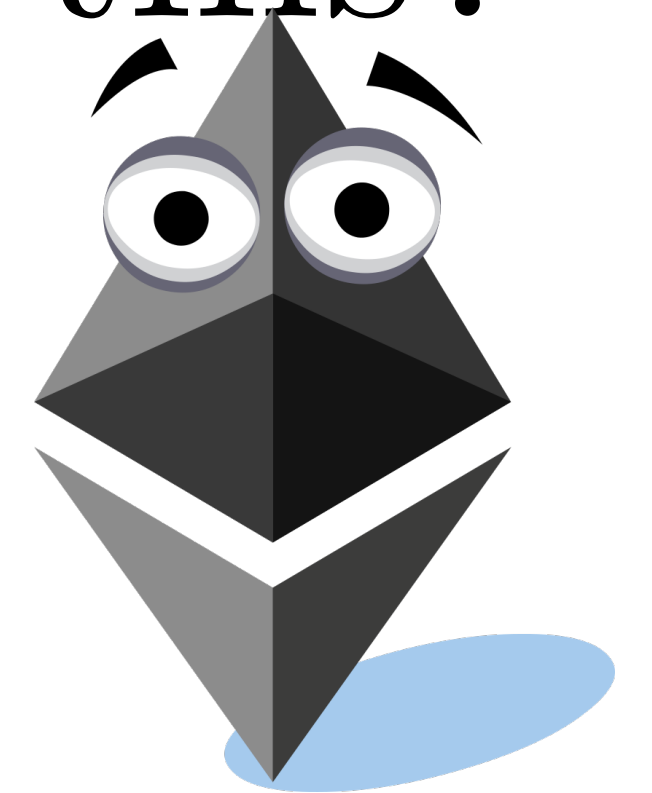
What does `foo[1337]` look like after this?



```
for (var i = 0; i < foo.length; ++i) {  
    foo[i] = i;  
}
```

*Lesson: Always use  
explicit types!*

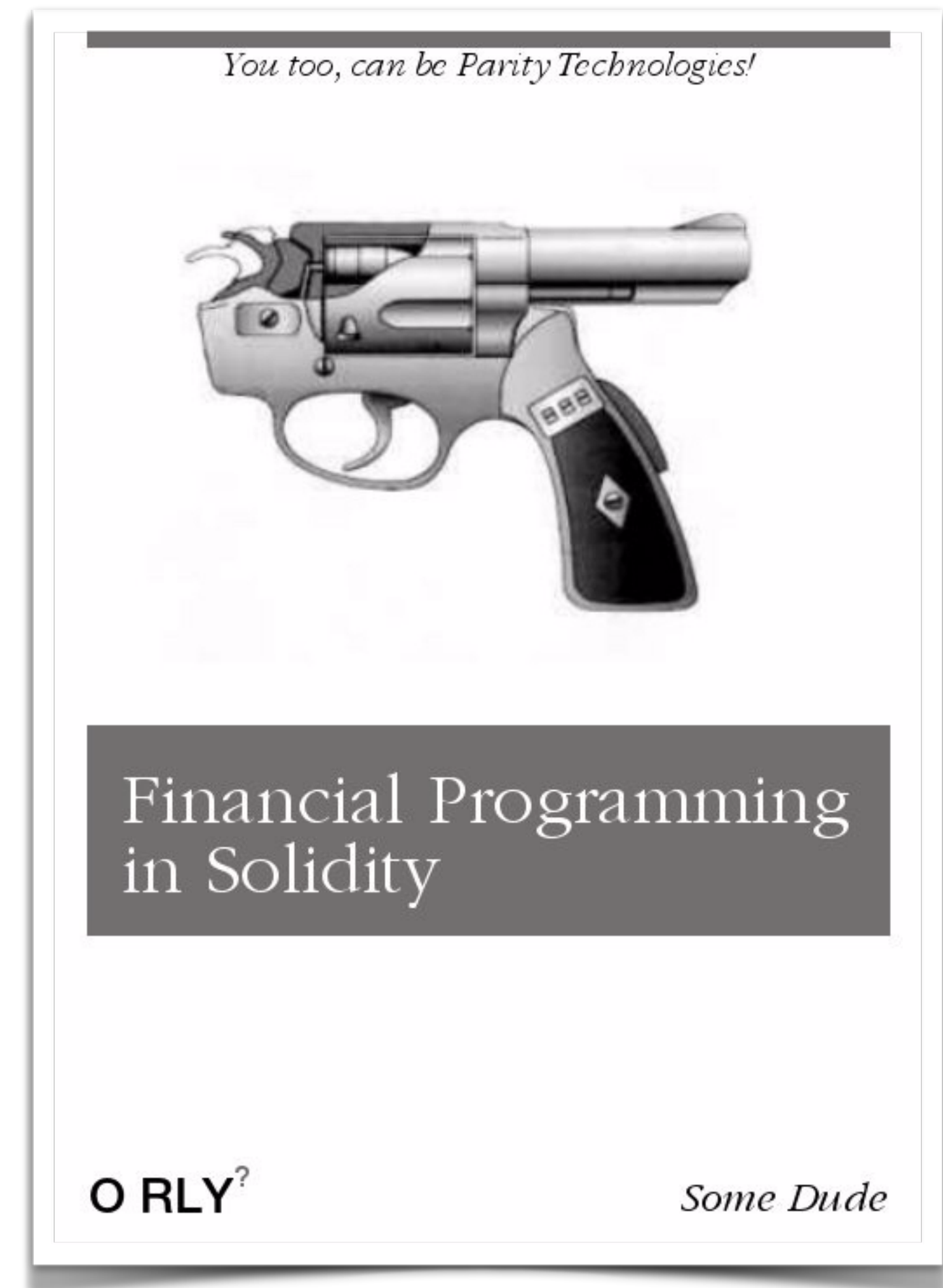
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# How to Write a Solidity Parser

(1) ☕

(2) Look up the official grammar



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- (8) Discover that existing parsers were #YOLO'd by hand

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- (9) 💉🪦



# One Does Not Simply Implement the Shunting Yard Algorithm





# One Does Not Simply Implement the Shunting Yard Algorithm





```
contract C{
    struct myStruct{
        function(uint) my_func;
    }

    function test(){
        myStruct m;

        m.my_func = call_log;
        m.my_func(0);

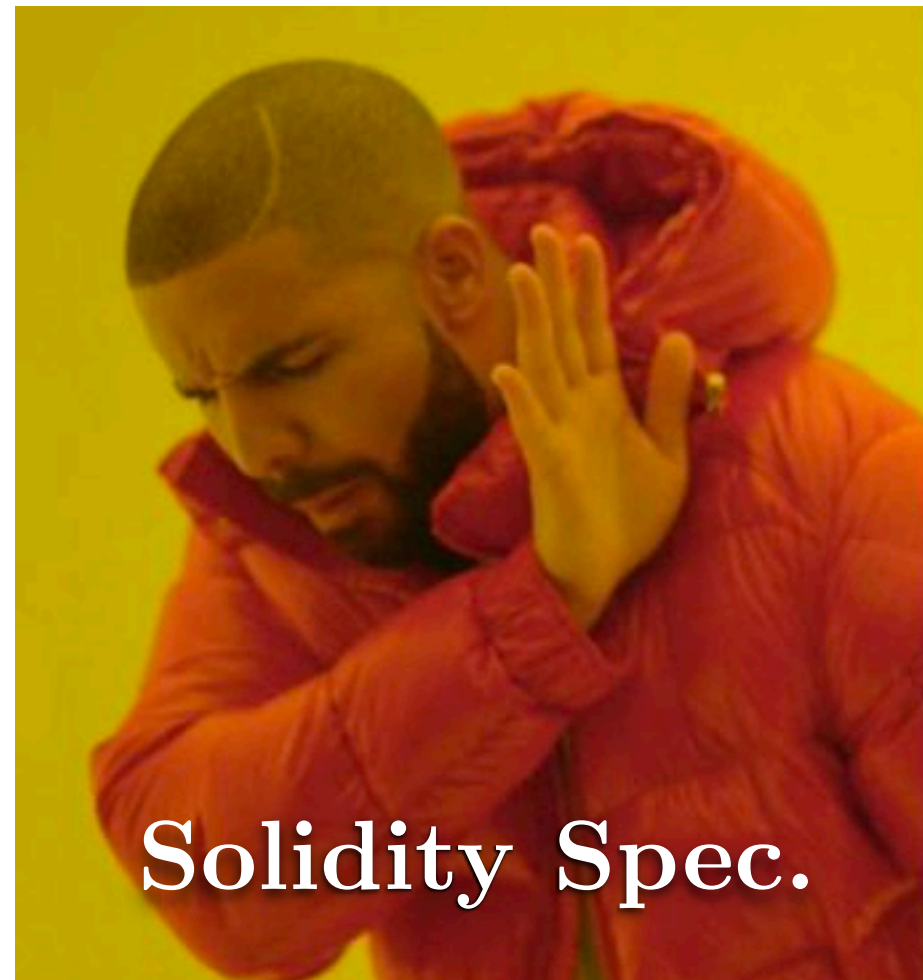
        m.my_func = call_log2;
        m.my_func(0);
    }

    function call_log(uint a){
        Log(a);
    }

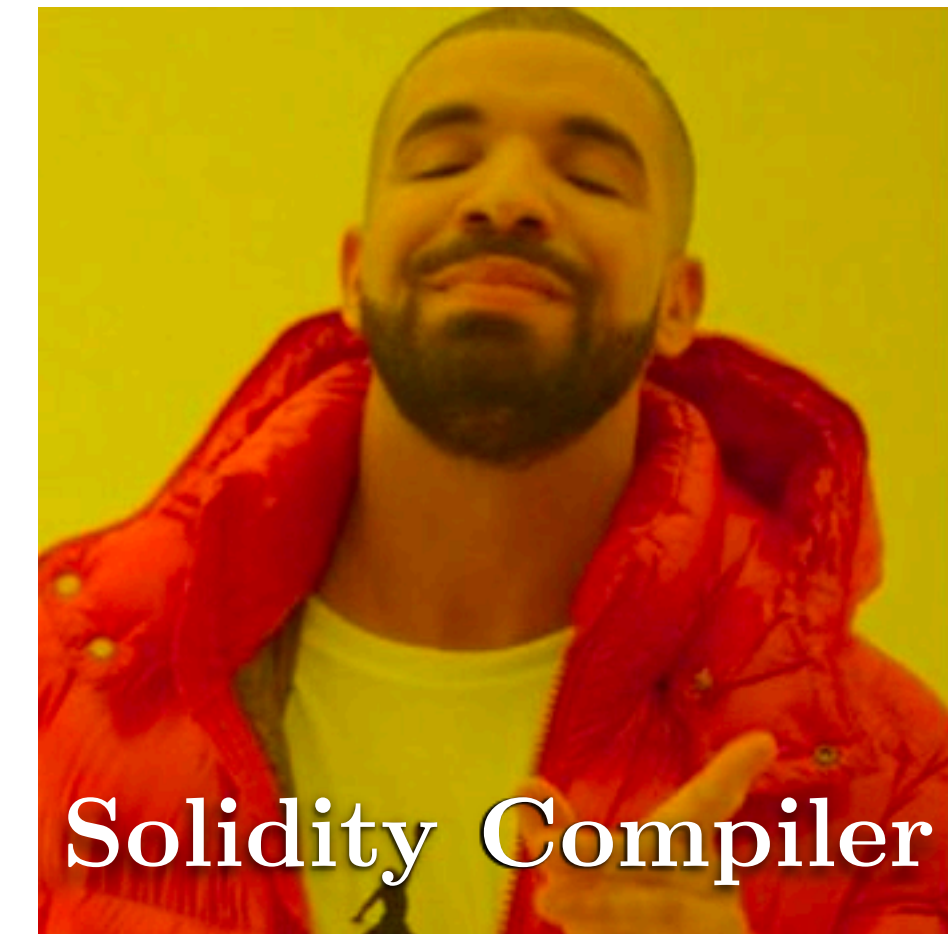
    function call_log2(uint a){
        Log2(a);
    }

    event Log(uint);
    event Log2(uint);
}
```

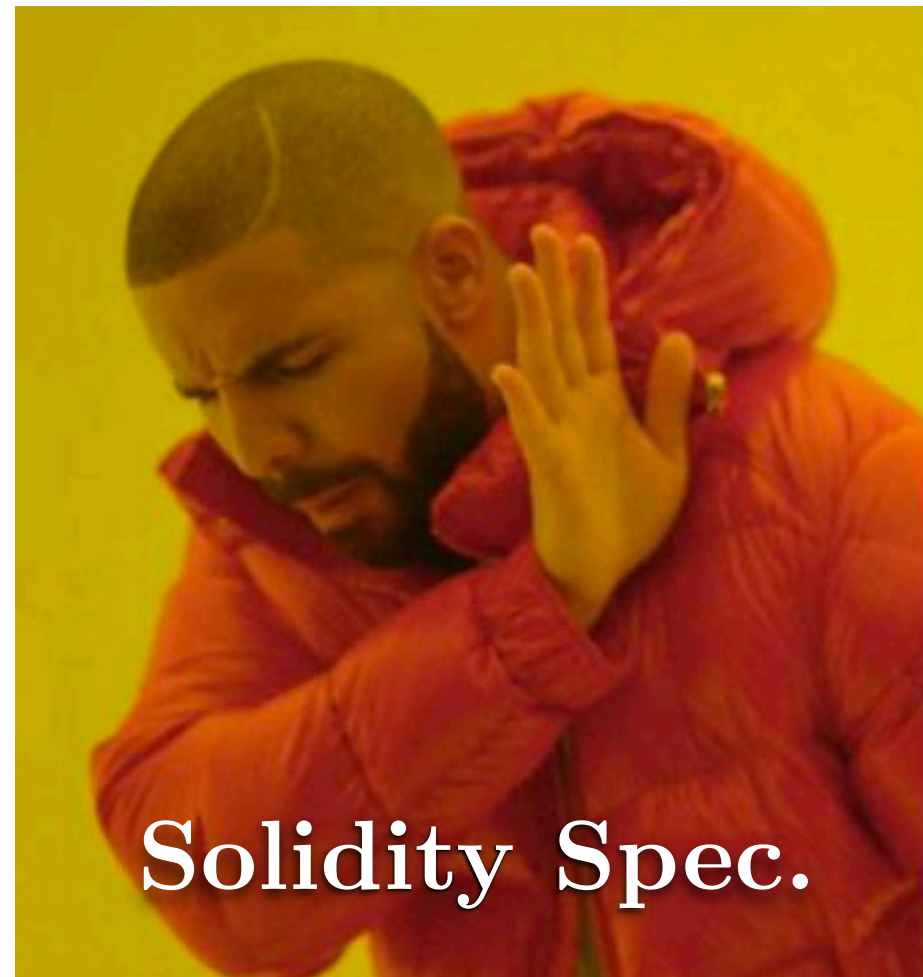
a struct that contains a pointer to a function



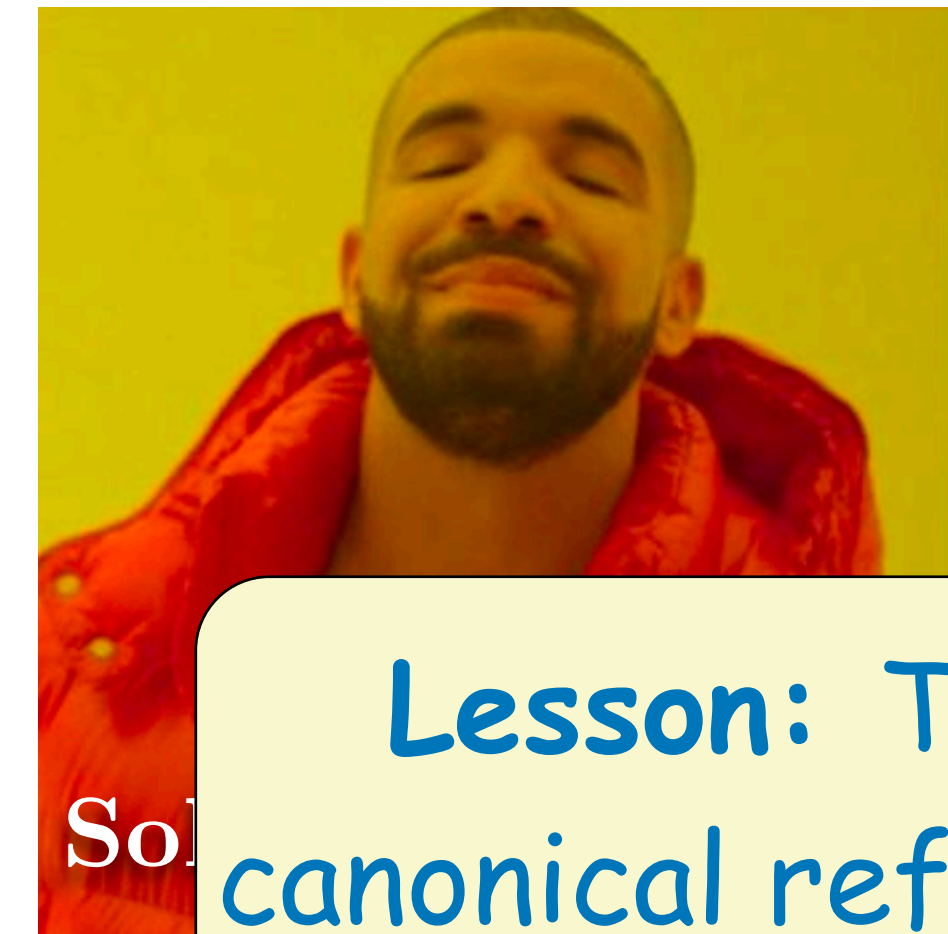
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    }  
  
    function test(){  
        myStruct m;  
  
        m.my_func = call_log;  
        m.my_func(0);  
  
        m.my_func = call_log2;  
        m.my_func(0);  
    }  
  
    function call_log(uint a){  
        Log(a);  
    }  
  
    function call_log2(uint a){  
        Log2(a);  
    }  
  
    event Log(uint);  
    event Log2(uint);  
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    }  
  
    event Log(uint);  
    event Log2(uint);  
}
```



**Lesson:** The sole,  
canonical reference for  
Solidity is the source  
code of its sole  
compiler.





# Solidity Parsing Using SmaCC: Challenges and Irregularities

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## Abstract

Solidity is a language used to implement smart contracts on a blockchain platform. Since its initial conception in 2014, Solidity has evolved into one of the major languages for the Ethereum platform as well as other blockchain technologies. Due to its popularity, there are many tools specifically designed to handle smart contracts written in Solidity. However, there is a lack of tools for Pharo to handle Solidity contracts. Therefore, we implemented a parser using SmaCC to serve as a base for further developing Solidity support in Pharo. In this paper we describe the parser creation, the irregularities we found in the Solidity grammar specification, and common practices on how to adapt the grammar to an LR type parser. Our experiences with parsing the Solidity language using SmaCC may help other developers trying to convert similar grammars.

**Keywords**   Solidity, Parser, SmaCC, Blockchain, Ethereum.

## 1. Introduction

The Blockchain technology attracted a lot attention recently [LCO<sup>+</sup>16]. Blockchain is a distributed database, managed by a peer-to-peer network that stores a list of blocks or records. Ethereum [Eth14], and BitCoin [Nak09] are examples of blockchain technologies. Blockchains can be used for many applications such as cryptocurrency, digital wallets, adhoc networks, and remote transactions [LCO<sup>+</sup>16, HL16, LMH16, Dzi15, Eth14, Nak09]. One notable application of blockchain is the execution of smart contracts [LCO<sup>+</sup>16].

Smart contracts are what embedded procedures are for databases: programs executed in the blockchain to manage and transfer digital assets. When used in platforms like Ethereum, the contract language is Turing-complete [BDLF<sup>+</sup>16]. Therefore, smart contracts can be used in many different scenarios [LCO<sup>+</sup>16]. For example, there are smart contracts employed to simple storage [Eth17], and outsourced computation [LTKS15].

Solidity [Eth17] is a programming language loosely based on JavaScript, and it is used to specify smart contracts on blockchain platforms. Solidity was originally designed to be the primary smart contract language for the Ethereum platform. Even though other contract languages have been created for Ethereum [DAKM15], Solidity is still one of the major ones. Moreover, Solidity can also be used in other blockchain platforms such as Monax<sup>1</sup> and Hyperledger<sup>2</sup>.

Probably because of its popularity, there are many tools to help integrate smart contracts written in Solidity with other languages and technologies [Eth17]. For example, we have Solidity compilers coded in C/C++ and NodeJs, third-party parsers and grammar specifications (JavaScript and ANTLR), plugins for IDEs and editors (IntelliJ, Visual Studio, Vim, Atom, and etc.). Such tool integration support developers of smart contracts. However, as far as we know, there is a lack of tools for Pharo Smalltalk to handle Solidity smart contracts. Moreover, most academic work towards smart contracts focuses on security [LCO<sup>+</sup>16, BDLF<sup>+</sup>16, DAK<sup>+</sup>15] and not in tool support.

In this paper, we plan to partially tackle this lack of tools problem by building a Solidity parser that runs on Pharo Smalltalk. We claim that with a parser and its generated AST (Abstract Syntax Tree), we will be able to develop strong tool support for Solidity contracts. For instance, it would be much easier to create code inspection tools on top of a functional parser than to rely on the purely textual content of the contract. To accomplish these goals, we used SmaCC

<sup>1</sup> <https://monax.io/>, verified 2017-06-19.

<sup>2</sup> <https://www.hyperledger.org/>, verified 2017-06-19.



## Solidity Parsing Using SmaCC: Challenges and Irregularities

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Another interesting challenge we found to parse Solidity was that the language uses the same symbol (comma) as a separator for expression lists but also as an operator for the expression itself. ... This causes a serious problem because when the parser finds a comma in the input it does not know if it is an operator for the current expression (matching the Expression rule) or a separator to the current expression and the beginning of a new one (matching ExpressionList). **This is a potential problem for any parser due to the ambiguity of matching either rule when encountering a comma.**

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[Copyright notice will appear here once 'preprint' option is removed.]

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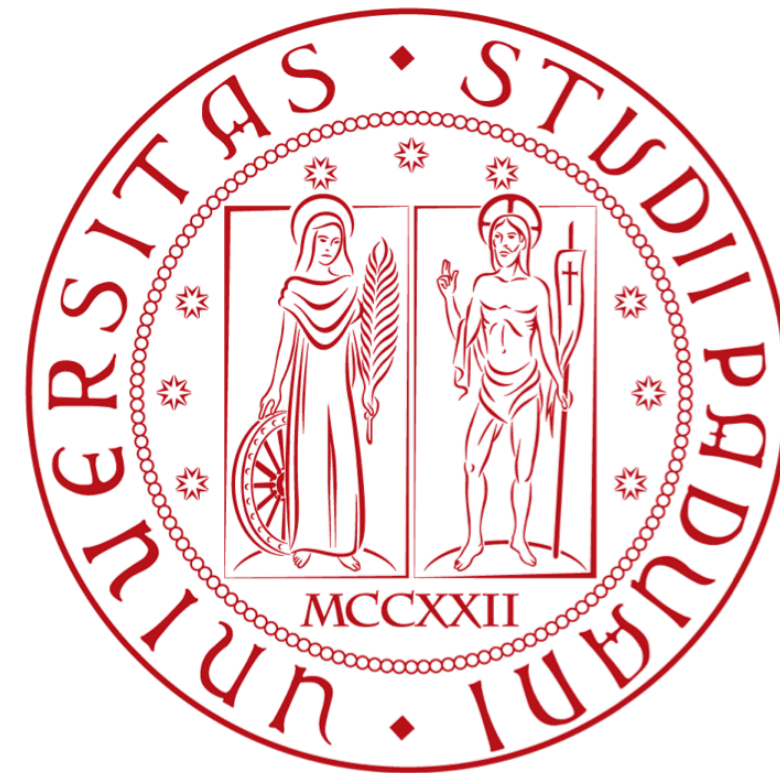
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**Università degli Studi di Padova**

DIPARTIMENTO DI MATEMATICA “TULLIO  
LEVI-CIVITA”

CORSO DI LAUREA MAGISTRALE IN INFORMATICA



## **How Solid is Solidity?**

### **An In-dept Study of Solidity’s Type Safety**

Master thesis

*Supervisor*

Prof. Silvia Crafa

*Author*

Matteo Di Pirro

---

SEPTEMBER 2018

...we found out that Solidity's type system is far from being safe with respect to any type of error:



in many occasions, contract interfaces are not consulted at compile-time, and this makes the execution raise an exception and the user waste money.

# Solidity

Implementation and Tooling



The difference between an amateur and a professional is: you write your own compiler.



**Ryan Stortz**

@withzombies

There are contracts on the blockchain that calculate 1 with exponentiation. This actually costs people money...

```
JUMPI(#0x200, %13),  
]>  
<SSA:BasicBlock ofs:0x24c insns:[  
  %14 = SLOAD(#0x3),  
  %15 = EXP(#0x100, #0x0),  
  %16 = DIV(%14, %15),  
  %17 = EXP(#0x2, #0xA0),  
  %18 = SUB(%17, #0x1),
```



# 16 Block Trace

by Martin Holst Swende

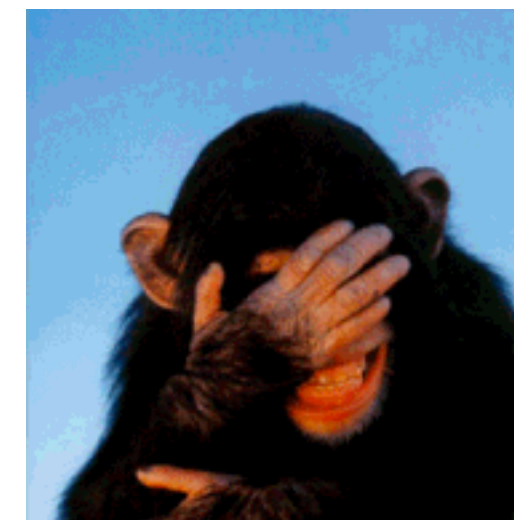


18,538 invocations of EXP

$$\frac{18538 \times 4}{20000 \times 16} = \sim 25\%$$

Well over half were calculating 160 raised to the power of 1

Martin's GitHub profile pic:



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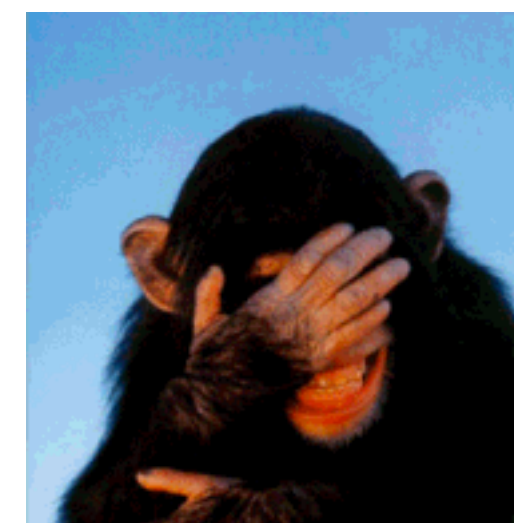


18,538 invocations of EXP

**Lesson:** Solidity is bad at optimization, but getting better, kinda (more on this later)



Well over half were calculating 160 raised to



Martin's GitHub profile pic:



# Exponentiation: How does it work?

// We need cleanup for EXP because  $0^{**}0 == 1$ , but  $0^{**}0x100 == 0$

4  libsolidity/codegen/ExpressionCompiler.cpp View 

	@@ -2069,7 +2069,9 @@ bool ExpressionCompiler::cleanupNeededForOp(Type::Category _type, Token::Value _
2069	{
2070	if (Token::isCompareOp(_op)    Token::isShiftOp(_op))
2071	return true;
2072	- else if (_type == Type::Category::Integer && (_op == Token::Div    _op == Token::Mod))
2073	return true;
2074	else
2075	return false;
	


2069	{
2070	if (Token::isCompareOp(_op)    Token::isShiftOp(_op))
2071	return true;
2072	+ else if (_type == Type::Category::Integer && (_op == Token::Div    _op == Token::Mod    _op == Token::Exp))
2073	+ // We need cleanup for EXP because $0^{**}0 == 1$ , but $0^{**}0x100 == 0$
2074	+ // It would suffice to clean the exponent, though.
2075	return true;
2076	else
2077	return false;

Using the `**` operator with an exponent of type shorter than 256 bits can result in unexpected values.



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✱

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2069 {

2070 if (Token::isCompareOp(\_op) || Token::isShiftOp(\_op))

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✱

2069 {

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2071 return true;

2072 + else if (\_type == Type::Category::Integer && (\_op == Token::Div || \_op

== Token::Mod || \_op == Token::Mod))

2073 + // We need cleanup for EXP because 0\*\*0 == 1, but 0\*\*0x100 == 0

2074 + // It would be better to have a separate cleanup function for EXP

2075 return true;

2076 else

2077 return false;

Lesson: The compiler is still immature



Using the `**` operator with an exponent of type shorter than 256 bits can result in unexpected values.

# Things Are Improving



 **Chris**  
@ethchris

Replying to @NicuOprisan

The breaking release before 0.5.0 was over two years ago. I think Solidity needs to get more flexible. We are planning breaking releases roughly every 6 months now. And I think it's fine, one reason being that you cannot change deployed code anyway.

4:12 AM - 30 Nov 2018



# Things Are ~~Improving~~ Changing



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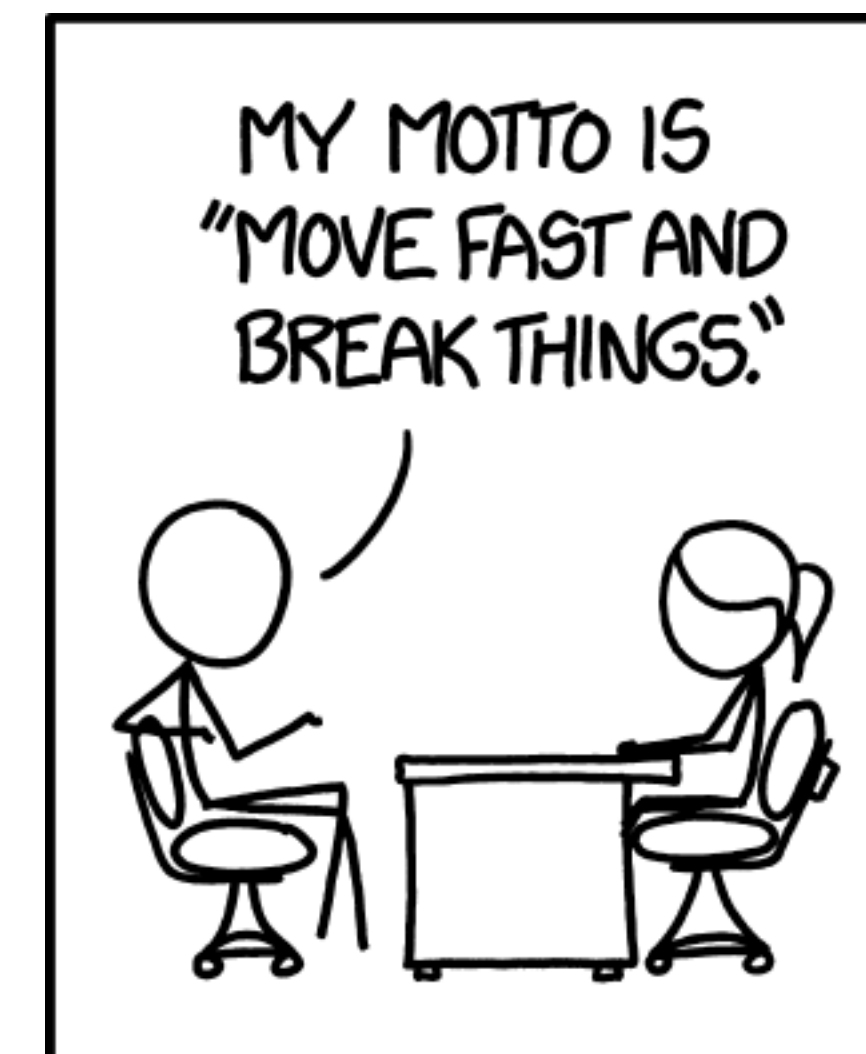


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@ethchris

Replying to @NicuOprisan

The breaking release before 0.5.0 was over two years ago. I think Solidity needs to get more flexible. We are planning breaking releases roughly every 6 months now. And I think it's fine, one reason being that you cannot change deployed code anyway.

4:12 AM - 30 Nov 2018



**JOBS I'VE BEEN  
FIRED FROM  
SOLIDITY DEV**

Adapted from <https://xkcd.com/1428/>



# Things Are ~~Improving~~

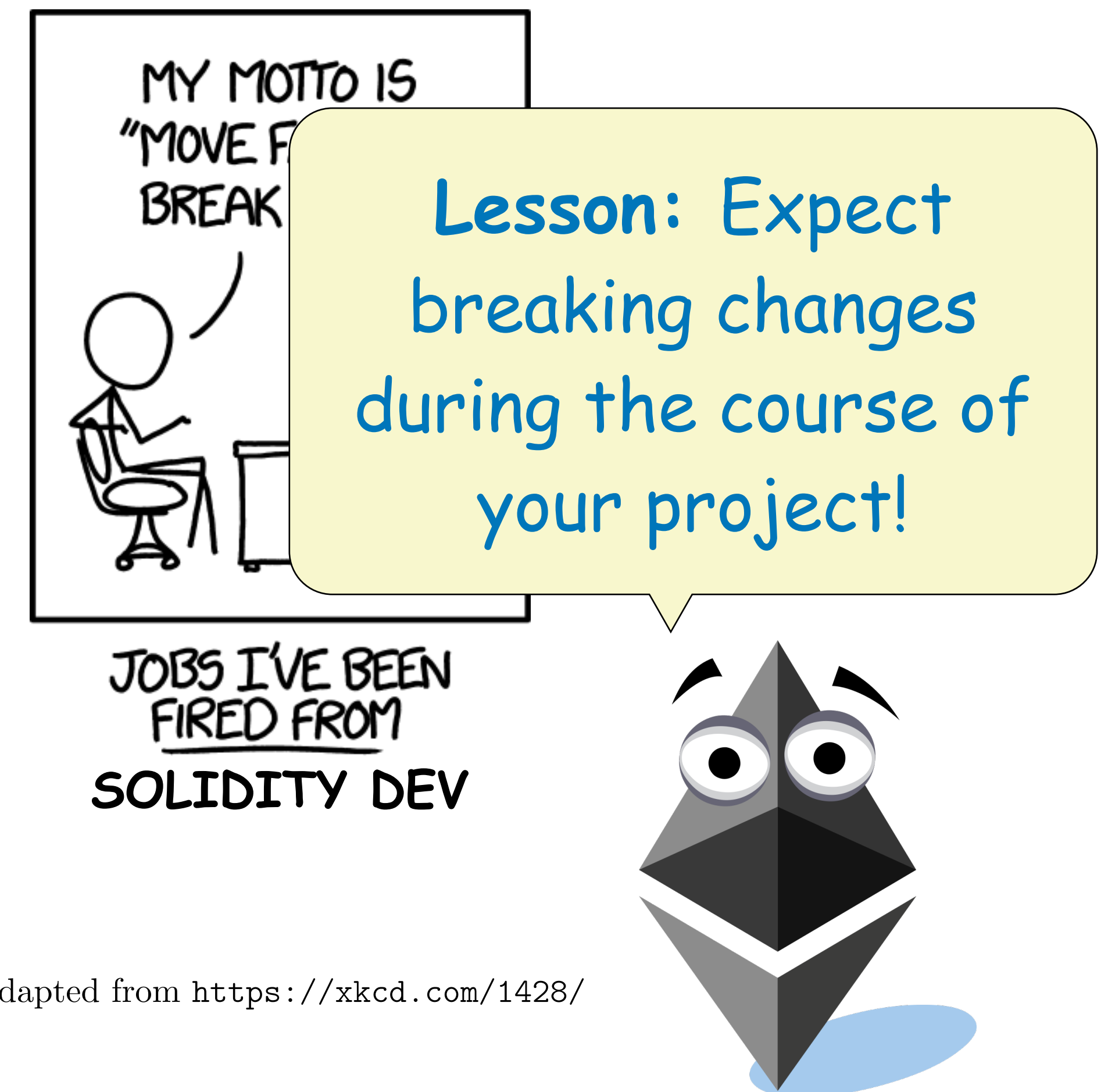


**Chris**  
@ethchris

Replying to @NicuOprisan

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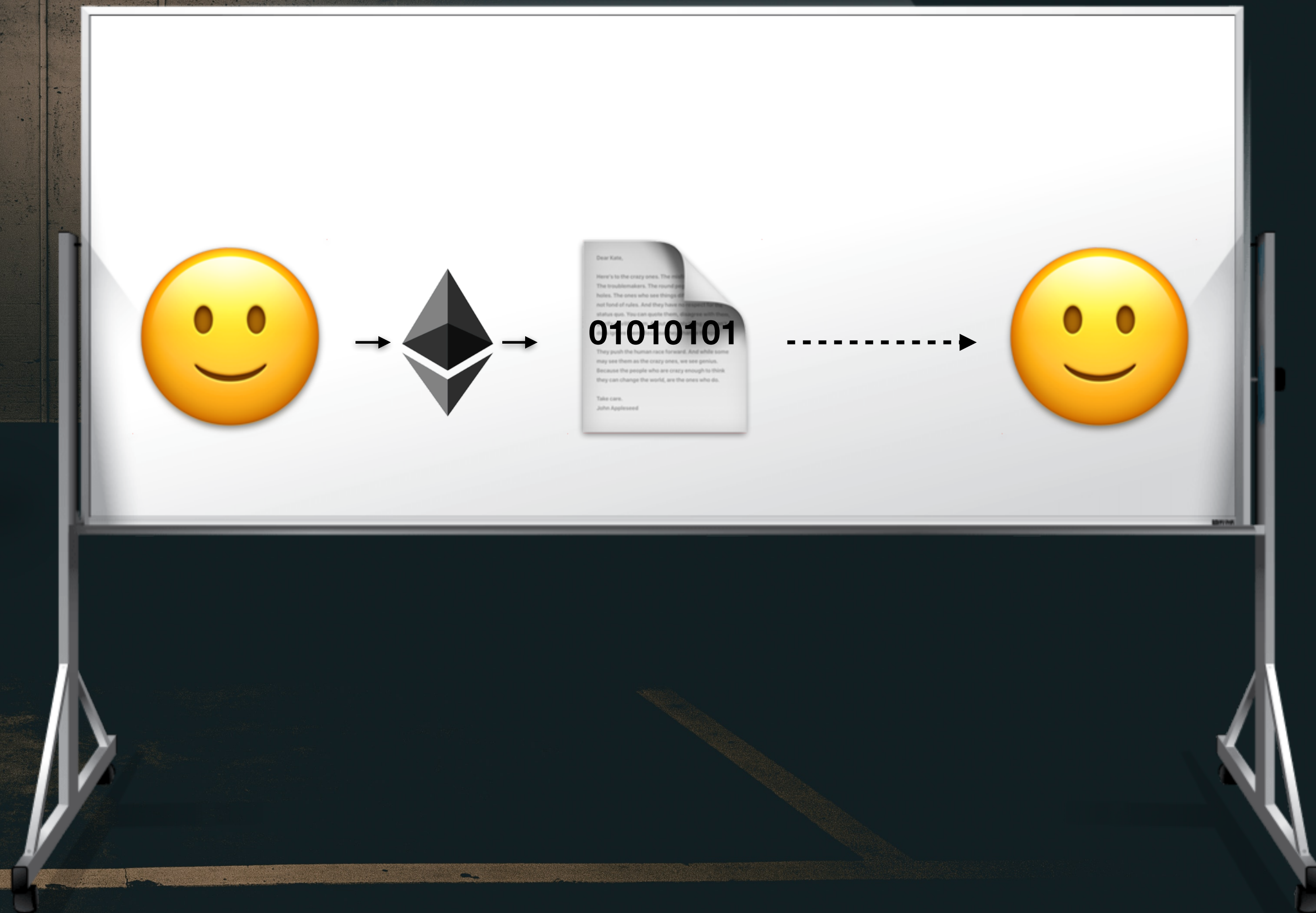


# Upgradable Contracts



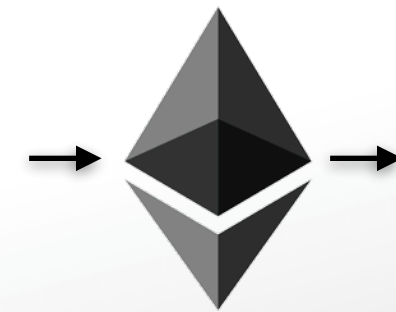


# Upgradable Contracts





# Upgradable Contracts





# Upgradable Contracts



**Lesson:** If you absolutely have to use the `DELEGATECALL` proxy upgrade pattern, then you must always make sure the storage layout of your new contract matches the old one!





# Backward Compatibility?





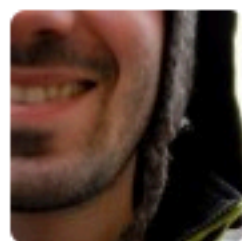
# Backward Compatibility?





# Backward Compatibility?





mo-seph commented 27 days ago • edited ▼



Compilation fails for `solidity` using recommended install method. I'm on macos 10.13.6, and I've just installed brew to compile solc.

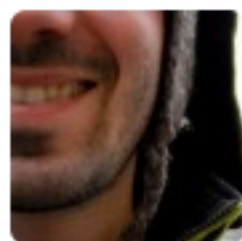
I've run

```
brew update  
brew upgrade  
brew tap ethereum/ethereum
```

I've tried installing the latest version, also tried 0.4.24. Here's a log with the latest version

- 
- 
-





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•  
•  
•



axic commented 26 days ago

Member



Since it works with 0.5.0, which has been released now, closing this issue.



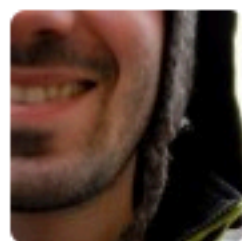
1



1



axic closed this 26 days ago



mo-seph commented 27 days ago • edited ▼



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I've run

```
brew update
brew upgrade
brew tap ethereum/ethereum
```

I've tried installing the latest version, also tried 0.4.24. Here's a log with the latest

⋮

**Lesson: Use  
solc-select!**



axic commented 26 days ago

Member



Since it works with 0.5.0, which has been released now, closing this issue.



1



1



axic closed this 26 days ago



# Optimizations are Dangerous

- Compiler optimization still in active development
- Independent compiler audit in November of 2018 concluded optimizations are dangerous
- Numerous high severity bugs related to the optimizer, many excluded from the changlog
- There are likely latent bugs



# Optimizations are Dangerous

- Compiler optimization still in active development
- Independent compiler audit in November of 2018 confirmed optimizations are dangerous
- Numerous high severity bugs related to the optimizations excluded from the changelog
- There are likely latent bugs

**Lesson:** Don't turn on  
solc optimizations  
unless you really, really  
know what you are  
doing



# On the Horizon

# They're Proposing a New Intermediate Representation, YUL

```
Block = '{' Statement* '}'
Statement =
    Block |
    FunctionDefinition |
    VariableDeclaration |
    Assignment |
    Expression |
    Switch |
    ForLoop |
    BreakContinue
FunctionDefinition =
    'function' Identifier '(' TypedIdentifierList? ')'
    ( '-'>' TypedIdentifierList )? Block
VariableDeclaration =
    'let' TypedIdentifierList ( ':' Expression )?
Assignment =
    IdentifierList ':' Expression
Expression =
    FunctionCall | Identifier | Literal
If =
    'if' Expression Block
Switch =
    'switch' Expression Case* ( 'default' Block )?
Case =
    'case' Literal Block
```

```
ForLoop =
    'for' Block Expression Block Block
BreakContinue =
    'break' | 'continue'
FunctionCall =
    Identifier '(' ( Expression ( ',' Expression )* )? ')'
Identifier = [a-zA-Z_$] [a-zA-Z_0-9]*
IdentifierList = Identifier ( ',' Identifier)*
TypeName = Identifier | BuiltinTypeName
BuiltinTypeName = 'bool' | [us] ( '8' | '32' | '64' | '128' | '256' )
TypedIdentifierList = Identifier ':' TypeName ( ',' Identifier ':' TypeName )*
Literal =
    (NumberLiteral | StringLiteral | HexLiteral | TrueLiteral | FalseLiteral) ':'
TypeName
NumberLiteral = HexNumber | DecimalNumber
HexLiteral = 'hex' ( '"' ([0-9a-fA-F]{2})* '"' | '\' ' ' ([0-9a-fA-F]{2})* '\' ' ' )
StringLiteral = '"' ([^"\r\n\\] | '\\ ' .)* '"'
TrueLiteral = 'true'
FalseLiteral = 'false'
HexNumber = '0x' [0-9a-fA-F]+
DecimalNumber = [0-9]+
```



# They're Proposing a New Intermediate Representation, YUL

```
Block = '{' Statement* '}'
Statement =
    Block |
    FunctionDefinition |
    VariableDeclaration |
    Assignment |
    Expression |
    Switch |
    ForLoop |
    BreakContinue
FunctionDefinition =
    'function' Identifier '(' TypedIdentifierList? ')'
    ( '->' TypedIdentifierList )? Block
VariableDeclaration =
    'let' TypedIdentifierList ( ':' Expression )?
Assignment =
    IdentifierList ':' Expression
Expression =
    FunctionCall | Identifier | Literal
If =
    'if' Expression Block
Switch =
    'switch' Expression Case* ( 'default' Block )?
Case =
    'case' Literal Block
```

```
ForLoop =
    'for' Block Expression Block Block
BreakContinue =
    'break' | 'continue'
FunctionCall =
    Identifier '(' ( Expression ( ',' Expression )* )? ')'
Identifier = [a-zA-Z_$] [a-zA-Z_0-9]*
IdentifierList = Identifier ( ',' Identifier)*
TypeName = Identifier | BuiltinTypeName
BuiltinTypeName = 'bool' | [us] ( '8' | '32' | '64' | '128' | '256' )
TypedIdentifierList = Identifier ':' TypeName ( ',' Identifier ':' TypeName )*
Literal =
    (NumberLiteral | StringLiteral | HexLiteral | TrueLiteral | FalseLiteral) ':'
TypeName
NumberLiteral = HexNumber | DecimalNumber
HexLiteral = 'hex' ( '"' ([0-9a-fA-F]{2})* '"' | '\\' ([0-9a-fA-F]{2})* '\\' )
StringLiteral = '"' ([^"r\n\\] | '\\' .)* '"'
TrueLiteral = 'true'
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```

All of this has happened before ... and will happen again.

# They're Proposing a New Intermediate Representation, YUL

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Block = '{' Statement* '}'  
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    Block |  
    FunctionDefinition |  
    VariableDeclaration |  
    Assignment |  
    Expression |  
    Switch |  
    ForLoop |  
    BreakContinue
```

**The “If” production rule is never used!**

```
VariableDeclaration =  
    'let' TypedIdentifierList ( ':' Expression )?  
Assignment =  
    IdentifierList ':' Expression  
Expression =  
    FunctionCall | Identifier | Literal  
If =  
    'if' Expression Block  
Switch =  
    'switch' Expression Case* ( 'default' Block )?  
Case =  
    'case' Literal Block
```

```
ForLoop =  
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FunctionCall =  
    Identifier '(' ( Expression ( ',' Expression )* )? ')'  
Identifier = [a-zA-Z_$] [a-zA-Z_0-9]*  
IdentifierList = Identifier ( ',' Identifier)*  
Identifier | BuiltinTypeName  
'bool' | [us] ( '8' | '32' | '64' | '128' | '256' )  
t = Identifier ':' TypeName ( ',' Identifier ':' TypeName )*  
Literal =  
    (NumberLiteral | StringLiteral | HexLiteral | TrueLiteral | FalseLiteral) ':'  
TypeName  
NumberLiteral = HexNumber | DecimalNumber  
HexLiteral = 'hex' ( '"' ([0-9a-fA-F]{2})* '"' | '\' ([0-9a-fA-F]{2})* '\' )  
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TrueLiteral = 'true'  
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Block = '{' Statement* '}'  
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    FunctionDefinition |  
    VariableDeclaration |  
    Assignment |  
    Expression |  
    Switch |  
    ForLoop |  
    BreakContinue
```

**The “If” production rule is never used!**

```
VariableDeclaration =  
    'let' TypedIdentifierList ( ':' Expression )?  
Assignment =  
    IdentifierList ':' Expression  
Expression =  
    FunctionCall | Identifier | Literal  
If =  
    'if' Expression Block  
Switch =  
    'switch' Expression Case* ( 'default' Block )?  
Case =  
    'case' Literal Block
```

**The default switch case isn't followed by a ':'**

```
ForLoop =  
    'for' Block Expression Block Block  
BreakContinue =  
    'break' | 'continue'  
FunctionCall =  
    Identifier '(' ( Expression ( ',' Expression )* )? ')'  
Identifier = [a-zA-Z_$] [a-zA-Z_0-9]*  
IdentifierList = Identifier ( ',' Identifier )*  
TypedIdentifier = Identifier | BuiltinTypeName  
BuiltinTypeName = 'bool' | [us] ( '8' | '32' | '64' | '128' | '256' )  
Type = Identifier ':' TypeName ( ',' Identifier ':' TypeName )*  
Literal =  
    (NumberLiteral | StringLiteral | HexLiteral | TrueLiteral | FalseLiteral) ':'  
HexLiteral = hex ( '0x' | '0X' ) ( [0-9a-fA-F] | 'x' )*  
StringLiteral = '"' ( [^"\\r\\n\\] | '\\.' )* '"'  
TrueLiteral = 'true'  
FalseLiteral = 'false'  
HexNumber = '0x' [0-9a-fA-F]+  
DecimalNumber = [0-9]+
```

All of this has happened before ... and will happen again.

# They're Proposing a New Intermediate Representation, YUL

“switch foo” is a legal production in this grammar

```
FunctionDefinition |
VariableDeclaration |
Assignment |
Expression |
Switch |
ForLoop |
BreakContinue
```

```
'for' Block Expression Block Block
```

```
BreakContinue =
```

```
'break' | 'continue'
```

FunctionCall =

Identifier '(' ( Expression ( ',' Expression ) \* ) ? ')'

Identifier = [a-zA-Z\_\$] [a-zA-Z\_0-9]\*

```
IdentifierList = Identifier ( ',' Identifier)*
```

```
ier | BuiltinTypeName
```

```
'bool' | [us] ( '8' | '32' | '64' | '128' | '256' )
```

```
t = Identifier ':' TypeName ( ',' Identifier ':' TypeName )*
```

Literal =

```
(NumberLiteral | StringLiteral | HexLiteral | TrueLiteral | FalseLiteral) ':'
```

```
VariableDeclaration =
```

```
'let' TypedIdentifierList ( ':= ' Expression )?
```

Assignment =

IdentifierList  $\mathrel{:=}$  Expression

Expression =

FunctionCall	Identifier	Literal
FunctionCall	Identifier	Literal

If =

### 'if' Expression Block

## Switch

```
'switch' Expression Case* ( 'default' Block )?
```

Case =

```
'case' Literal Block
```

## The default switch case isn't followed by a ':'

```
hexLiteral = hex ('' ([0-9a-fA-F]{2})* | \ ([0-9a-fA-F]{2})* \ ')
```

```
StringLiteral = ''' ([^"\\r\\n\\| | '\\\\' .])* '''
```

```
TrueLiteral = 'true'
```

```
FalseLiteral = 'false'
```

```
HexNumber = '0x' [0-9a-fA-F]+
```

DecimalNumber = [0-9]+

All of this has happened before ... and will happen again.



# They're Proposing a New Intermediate Representation, YUL

“switch foo” is a legal production in this grammar

StringLiteral can't be represented without casting

The “If” production rule is never used!

The default switch case isn't followed by a ‘:’

Block = '{' Statement\* '}'

FunctionDefinition |  
VariableDeclaration |  
Assignment |  
Expression |  
Switch |  
ForLoop |  
BreakContinue

'for' Block Expression Block Block  
BreakContinue =

Identifier = [a-zA-Z\_\$] [a-zA-Z\_0-9]\*  
IdentifierList = Identifier ( ',' Identifier)\*

Identifier | BuiltinTypeName  
'bool' | [us] ( '8' | '32' | '64' | '128' | '256' )  
t = Identifier ':' TypeName ( ',' Identifier ':' TypeName )\*

VariableDeclaration =  
'let' TypedIdentifierList ( ':' Expression )?

Assignment =  
IdentifierList ':' Expression

Expression =  
FunctionCall | Identifier | Literal

If =  
'if' Expression Block

Switch =  
'switch' Expression Case\* ( 'default' Block )?

Case =  
'case' Literal Block

Literal =  
(NumberLiteral | StringLiteral | HexLiteral | TrueLiteral | FalseLiteral) ':'

HexLiteral = hex ( '0x' | '0X' ) ( [0-9a-fA-F] )\*

StringLiteral = '"' ( [^"\\r\\n\\] | '\\\\' | '\\\'' )\* '"'

TrueLiteral = 'true'

FalseLiteral = 'false'

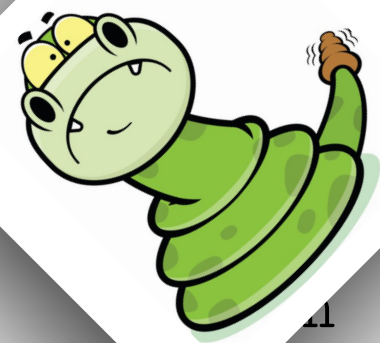
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All of this has happened before ... and will happen again.

# They're Proposing a New Intermediate Representation, Y

EVM::SSA



SlithIR

is a legal production in this grammar

StringLiteral can't be represented without casting

The "If" production rule is never used!

The default switch case isn't followed by a ':'

```
Block =  
"switch" Expression Block Block  
"for" Block Expression Block Block  
BreakContinue =  
Identifier = [a-zA-Z_$] [a-zA-Z_0-9]*  
IdentifierList = Identifier ( ',' Identifier ) *  
Identifier | BuiltinTypeName  
'bool' | [us] ( '8' | '32' | '64' | '128' | '256' )  
t = Identifier ':' TypeName ( ',' Identifier ':' TypeName ) *  
Literal =  
(NumberLiteral | StringLiteral | HexLiteral | TrueLiteral | FalseLiteral) ':'  
HexLiteral = hex ( [0-9a-fA-F] {2} ) *  
StringLiteral = '"' ( [^"\\r\\n\\] | '\\ ' . ) * '"'  
TrueLiteral = 'true'  
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HexNumber = '0x' [0-9a-fA-F] +  
DecimalNumber = [0-9] +  
VariableDeclaration =  
  'let' TypedIdentifierList ( ':' Expression ) ?  
Assignment =  
  IdentifierList ':' Expression  
Expression =  
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If =  
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Switch =  
  'switch' Expression Case* ( 'default' Block ) ?  
Case =  
  'case' Literal Block
```

All of this has happened before ... and will happen again.



# Solidity Alternatives

- Even more immature
- Lack of security tooling
- Different semantics!



Bugs!



# Compiler Warnings

```
1 pragma solidity ^0.4.9;
2
3 contract SafeMath {
4     /**
5      * @dev Adds two numbers, throws on overflow.
6      */
7     function add(uint256 _a, uint256 _b) internal pure returns (uint256) {
8         uint256 c = _a + _b;
9         if(c < _a) {
10             throw;
11         }
12
13         return c;
14     }
15 }
```

# Compiler Warnings

```
1 pragma solidity ^0.4.9;
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3 contract SafeMath {
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13         return c;
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15 }
```

**Warning:** "throw" is deprecated in favour of "revert()", "require()" and "assert()"





# Compiler Warnings

```
1 pragma solidity ^0.4.9;
2
3 contract SafeMath {
4     /**
5      * @dev Adds two numbers, throws on overflow.
6      */
7     function add(uint256 _a, uint256 _b) internal pure returns (uint256) {
8         uint256 c = _a + _b;
9         if(c < _a) {
10             throw; assert;
11         }
12
13         return c;
14     }
15 }
```



# Compiler Warnings

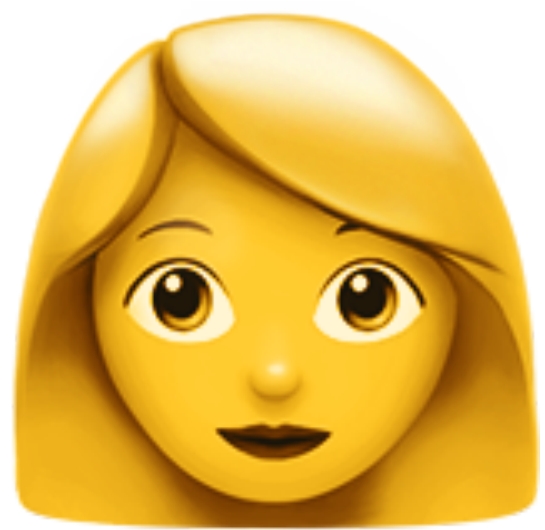


# Reentrancy

```
1 mapping (address => uint) private userBalances;  
2  
3 function withdrawBalance() public {  
4     uint amountToWithdraw = userBalances[msg.sender];  
5     require(msg.sender.call.value(amountToWithdraw)());  
6     userBalances[msg.sender] = 0;  
7 }
```

# Reentrancy

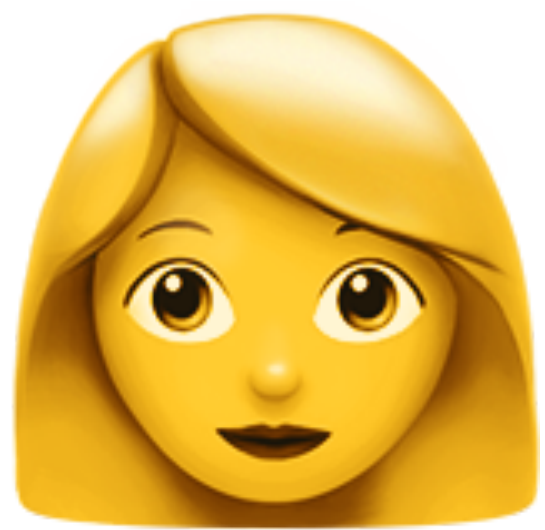
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# Reentrancy

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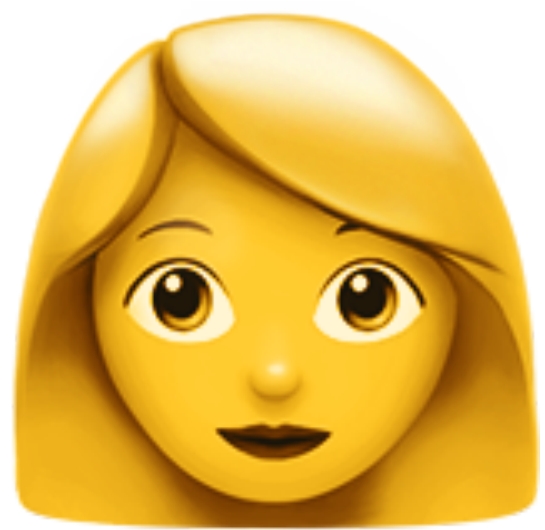


Deploy Attack Contract

```
1 uint8 toWithdraw = 2;  
2 address target = ...;  
3  
4 function () public payable {  
5     if (--toWithdraw > 0) {  
6         pwn();  
7     }  
8 }  
9  
10 function pwn() public {  
11     target.withdrawBalance()  
12 }
```

# Reentrancy

```
1 mapping (address => uint) private userBalances;  
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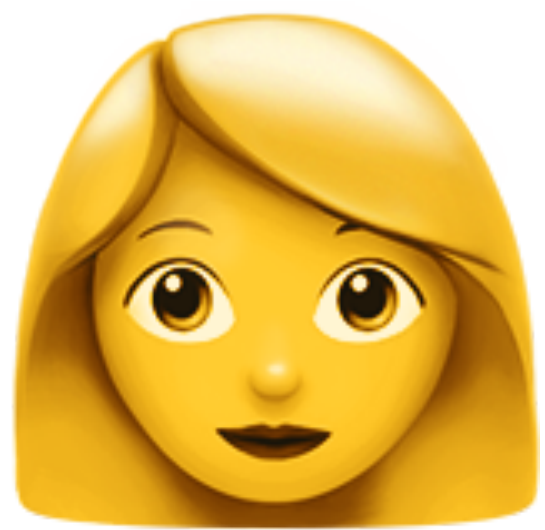
*pwn()*

```
1 uint8 toWithdraw = 2;  
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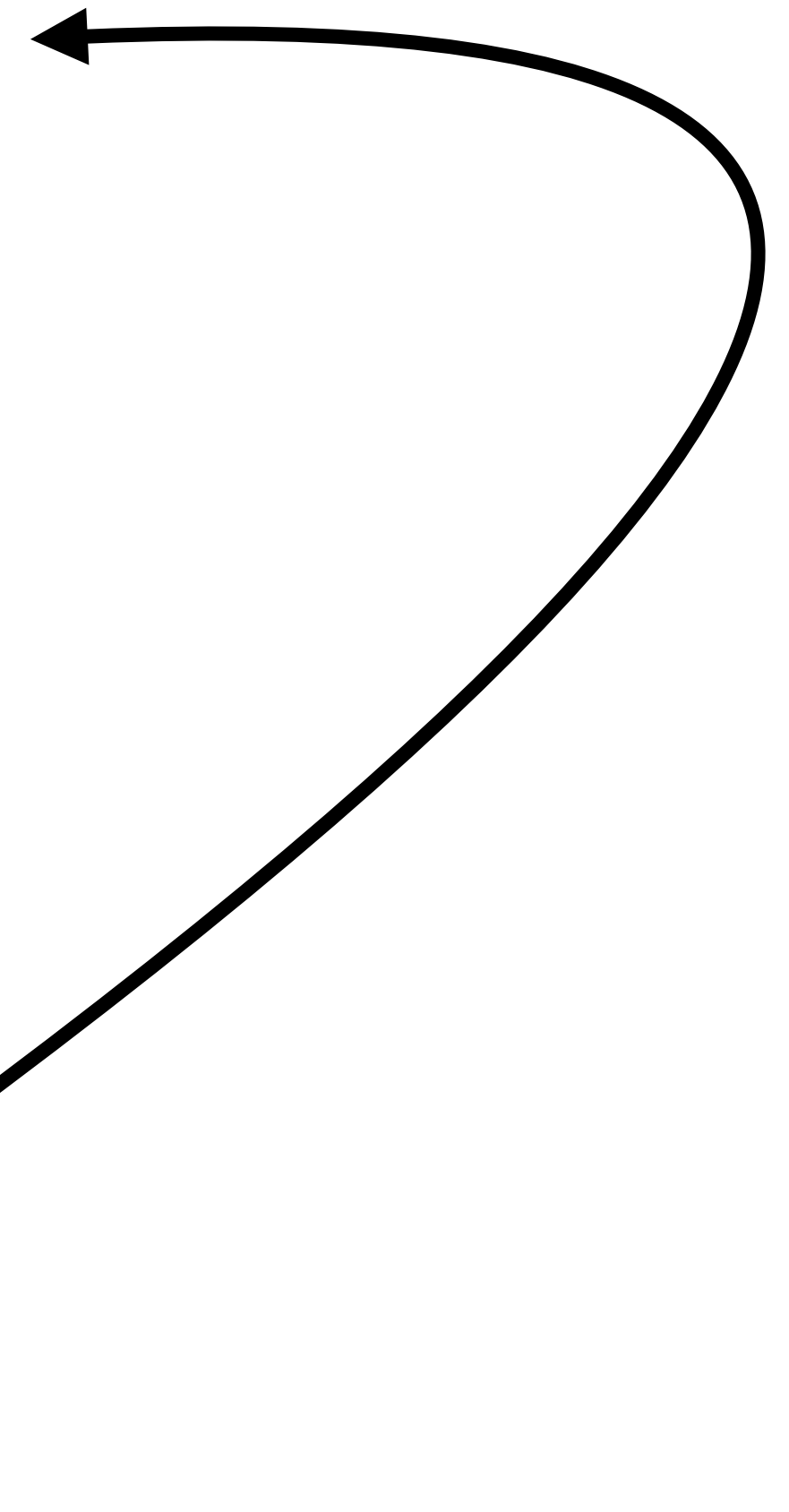
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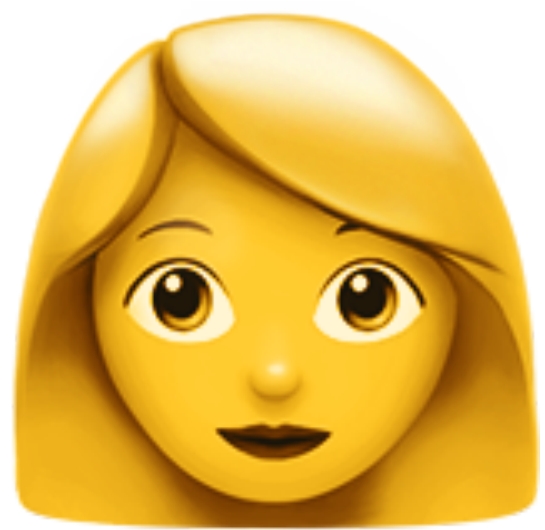
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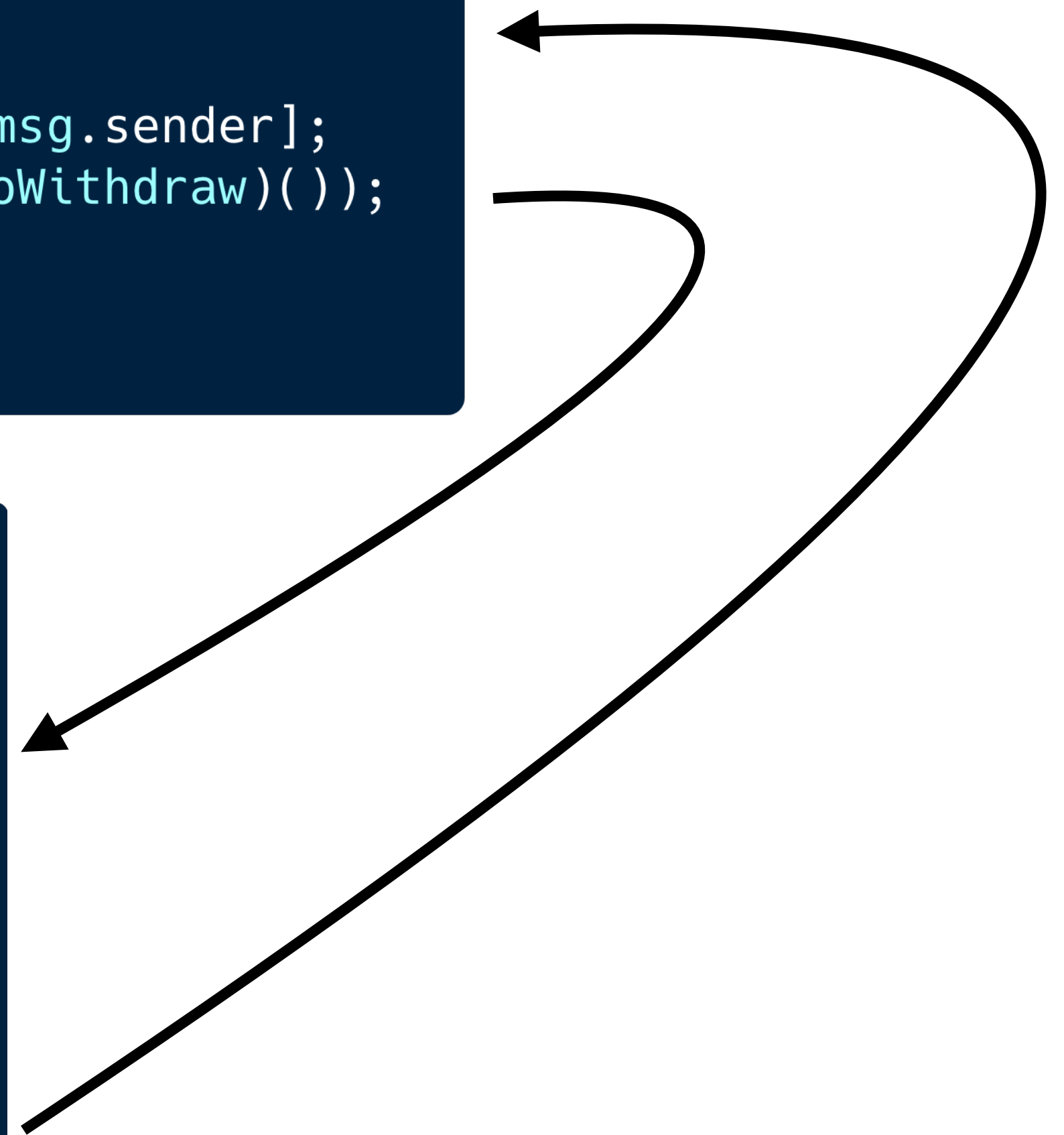
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*pwn()*

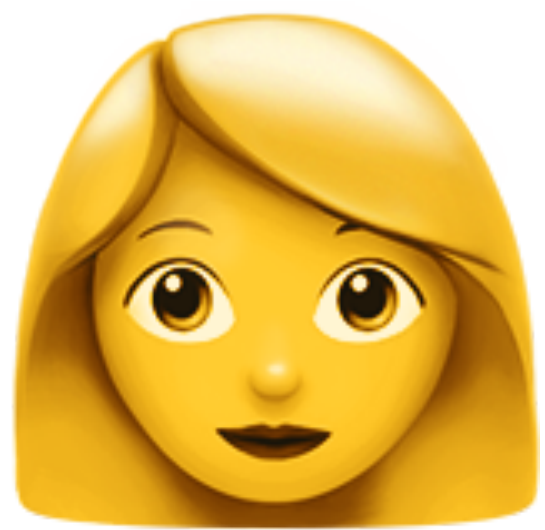
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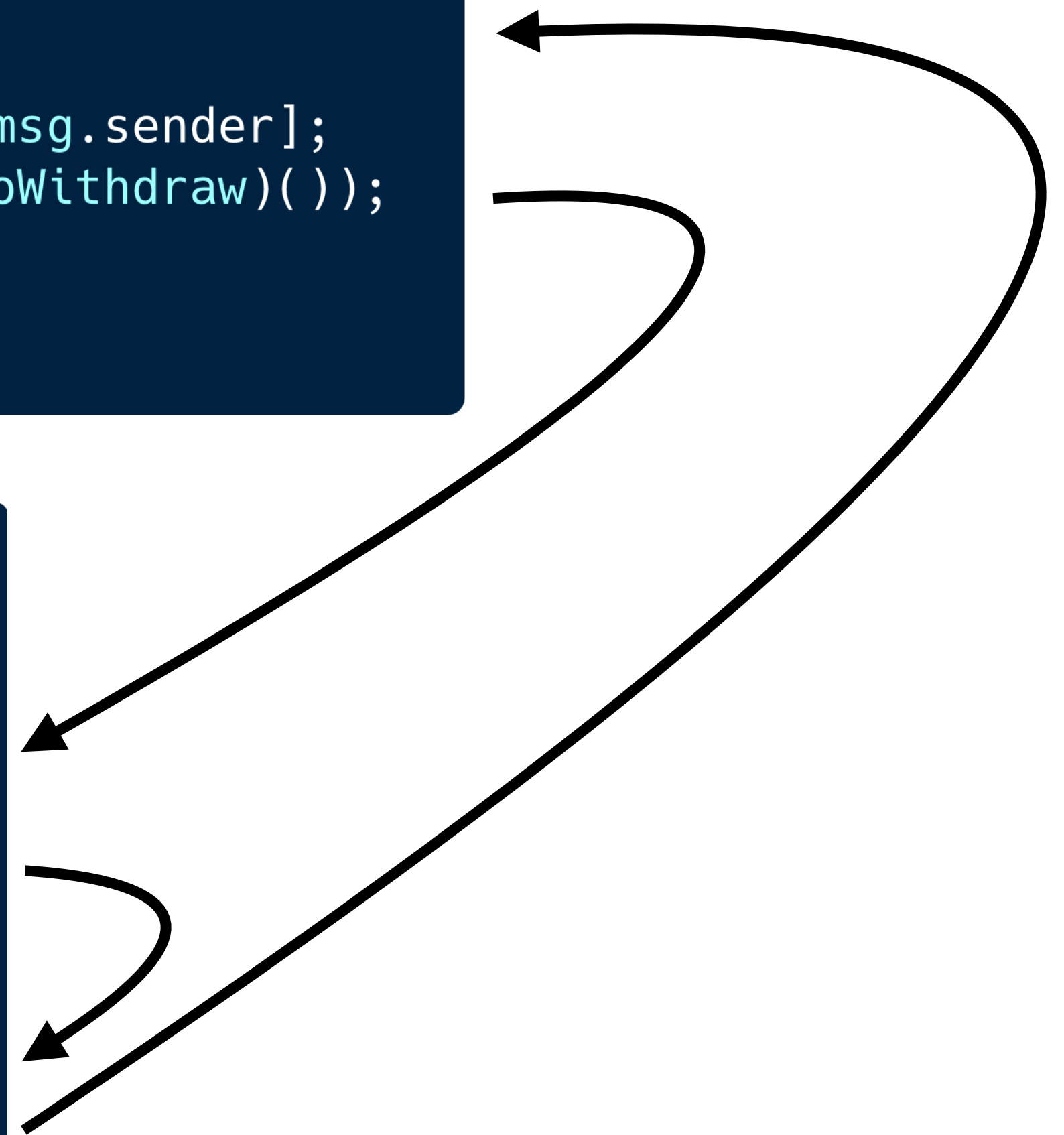
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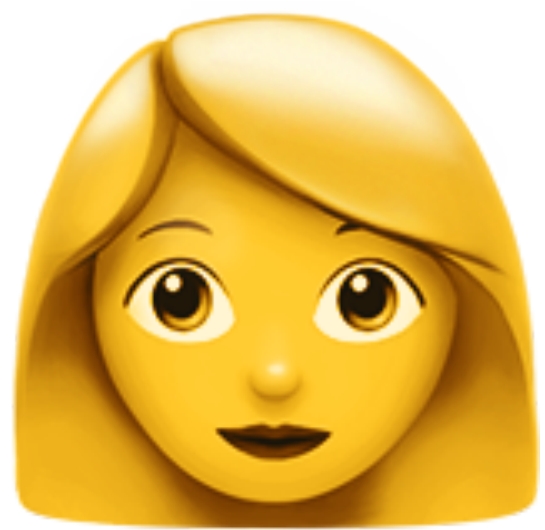
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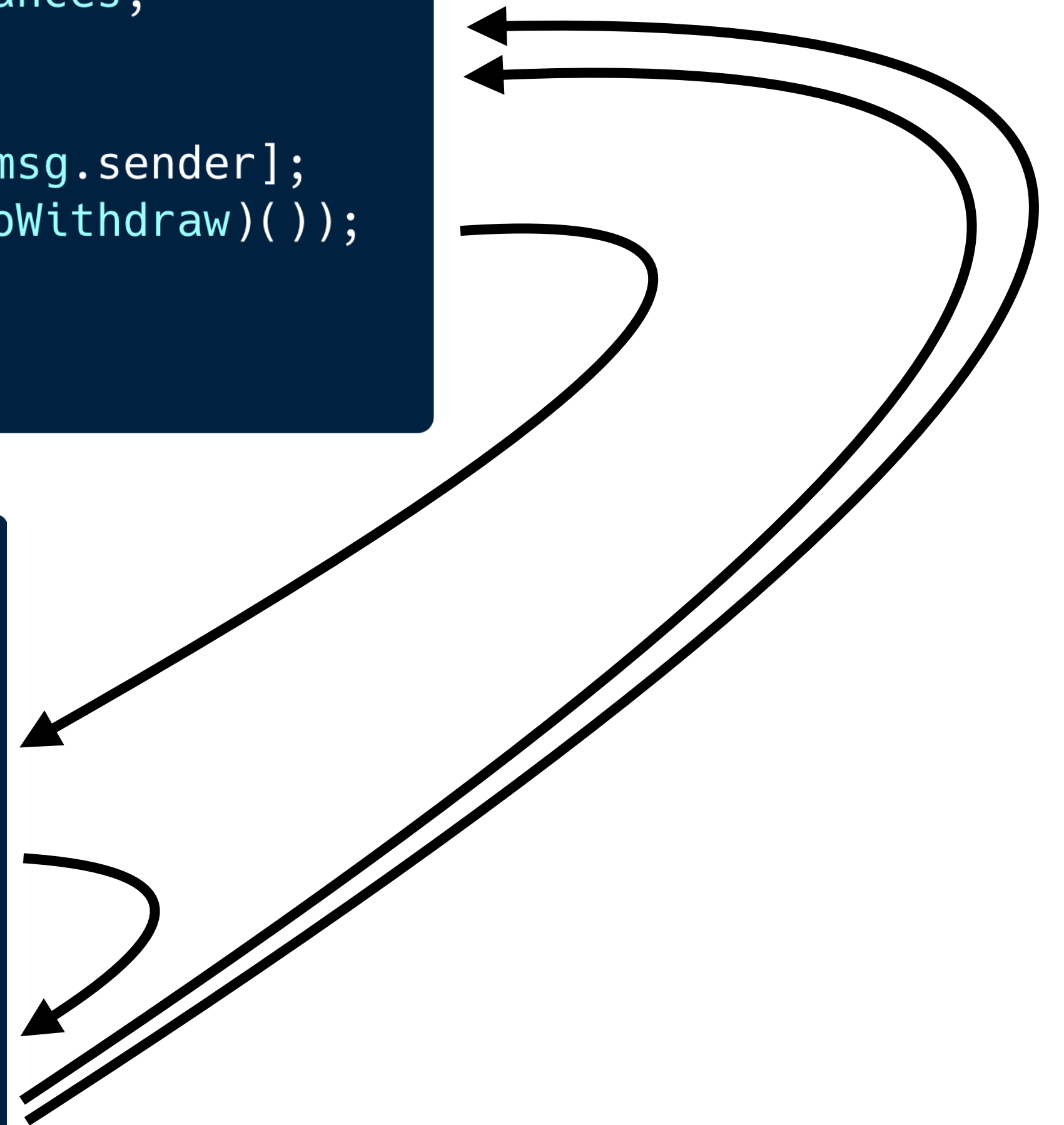
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*pwn()*

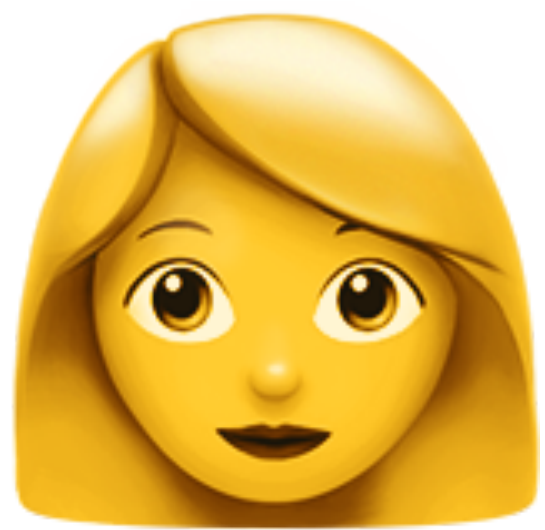
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8 }  
9  
10 function pwn() public {  
11     target.withdrawBalance()  
12 }
```





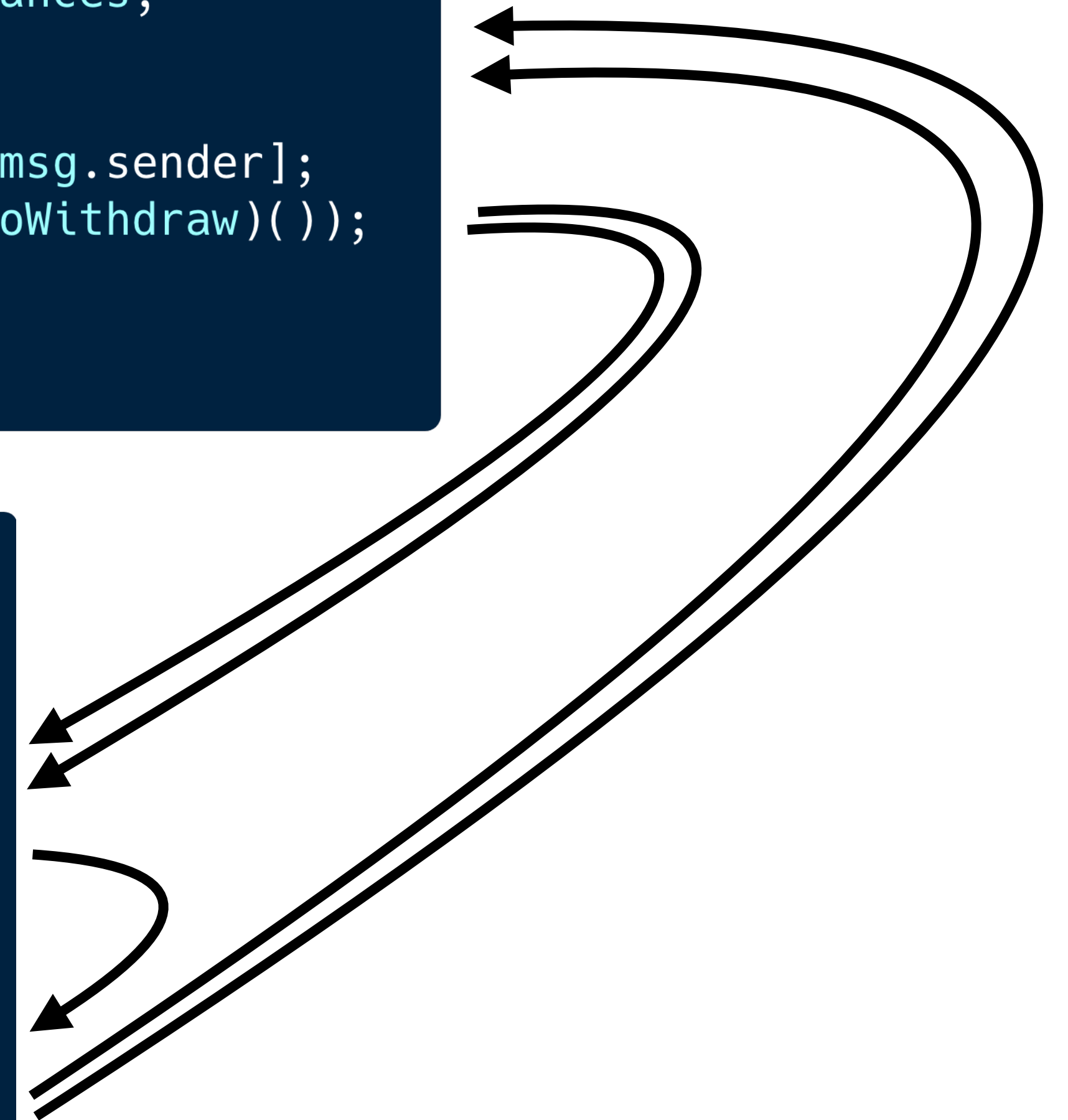
# Reentrancy

```
1 mapping (address => uint) private userBalances;  
2  
3 function withdrawBalance() public {  
4     uint amountToWithdraw = userBalances[msg.sender];  
5     require(msg.sender.call.value(amountToWithdraw)());  
6     userBalances[msg.sender] = 0;  
7 }
```



*pwn()*

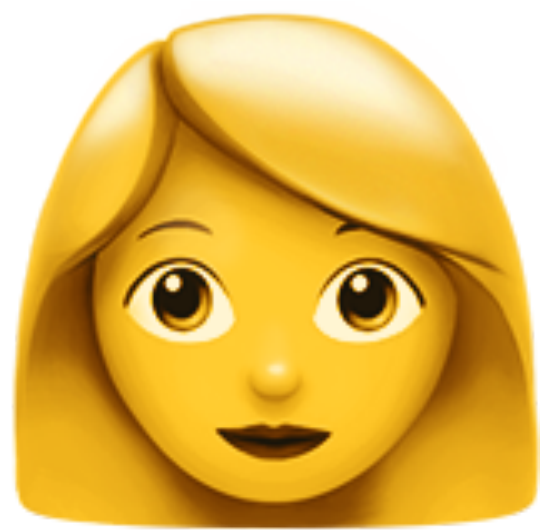
```
1 uint8 toWithdraw = 2;  
2 address target = ...;  
3  
4 function () public payable {  
5     if (--toWithdraw > 0) {  
6         pwn();  
7     }  
8 }  
9  
10 function pwn() public {  
11     target.withdrawBalance()  
12 }
```



# Reentrancy

```
1 mapping (address => uint) private userBalances;  
2  
3 function withdrawBalance() public {  
4     uint amountToWithdraw = userBalances[msg.sender];  
5     require(msg.sender.call.value(amountToWithdraw)());  
6     userBalances[msg.sender] = 0;  
7 }
```

Lesson: Use the  
"checks, effects,  
interactions" pattern!



*pwn()*

```
1 uint8 toWithdraw = 2;  
2 address target = ...;  
3  
4 function () public payable {  
5     if (--toWithdraw > 0) {  
6         pwn();  
7     }  
8 }  
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10 function pwn() public {  
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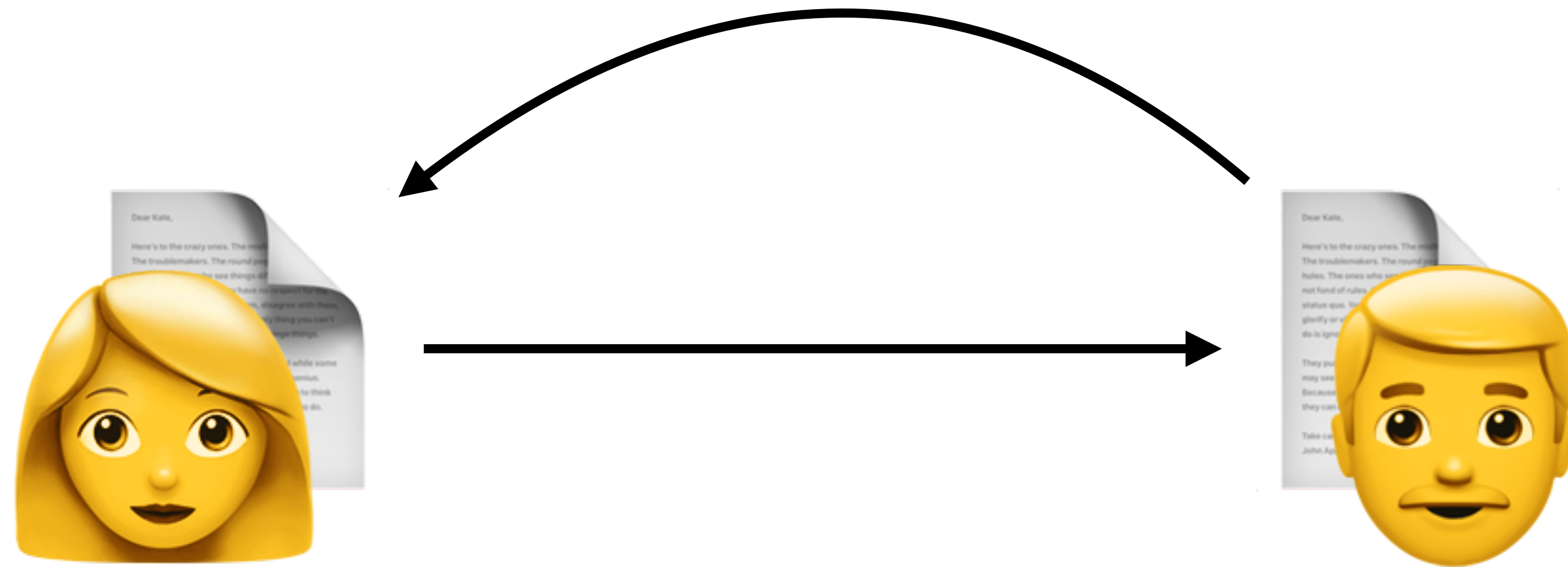
# Malicious External Calls



# Malicious External Calls

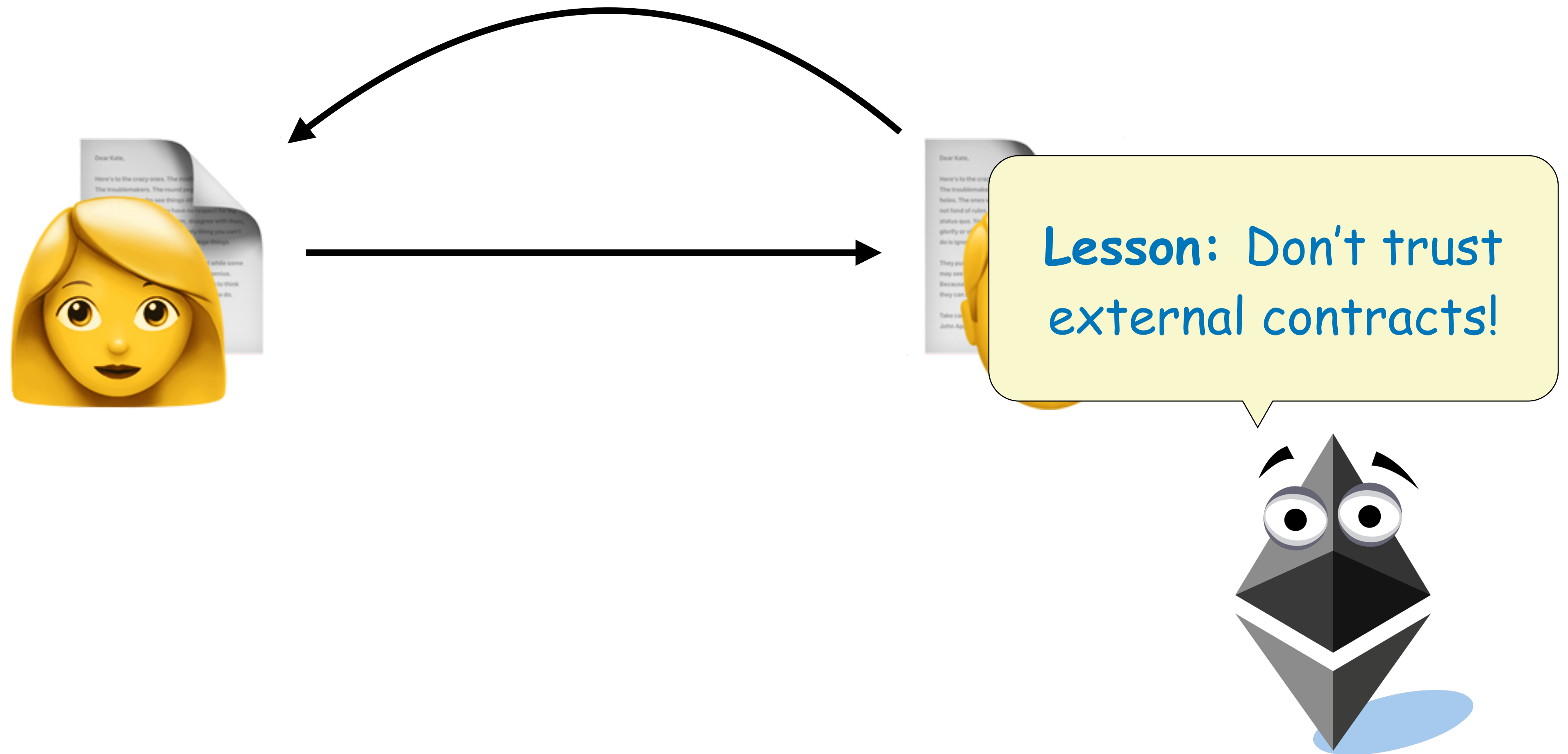


# Malicious External Calls





# Malicious External Calls



# Zero Initialization

```
1 /**
2  * Checks if the address is an admin
3  */
4 modifier onlyAdmin() {
5     require(admins[msg.sender].revokedTimeStamp == 0,
6             'Admin was revoked. ');
7     _;
8 }
```

# Zero Initialization

```
1 /**
2  * Checks if the address is an admin
3  */
4 modifier onlyAdmin() {
5     require(admins[msg.sender].revokedTimeStamp == 0,
6         'Admin was revoked.');
```

```
1 enum UserStatus { Registered, Approved, Denied }
2
3 modifier onlyRegisteredUsers(address userAddress) {
4     require(users[userAddress].status == UserStatus.Registered);
5     _;
6 }
```



# Zero Initialization

```
1 /**
2  * Checks if the address is an admin
3  */
4 modifier onlyAdmin() {
5     require(admins[msg.sender].revokedTimeStamp == 0,
6             'Admin was revoked.');
```

```
1 enum UserStatus { Registered, Approved, Banned }
2
3 modifier onlyRegisteredUsers(address userAddress) {
4     require(users[userAddress].status == UserStatus.Registered);
5     _;
6 }
```

**Lesson:** Unlike in most other languages, uninitialized keys will result in uninitialized memory, which is zeroed.



# Array Length Manipulation

```
1 contract Vulnerable {
2     address public owner = msg.sender;
3     uint256[] map;
4
5     function set(uint256 key, uint256 value) public payable {
6         if (map.length <= key) {
7             map.length = key + 1;
8         }
9
10        map[key] = value;
11    }
12 }
```

# Array Length Manipulation

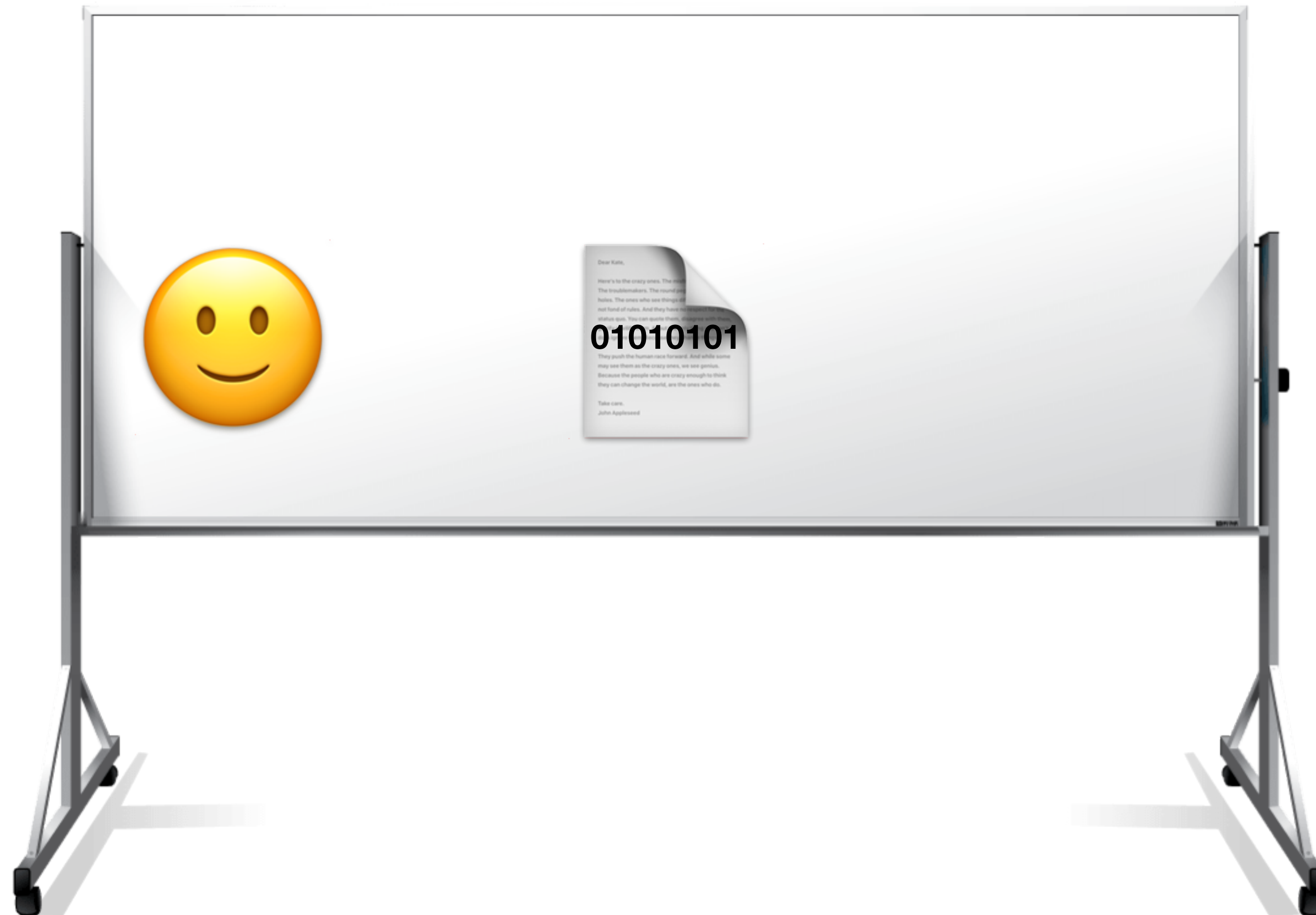
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2     address public owner = msg.sender;  
3     uint256[] map;  
4  
5     function set(uint256 key, uint256 value)  
6         if (map.length <= key) {  
7             map.length = key + 1;  
8         }  
9  
10    map[key] = value;  
11 }  
12 }
```

**Lesson:** Never manually  
manipulate the length  
of an array!

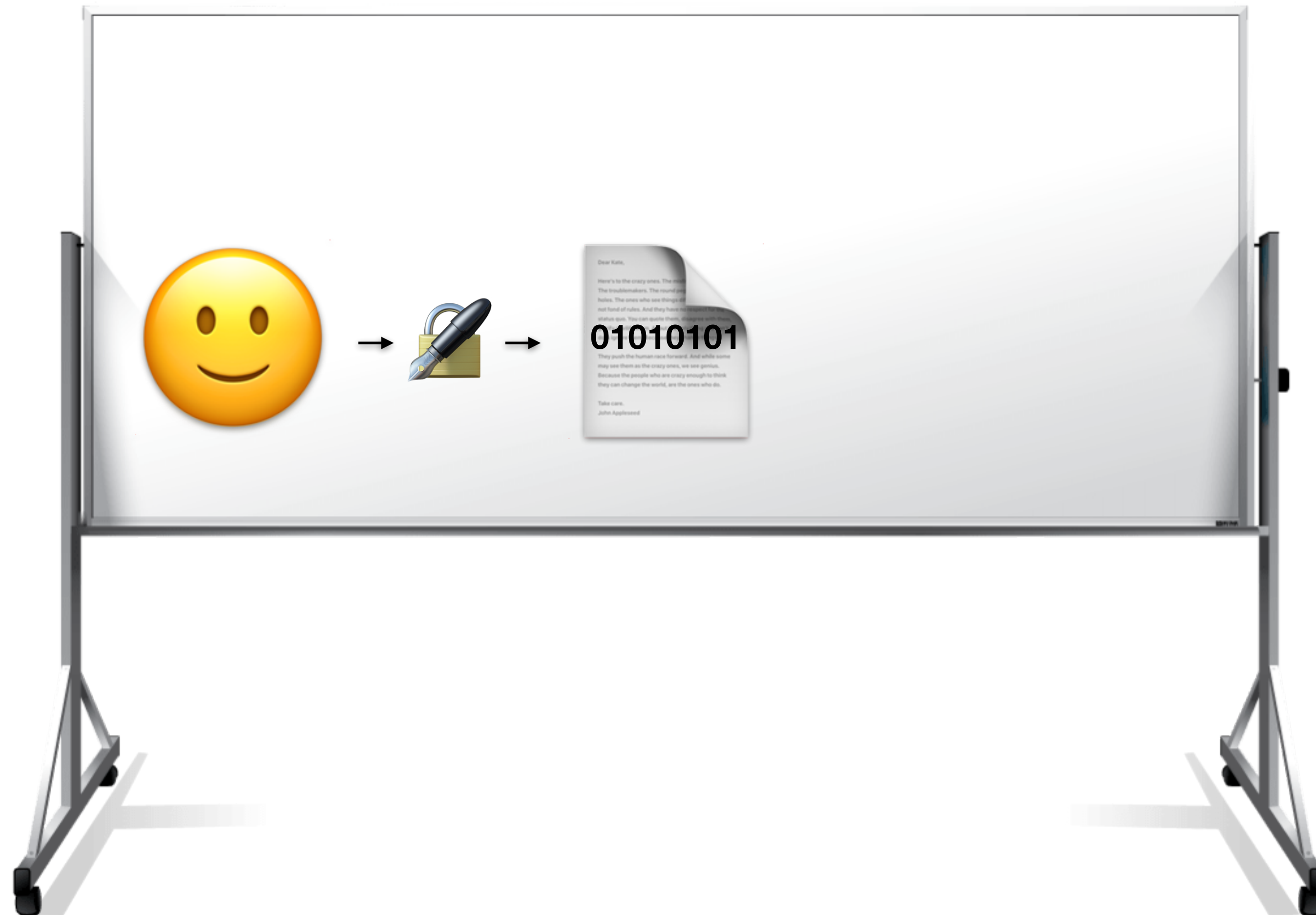




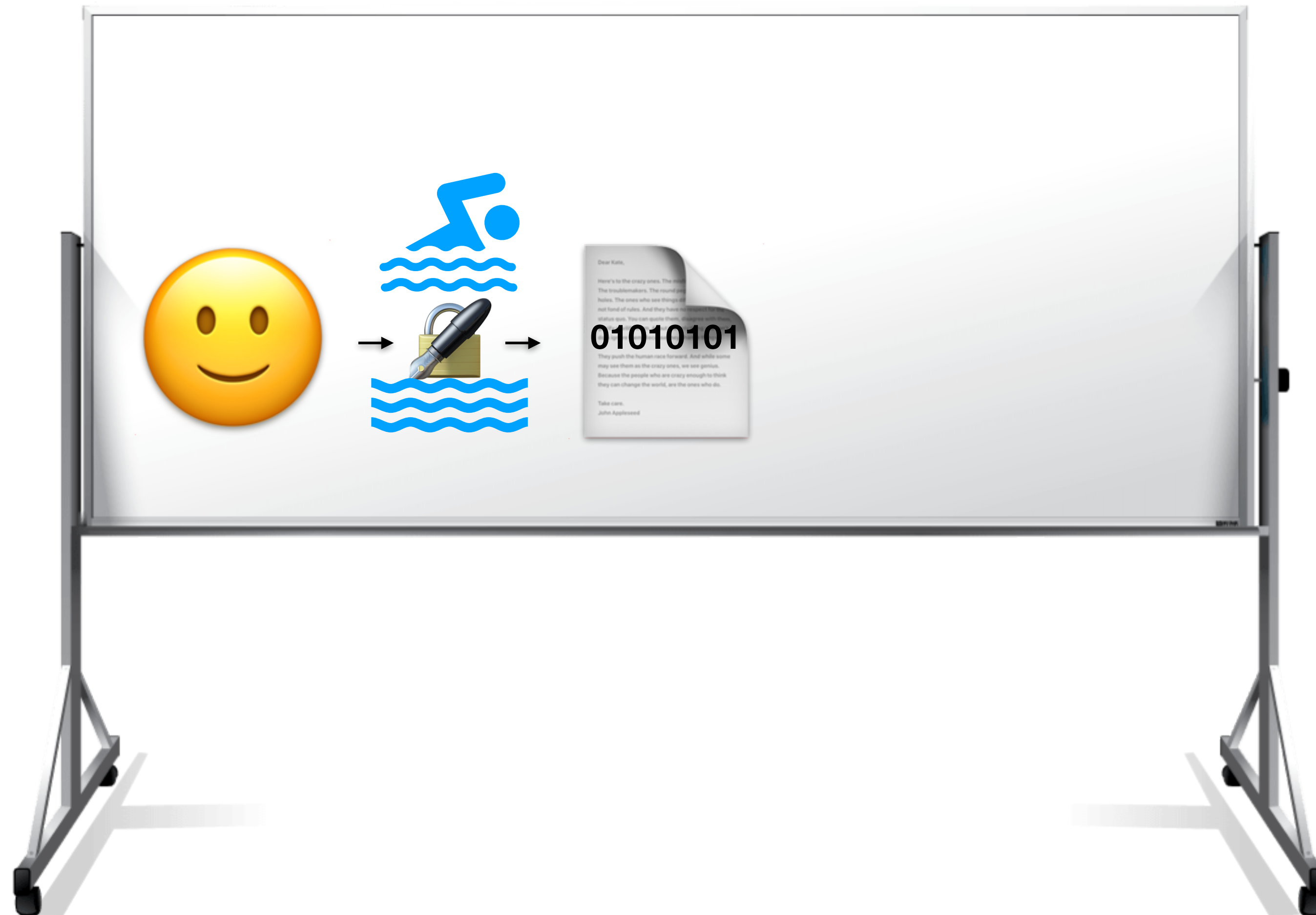
# Transaction “Frontrunning”



# Transaction “Frontrunning”

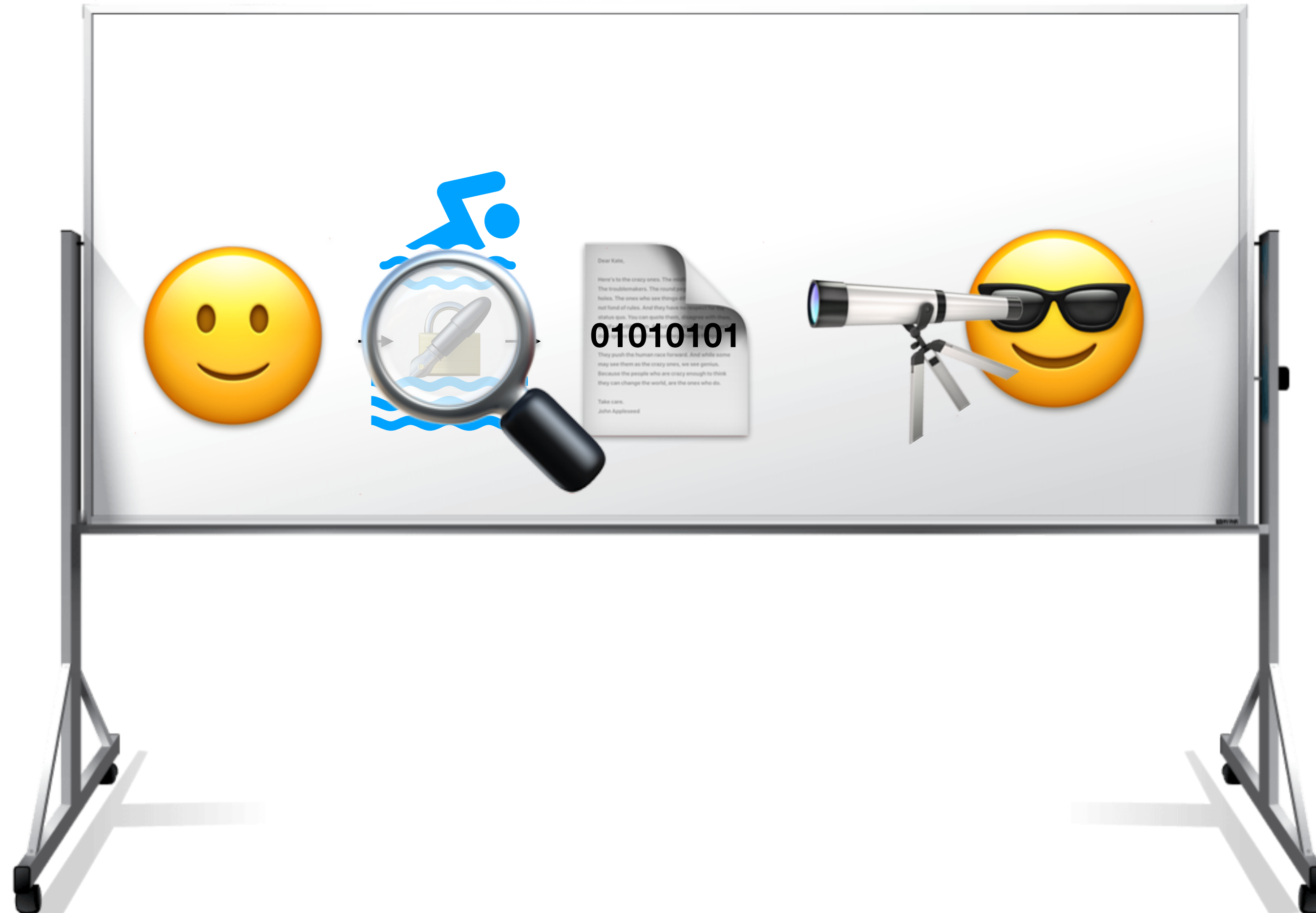


# Transaction “Frontrunning”

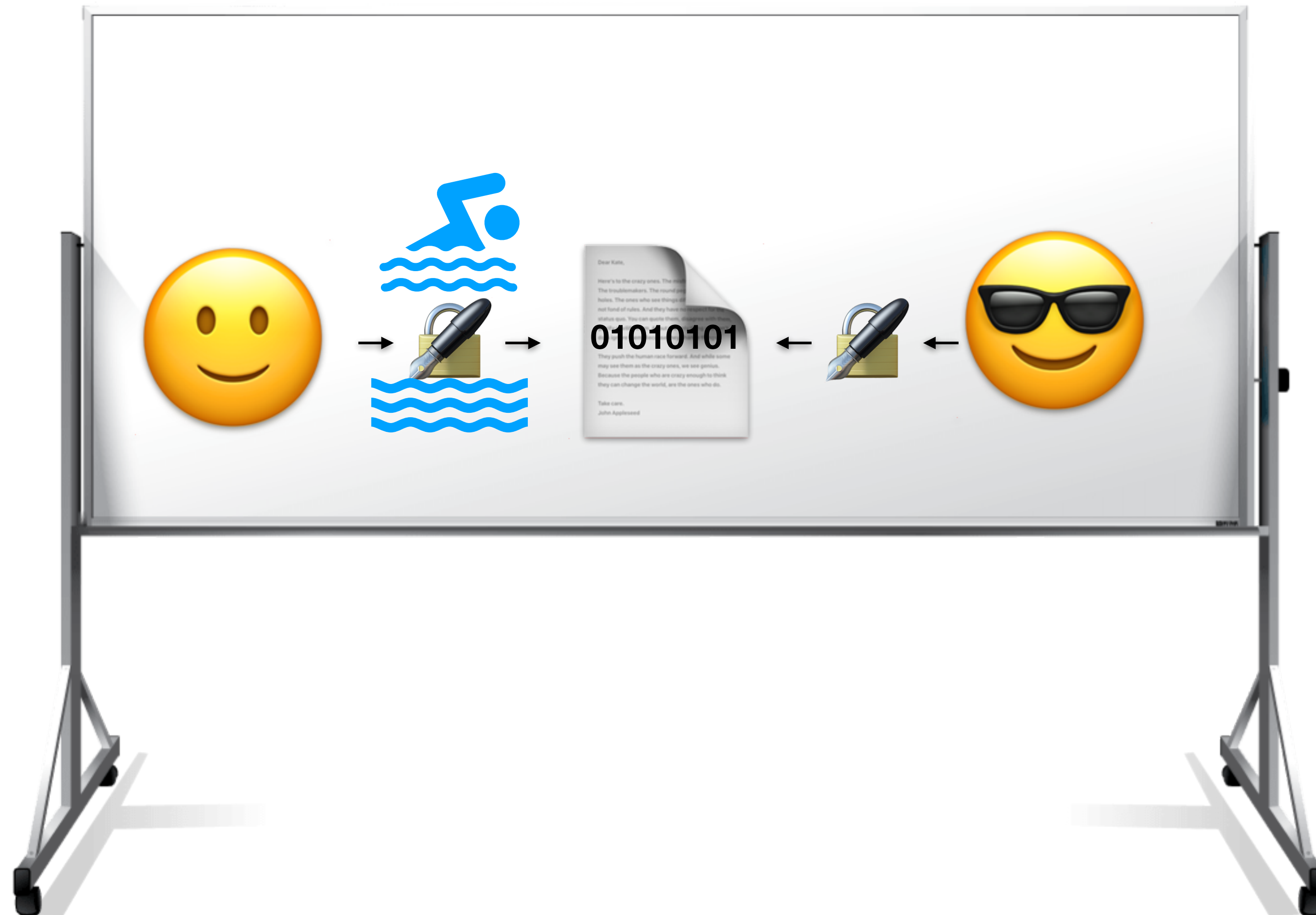




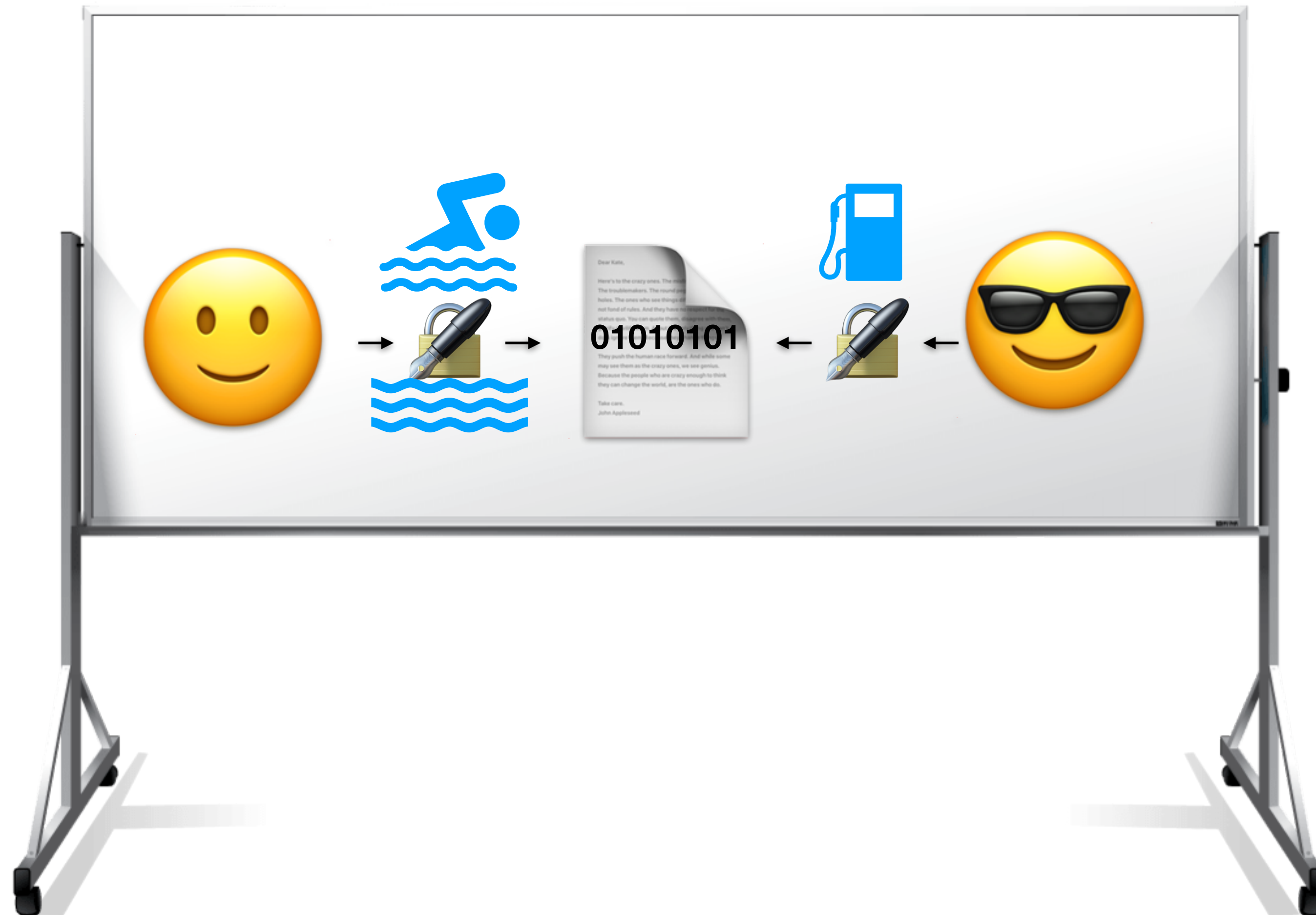
# Transaction “Frontrunning”



# Transaction “Frontrunning”

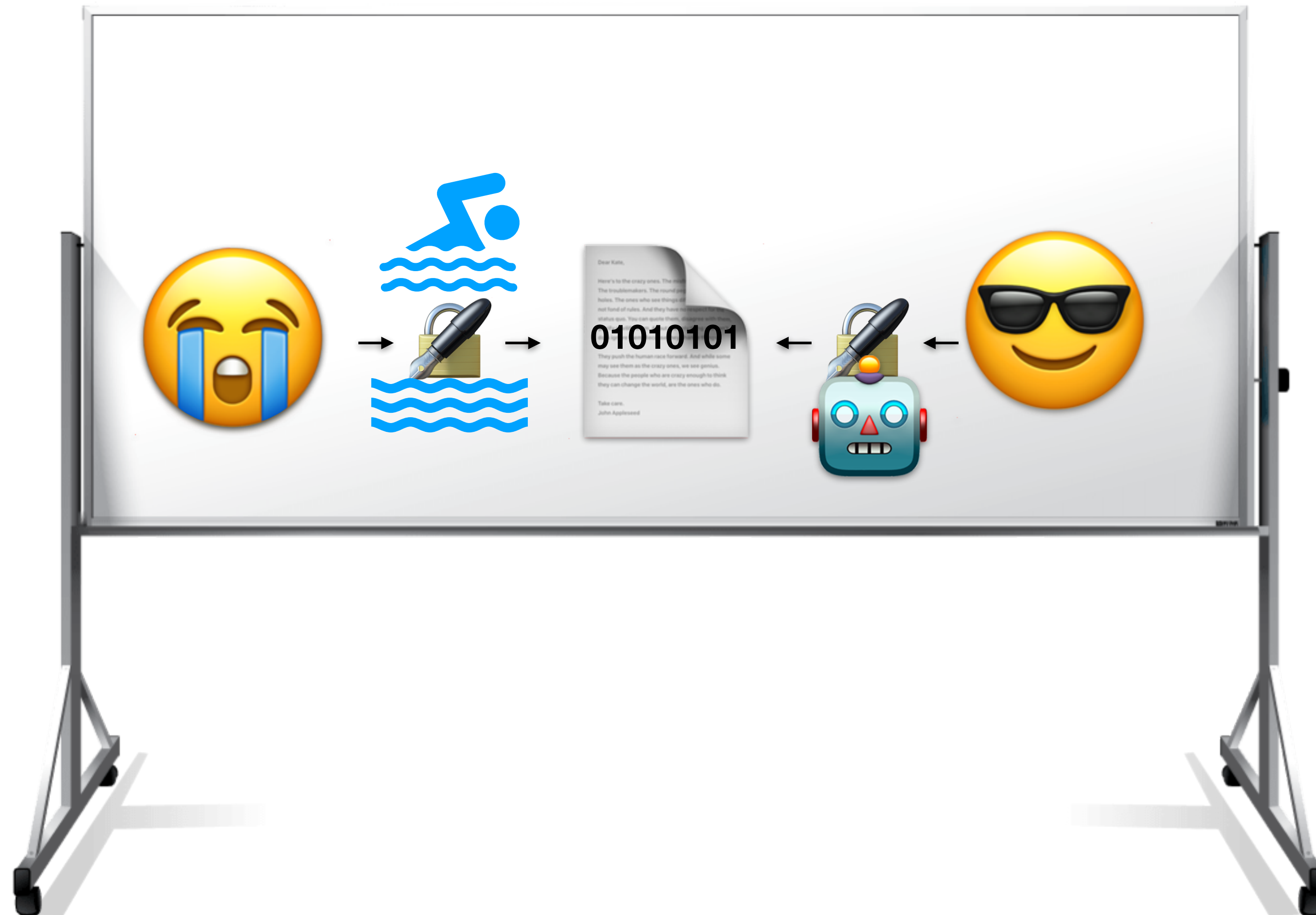


# Transaction “Frontrunning”

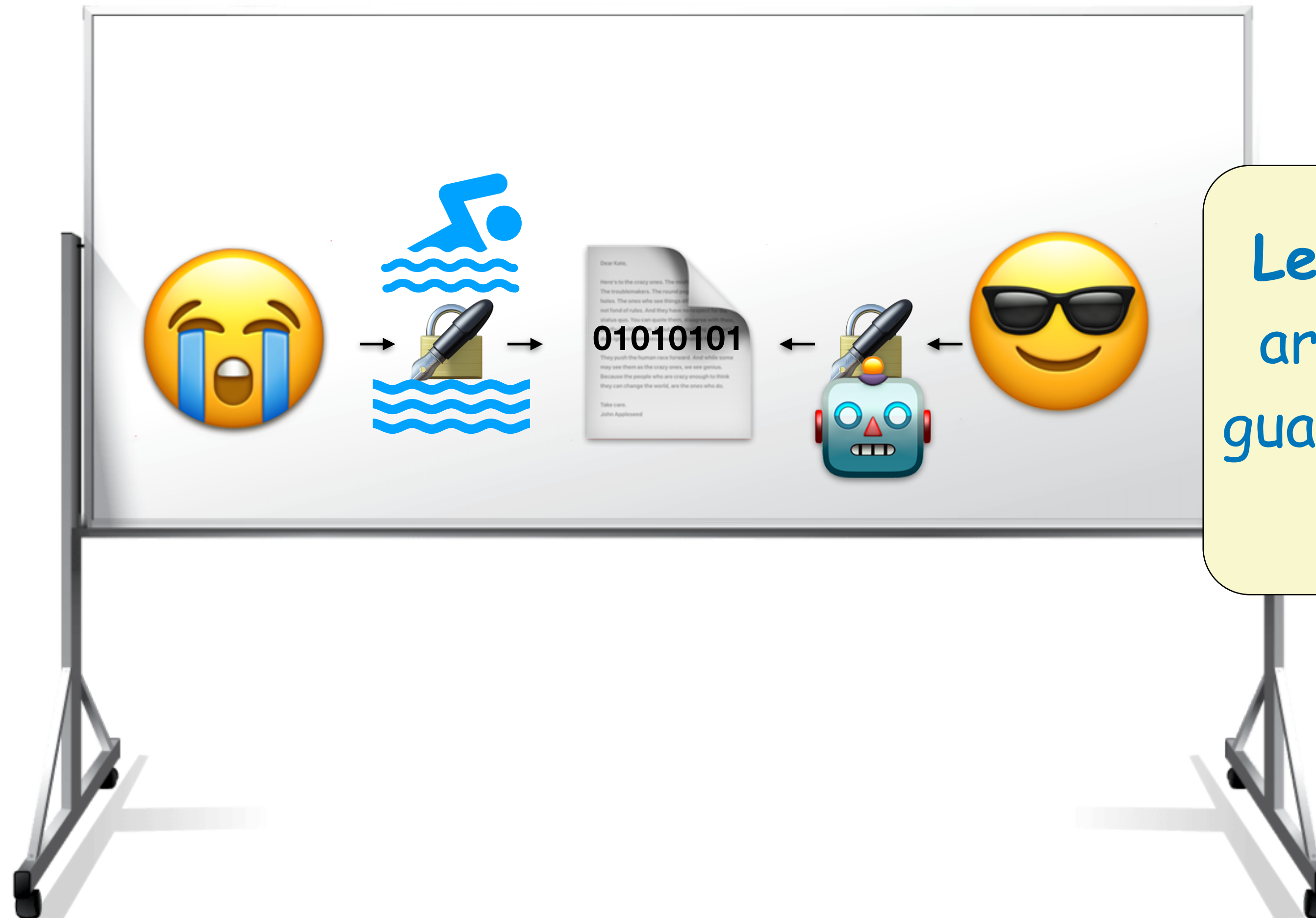




# Transaction “Frontrunning”



# Transaction “Frontrunning”



**Lesson:** Transactions are public, and aren't guaranteed to be mined in order



# Randomness

- The blockchain does not provide any cryptographically secure source of randomness
  - ▶ Block hashes are random, but miners can manipulate them
  - ▶ Miners can also influence timestamps





# Randomness

- **The blockchain does not provide any cryptographically secure source of randomness**
  - Block hashes are random, but miners can manipulate them
  - Miners can also influence timestamps
- **Everything in a contract is publicly visible**
  - Random numbers can't be generated until after all lottery entries have been recorded



# Randomness

- **The blockchain does not provide any cryptographically secure source of randomness**
  - Block hashes are random, but miners can manipulate them
  - Miners can also influence timestamps
- **Everything in a contract is publicly visible**
  - Random numbers can't be generated until after all lottery entries have been recorded
- **Computers will always be faster than the blockchain**
  - Any number a contract could generate can be pre-calculated off-chain faster



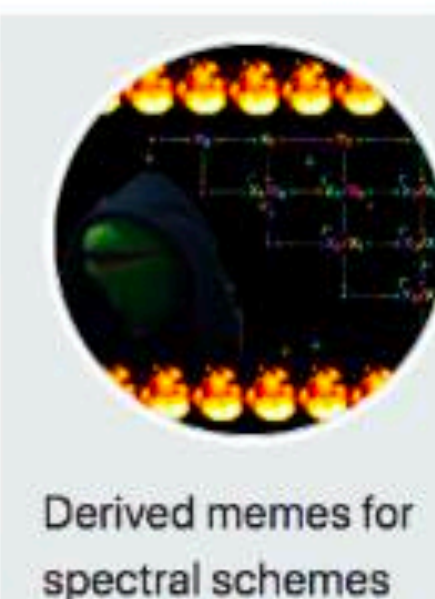
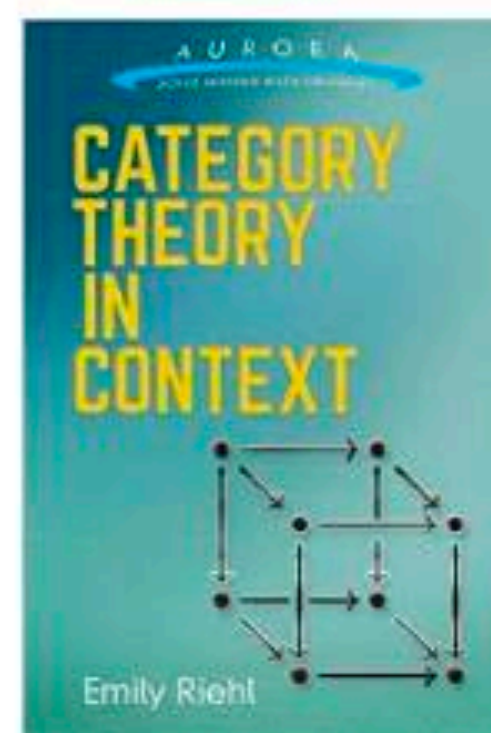


# Don't try to be clever with number theory

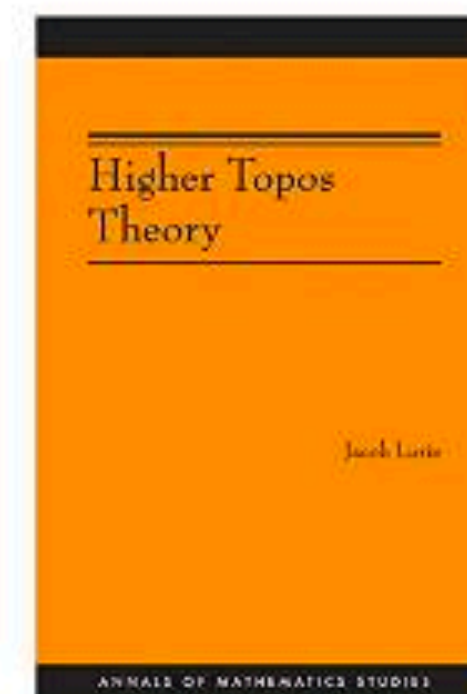
## category theory student starter pack



"\_\_\_ is just \_\_\_"



Derived memes for  
spectral schemes



136  
✓

I find some of this exchange truly depressing. There is a subject of "brave new algebra" and there are myriads of past and present constructions and calculations that depend on having concrete and specific constructions. People who actually compute anything do not use  $(\infty, 1)$  categories when doing so. To lay down a challenge, they would be of little or no use there. One can sometimes use  $(\infty, 1)$  categories to construct things not easily constructed otherwise, and then one can compute things about them (e.g. work of Behrens and Lawson). But the tools of computation are not  $(\infty, 1)$  categorical, and often not even model categorical. People should learn some serious computations, do some themselves, before totally immersing themselves in the formal theory. Note that  $(\infty, 1)$  categories are in principle intermediate between the point-set level and the homotopy category level. It is easy to translate into  $(\infty, 1)$  categories from the point-set level, whether from model categories or from something weaker. Then one can work in  $(\infty, 1)$  categories. But the translation back out to the "old-fashioned" world that some writers seem to imagine expendable lands in homotopy categories. That is fine if that is all that one needs, but one often needs a good deal more. One must be eclectic. Just one old man's view.



Top  $\simeq$   $\infty$ Grpd

"\_\_\_ is just \_\_\_,  
which is just an instance  
of \_\_\_ in the category of \_\_\_"



answered Dec 13 '11 at 2:06  
Peter May  
24.5k 3 76 113



$(\infty, 1)$ -categories

@GrothendieckSchemes




# Don't try to be clever with number theory

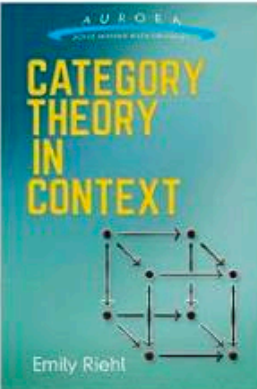
```
winner = entries[blockHash % entries.length];
```

## category theory student starter pack

"\_\_\_ is just \_\_\_"

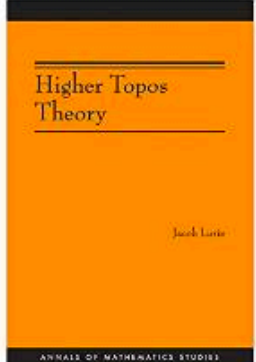


Kerodon



136


I find some of this exchange truly depressing. There is a subject of "brave new algebra" and there are myriads of past and present constructions and calculations that depend on having concrete and specific constructions. People who actually compute anything do not use  $(\infty, 1)$  categories when doing so. To lay down a challenge, they would be of little or no use there. One can sometimes use  $(\infty, 1)$  categories to construct things not easily constructed otherwise, and then one can compute things about them (e.g. work of Behrens and Lawson). But the tools of computation are not  $(\infty, 1)$  categorical, and often not even model categorical. People should learn some serious computations, do some themselves, before totally immersing themselves in the formal theory. Note that  $(\infty, 1)$  categories are in principle intermediate between the point-set level and the homotopy category level. It is easy to translate into  $(\infty, 1)$  categories from the point-set level, whether from model categories or from something weaker. Then one can work in  $(\infty, 1)$  categories. But the translation back out to the "old-fashioned" world that some writers seem to imagine expendable lands in homotopy categories. That is fine if that is all that one needs, but one often needs a good deal more. One must be eclectic. Just one old man's view.



nLab

Top  $\simeq$   $\infty$ Grpd

"\_\_\_ is just \_\_\_,  
which is just an instance  
of \_\_\_ in the category of \_\_\_"



$(\infty, 1)$ -categories

@GrothendieckSchemes

Everybody with me!

You can't do random  
on a blockchain

# Everybody with me!

## You can't do random on a blockchain

**Lesson:** If you really  
need randomness, use a  
trusted off-chain oracle.





# Pre-Signed Transfers

```
1 function signedTransfer(  
2     address to,  
3     uint256 tokens,  
4     address feeRecipient,  
5     uint256 fee,  
6     uint256 expiry,  
7     bytes32 nonce,  
8     uint8 v, bytes32 r, bytes32 s) external returns (bool success) {  
9     bytes32 releaseHash = keccak256(abi.encodePacked(  
10         "\x19\x01",  
11         DOMAIN_SEPARATOR,  
12         keccak256(abi.encode(SIGNEDTRANSFER_TYPEHASH, to, tokens, feeRecipient, fee, expiry, nonce))  
13     ));  
14  
15     address from = ecrecover(releaseHash, v, r, s);  
16  
17     approvals[from][msg.sender] = add(tokens, fee);  
18  
19     transferFrom(from, to, tokens);  
20     transferFrom(from, feeRecipient, fee);  
21  
22     return true;  
23 }
```

# Pre-Signed Transfers

```
1 function signedTransfer(  
2     address to,  
3     uint256 tokens,  
4     address feeRecipient,  
5     uint256 fee,  
6     uint256 expiry,  
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9     bytes32 releaseHash = keccak256(abi.encodePacked(  
10         "\x19\x01",  
11         DOMAIN_SEPARATOR,  
12         keccak256(abi.encode(SIGNEDTRANSFER_TYPEHASH, to, tokens, f  
13     ));  
14  
15     address from = ecrecover(releaseHash, v, r, s);  
16  
17     approvals[from][msg.sender] = add(tokens, fee);  
18  
19     transferFrom(from, to, tokens);  
20     transferFrom(from, feeRecipient, fee);  
21  
22     return true;  
23 }
```

**Lesson:** Always check the  
return value of ecrecover!  
Better yet, avoid it!



What Can You Do  
About It?

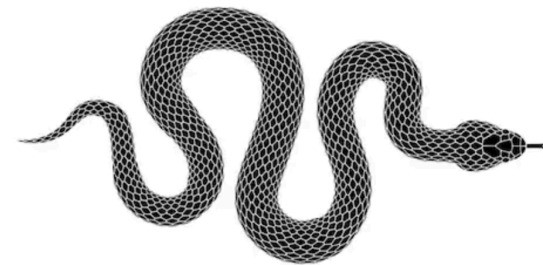


# What can be done?

<https://github.com/trailofbits/...>

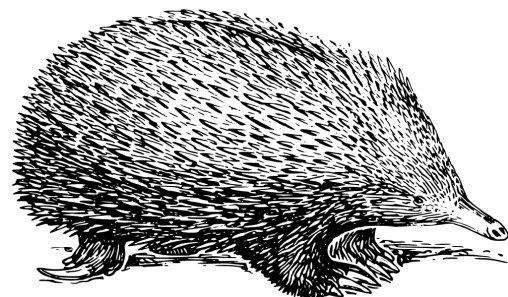


Manticore Symbolic Execution

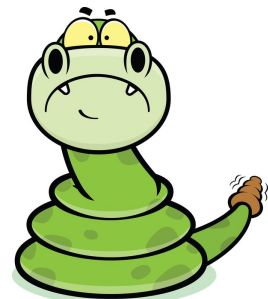


Slither Static Analysis

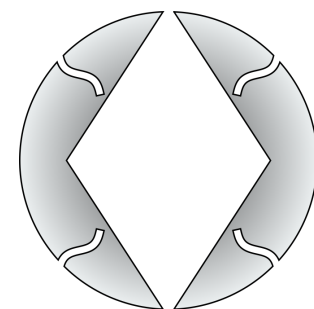
**evm-opcodes**  
VM Reference



Echidna Property Based Fuzzer



Rattle EVM to SSA Lifter



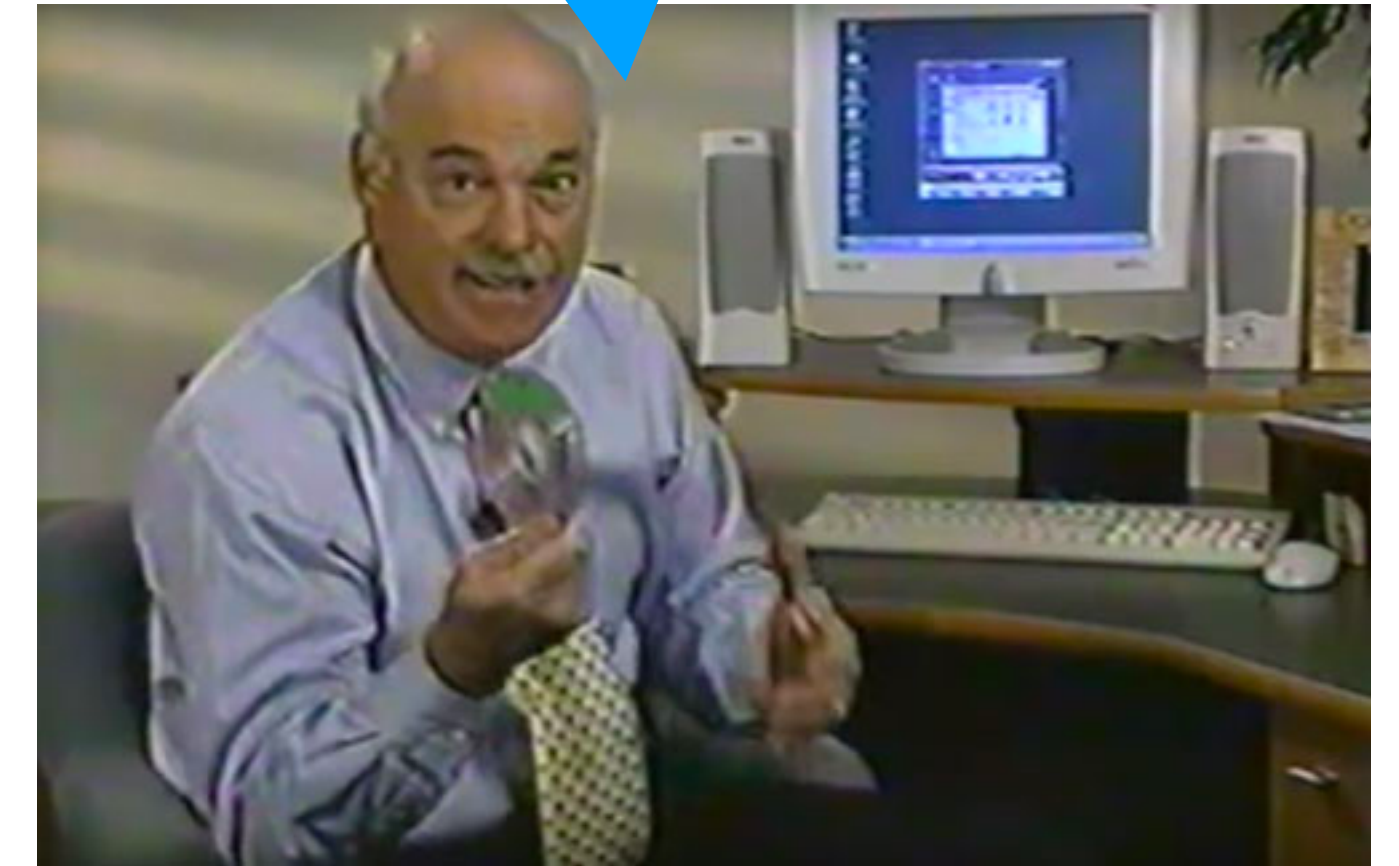
Etheno Test Framework Integration

Ethersplay Visual EVM Disassembler

pyevmasm Bytecode Analysis



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open source products.



**blockchain-security-contacts**  
it's surprisingly hard to disclose bugs

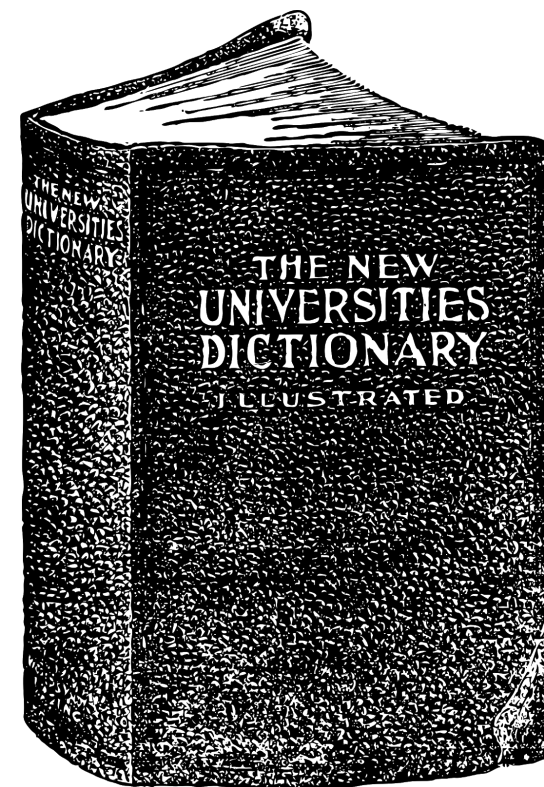
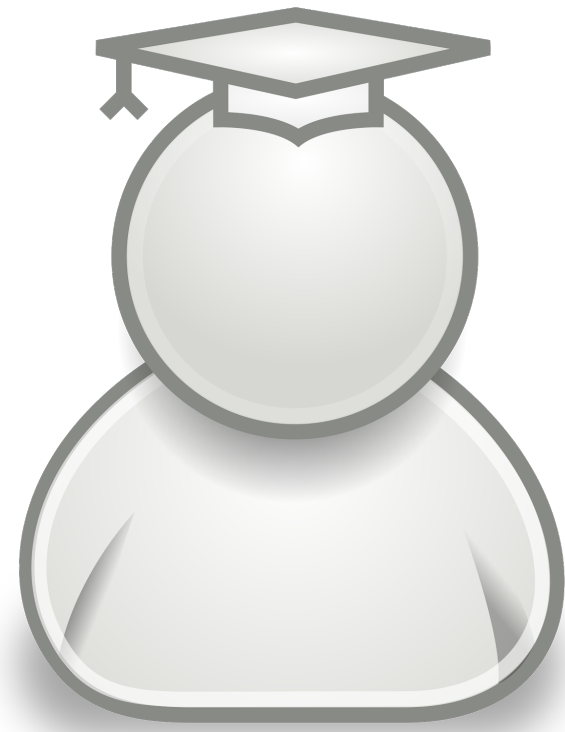
**not-so-smart-contracts**  
common vulnerability database  
**awesome-ethereum-security**

security best practices

# (Not So) Smart Contracts

## Educational Tool

Learn about EVM and Solidity Vulnerabilities



## Reference Material

Useful when Auditing Code

## Working Examples of Contracts

Real Vulnerabilities Found in the Wild



<https://github.com/crytic/not-so-smart-contracts>

# Community Information

## Awesome Ethereum Security

- **What?** Curated list of community-maintained and open-source references
- **Why?** Everything in one place: no more searching through stack overflow, github, and reddit
- **Features**
  - ▶ Resources for secure development, CTFs & wargames, and even specific podcast episodes
  - ▶ Identifies security tools for visualization, linting, bug finding, verification, and reversing
  - ▶ Pointers to related communities

## Blockchain Security Contacts

- **What?** Comprehensive list of security contacts for blockchain applications
- **Why?** Projects worth \$10MM+ should have a way to engage with security researchers
- **Features**
  - ▶ Vuln disclosure program best practices
  - ▶ Deployed addresses template for dapps
  - ▶ Existing contact info for over 100 projects (Blockchains, dapps, ERC20 and 721 tokens, Exchanges, Wallet software)

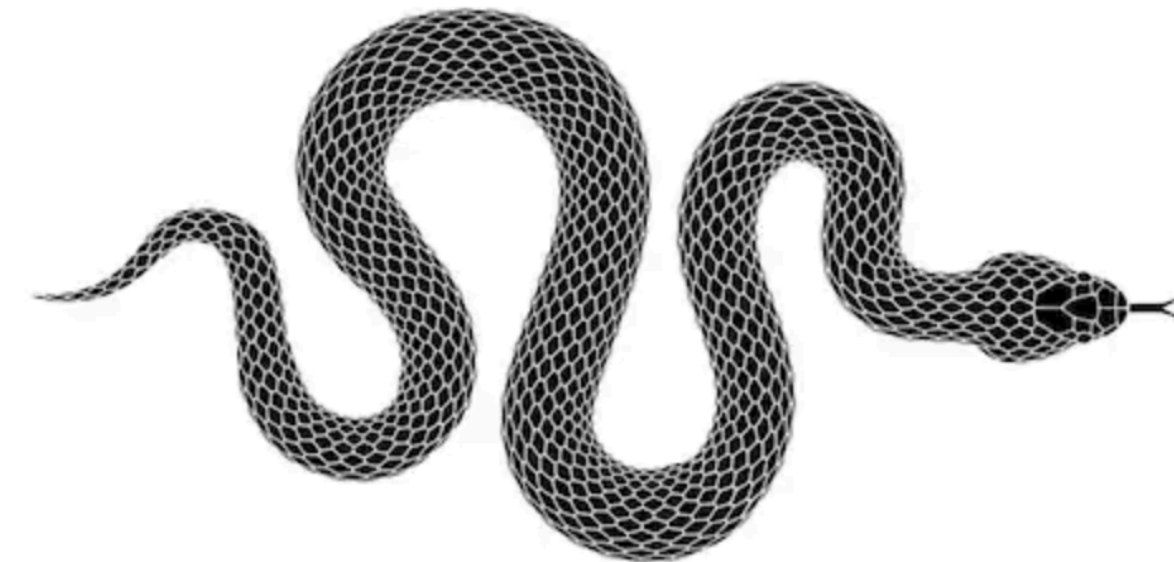
<https://github.com/crytic/awesome-ethereum-security> and [/blockchain-security-contacts](https://github.com/crytic/blockchain-security-contacts)



# Slither

## Smart Contract Static Analysis

- Solidity and Vyper vulnerability detection
  - Low false positives
  - Easily integrates into CI pipeline
  - Very fast (milliseconds)
  - Supports advanced value- and taint-tracking
  - Python-based detector API
- **Inputs:** Solidity code
  - **Outputs:**
    - Detected errors (extensive list of vulnerability detectors included)
    - Warnings of poor coding practices
    - Inheritance graph and contract summary



Slither is open source!

<https://github.com/crytic/slither>

# Slither Installation and Usage

```
$ pip3 install slither-analyzer
```

then

```
$ slither contract.sol
```

or

```
$ truffle compile; slither .
```

That's literally it!

# Slither Installation and Usage

```
$ pip3 install slither-analyzer
```

then

```
$ slither contract.sol
```

or

```
$ truffle compile; slither .
```

That's literally it!

**Lesson:** Slither is super easy and quick! No excuse not to integrate it in your CI pipeline.





# Problem: Test for New Bugs

```
contract Simple {  
    function f(uint a){  
        // .. lot of paths and conditions  
  
        if (a == 65) {  
            // leads to a bug here  
        }  
    }  
}
```

# Problem: Test for New Bugs

```
contract Simple {  
    function f(uint a){  
        // .. lot of paths and  
  
        if (a == 65) {  
            // leads to a bug h  
        }  
    }  
}
```

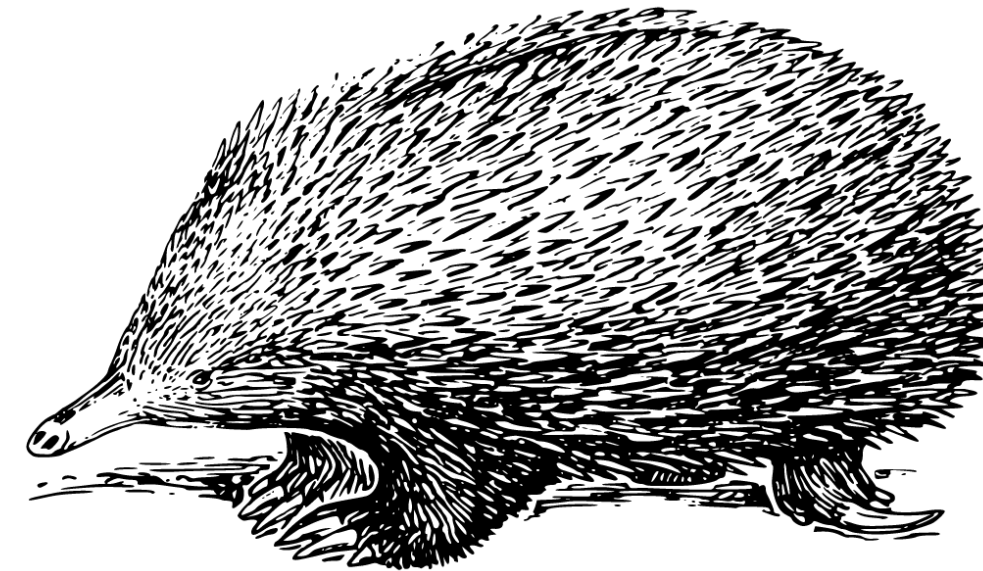
It looks like you want to detect classes of bugs that have never been seen before!



# Echidna

## Smart Contract Property Tester

- Generates and execute many contract inputs
  - Generate intelligent, grammar-based inputs
  - Seamlessly integrate into developer workflows
  - Run thousands of generated inputs per second
  - Automatically generate minimal testcases
  - Highly extensible via Haskell API
- **Inputs:** Solidity code and tests
  - **Outputs:**
    - ▶ List of invariants Echidna was able to violate
    - ▶ Minimal call sequence to trigger discovered violations



Echidna is open source!

<https://github.com/crytic/echidna>



# Echidna Example

```
→ echidna git:(master) ✗ cat solidity/cli.sol
pragma solidity ^0.4.16;

contract Test {
    bool private flag0=true;
    bool private flag1=true;

    function set0(int val) returns (bool){
        if (val % 10 == 0) {flag0 = false;}
    }
    function set1(int val) returns (bool){
        if (val % 10 == 0 && flag0) {flag1 = false;}
    }
    function echidna_alwaystrue() returns (bool){
        return(true);
    }
    function echidna_sometimesfalse() returns (bool){
        return(flag0 || flag1);
    }
}

→ echidna git:(master) ✗ ./echidna-test solidity/cli.sol
— solidity/cli.sol —
✗ "echidna_sometimesfalse" failed after 36 tests and 681 shrinks.

    | Call sequence: set0(7946810797001355118938603703351564369838113269809310950469780);
    |                  set1(8045329803519652513052969161362647695379403994810754718464019950667760);

✓ "echidna_alwaystrue" passed 100 tests.
✗ 1 failed, 1 succeeded.
→ echidna git:(master) ✗
```

# Echidna Example

```
→ echidna git:(master) ✗ cat solidity/cli.sol
pragma solidity ^0.4.16;

contract Test {
    bool private flag0=true;
    bool private flag1=true;

    function set0(int val) returns (bool){
        if (val % 10 == 0) {flag0 = false;}
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    | Call sequence: set0(7946810797001355118938603703351564369838113269809310950469780);
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✓ "echidna_alwaystrue" passed 100 tests.
✗ 1 failed, 1 succeeded.
→ echidna git:(master) ✗
```

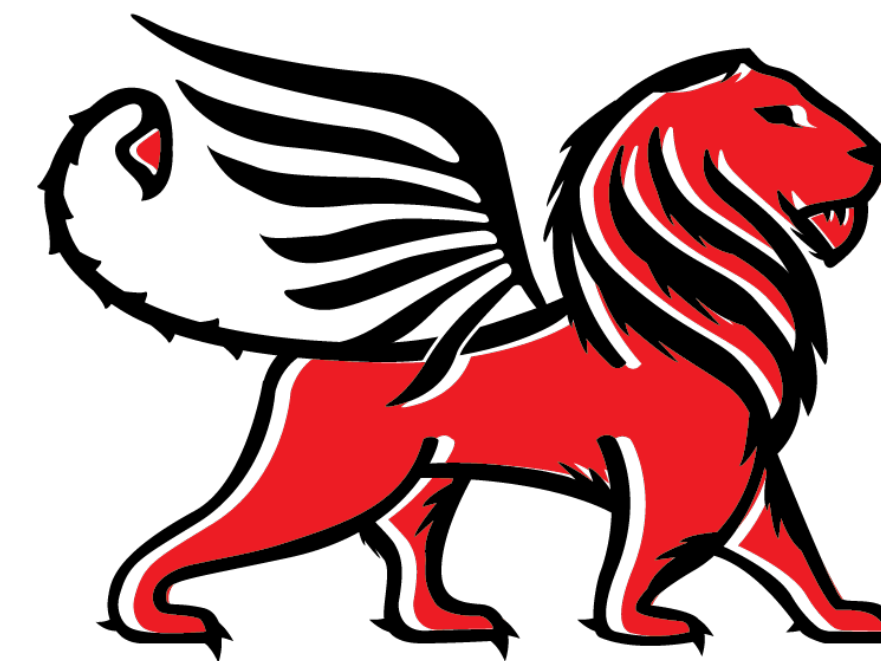
**Lesson:** Echidna is not as fast as Slither, but it is still quick enough to be useful in your CI pipeline. Unlike Slither, it is capable of discovering wholly new classes of bugs.



# Manticore

## Smart Contract Verifier

- Uses symbolic execution of EVM
  - Deeply explores possible contract states across multiple transactions and contracts
  - Discover functions directly from bytecode
  - Detect contract flaws like int overflows, uninitialized memory/storage usage, and more
  - Verify customized program assertions
  - Highly scriptable and extensible via Python API
- **Inputs:** Solidity code (optional) or raw EVM bytecode
  - **Outputs:**
    - List of detected flaws
    - Verified properties
    - Execution traces of discovered paths



Manticore is open source!

<https://github.com/trailofbits/manticore>



# Manticore Example

```
contract Simple {  
    function f(uint a){  
        // .. lot of paths and conditions  
  
        if (a == 65) {  
            revert();  
        }  
    }  
}
```

# Manticore Example

```
contract Simple {  
    function f(uint a){  
        // .. lot of paths and conditions  
  
        if (a == 65) {  
            revert();  
        }  
    }  
}
```

```
$ manticore simple.sol
```

```
2018-02-28 17:06:21,650: [25981] m.main:INFO: Beginning analysis
```

```
2018-02-28 17:06:21,803: [25981] m.ethereum:INFO: Starting symbolic transaction: 1
```

```
2018-02-28 17:06:22,098: [25981] m.ethereum:INFO: Generated testcase No. 0 - REVERT
```

```
2018-02-28 17:06:23,185: [25981] m.ethereum:INFO: Generated testcase No. 1 - REVERT
```

```
2018-02-28 17:06:24,206: [25981] m.ethereum:INFO: Finished symbolic transaction: 1 | Code Coverage: 100% |
```

```
Terminated States: 3 | Alive States: 1
```

```
2018-02-28 17:06:24,213: [32058] m.ethereum:INFO: Generated testcase No. 2 - STOP
```

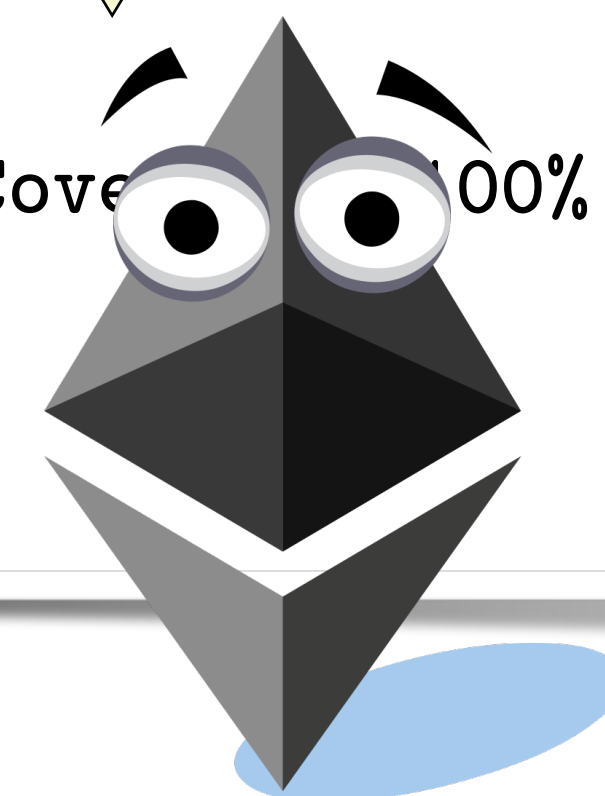
```
2018-02-28 17:06:25,269: [25981] m.ethereum:INFO: Results in /examples/mcore_zua0Yl
```

# Manticore Example

```
contract Simple {  
    function f(uint a){  
        // .. lot of paths and conditions  
  
        if (a == 65) {  
            revert();  
        }  
    }  
}
```

**Manticore can verify that  
your code satisfies its  
invariants, but it can take a  
long time to run!**

```
$ manticore simple.sol  
2018-02-28 17:06:21,650: [25981] m.main:INFO: Beginning analysis  
2018-02-28 17:06:21,803: [25981] m.ethereum:INFO: Starting symbolic t  
2018-02-28 17:06:22,098: [25981] m.ethereum:INFO: Generated testcase No. 0 - REVERT  
2018-02-28 17:06:23,185: [25981] m.ethereum:INFO: Generated testcase No. 1 - REVERT  
2018-02-28 17:06:24,206: [25981] m.ethereum:INFO: Finished symbolic transaction: 1 | Code Cover 100% |  
Terminated States: 3 | Alive States: 1  
2018-02-28 17:06:24,213: [32058] m.ethereum:INFO: Generated testcase No. 2 - STOP  
2018-02-28 17:06:25,269: [25981] m.ethereum:INFO: Results in /examples/mcore_zua0Yl
```





# Conclusions

- Solidity isn't a great language, but we're stuck with it (for now)
- Don't assume Solidity behaves like a "normal" language
- Don't trust the Solidity documentation; the sole compiler implementation is canon
- Don't enable Solidity compiler optimizations
- Avoid the "DELEGATECALL" upgrade pattern
- Don't trust calls to external contracts
- Remember that everything on the blockchain is public
- Don't assume transactions will be mined in order (or at all!)
- Read "(Not So) Smart Contracts"
- Add Slither and Echidna into your CI pipeline
- Use Manticore to verify the correctness of your contracts



Thanks!

@ESultanik



# Thanks!

@ESultanik



**TRAIL**  
*OF*  
**BITS**

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